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Logical Equivalence

CMSC250

Spring 2020

Equivalences

• Let's observe the following truth table

p	q	$p \wedge q$	$q \wedge p$
F	F	F	F
F	Т	F	F
T	F	F	F
T	Т	Т	Т

Equivalences

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	•		•

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• Conclusion: $p \land q \equiv q \land p$

This symbol means "logical equivalence"

• Please fill – in the following truth table:

p	q	$p \implies q$	$(\sim p) \vee q$
F	F	?	?
F	Т	?	?
T	F	?	?
T	Т	?	?

$$\Leftrightarrow$$
 vs \equiv

- \Leftrightarrow ("if and only if") is used to form statements, e.g.
 - $p \Leftrightarrow (q \land (\sim r))$
- = ("logically equivalent to") compares two statements, e.g
 - $(p \land q) \equiv (q \land p)$

• Let's fill in the following truth table :

a	b	$\sim (a \wedge b)$	$(\sim a) \lor (\sim b)$
F	F	?	?
F	Τ	?	?
T	F	?	?
Т	Т	?	?

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а	b	$\sim (a \wedge b)$	$(\sim a) \lor (\sim b)$
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This result is known as

De Morgan's law

• Conclusion:
$$\sim (a \land b) \equiv (\sim a) \lor (\sim b)$$

Understanding De Morgan's Law

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• (~"Alice is Blonde") \((~"Alice wears Green Dress"):
Also true!



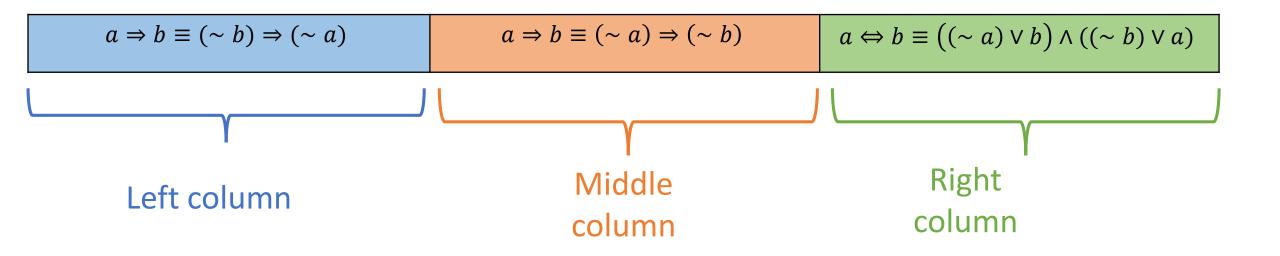
De Morgan's Laws (there's two of them)

$$\sim (a \lor b) \equiv (\sim a) \land (\sim b)$$

$$\sim (a \land b) \equiv (\sim a) \lor (\sim b)$$

- Conjunctions flipped to disjunctions, and vice versa
- Negation operator (~) distributed across terms
- These laws give us our first pair of equivalent expressions!

Are these correct equivalences?



• How do we prove an equivalence? (e.g $\sim (a \land b) \equiv (\sim a) \lor (\sim b)$)

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- Can we do better?

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1. Truth tables

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2. Laws of logical equivalence in a chain, one after the other!

- We no longer have to compare 2^n input combinations to ensure that they all map to the same truth value (**T** or **F**). ©
- But somebody needs to code the system up!

Table of equivalences

Commutativity of binary operators	$p \wedge q \equiv q \wedge p$	$p \lor q \equiv q \lor p$
Associativity of binary operators	$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	$(p \lor q) \lor r \equiv p \lor (q \lor r)$
Distributivity of binary operators	$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	$p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$
Identity laws	$p \wedge T \equiv p$	$p \lor F \equiv p$
Negation laws	$p \lor (\sim p) \equiv T$	$p \wedge (\sim p) \equiv F$
Double negation	$\sim (\sim p) \equiv p$	
Idempotence	$p \wedge p \equiv p$	$p \vee p \equiv p$
De Morgan's axioms	$\sim (p \land q) \equiv (\sim p) \lor (\sim q)$	$\sim (p \lor q) \equiv (\sim p) \land (\sim q)$
Universal bound laws	$p \lor T \equiv T$	$p \wedge F \equiv F$
Absorption laws	$p \lor (p \land q) \equiv p$	$p \land (p \lor q) \equiv p$
Negations of	${\scriptscriptstyle \sim} F \equiv T$	$\sim T \equiv F$
contradictions / tautologies		
Contrapositive	$(a \Rightarrow b) \equiv ($	$(\sim b) \Rightarrow (\sim a))$
Equivalence between	$a \Leftrightarrow b \equiv (a \Rightarrow b) \land (b \Rightarrow a)$	
biconditional and		
implication		
Equivalence between	$a \Rightarrow b \equiv {\scriptstyle \sim} a \lor b$	
implication and		
disjunction		
	l .	

• This exact table will be **given to you** during **all** exams where you might need it, so that you don't have to remember some "exotic" names

Proving equivalences using laws

Suppose we want to investigate if

$$(((a \land b) \lor q) \land (b \land a)) \equiv (p \lor \sim p) \land ((a \land b) \lor ((\sim r) \land r))$$

How many rows would the truth table have?

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 - $2^5 = 32 \odot$ Too much time!

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- How many rows would the truth table have?
 - $2^5 = 32 \otimes \text{Too much time!}$
- Let's see how we could use the laws of logical equivalence to prove this equivalence
 - Important: document the laws!

More equivalences

Let's prove the following equivalences true or false together.

$$a\Rightarrow b\equiv (\sim b)\Rightarrow (\sim a)$$
 (Contrapositive) $a\Rightarrow b\equiv (\sim a)\Rightarrow (\sim b)$ (Inverse Error) $a\Leftrightarrow b\equiv ((\sim a)\vee b)\wedge ((\sim b)\vee a)$

Simplifying expressions

- Large expressions can often be **simplified** using the equivalences we discussed earlier.
- Example: Let's simplify $p \land (p \lor q) \land (p \land q)$

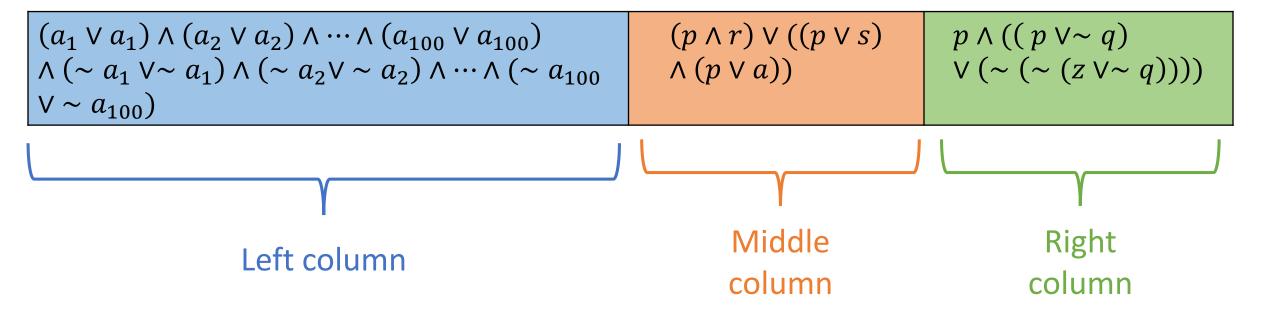
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Here's one way $p \land (p \lor q) \land (p \land q) \text{ (Original expression)}$ $\equiv p \land (p \land q) \text{ (How?)}$ $\equiv (p \land p) \land q \text{ (How?)}$ $\equiv p \land q \text{ (How?)}$

Your turn, class!

• Let's simplify the following three expressions.



Solution to 1

$$(a_1 \vee a_1) \wedge (a_2 \vee a_2) \wedge \cdots \wedge (a_{100} \vee a_{100}) \wedge (\sim a_1 \vee \sim a_1) \wedge (\sim a_2 \vee \sim a_2) \wedge \cdots \wedge (\sim a_{100} \vee \sim a_{100})$$

$$\equiv a_1 \wedge a_2 \wedge \cdots \wedge (a_{100}) \wedge (\sim a_1) \wedge (\sim a_2) \wedge \cdots \wedge (\sim a_{100})$$

$$\equiv a_1 \wedge (\sim a_1) \wedge a_2 \wedge (\sim a_2) \wedge \cdots \wedge (a_{999}) \wedge (\sim a_{999}) \dots \wedge (a_{100}) \wedge (\sim a_{100})$$

$$\equiv F \wedge F \wedge \cdots \wedge F \dots \wedge F$$

$$(Negation 100 times)$$

$$\equiv F$$

$$(Idempotence 99 times)$$

Solution to 2

$$(p \land r) \lor ((p \lor s) \land (p \lor a))$$

$$\equiv (p \wedge r) \vee (p \vee (s \wedge a))$$

$$\equiv ((p \land r) \lor p) \lor (s \land a)$$

$$\equiv (p \lor (p \land r)) \lor (s \land a)$$

$$\equiv p \vee (s \wedge a)$$

(Distributivity)

(Associativity)

(Commutativity)

(Absorption)

Solution to 3

$$p \land ((p \lor \sim q) \lor (\sim (\sim (z \lor \sim q))))$$

$$\equiv p \land ((p \lor \sim q) \lor (z \lor \sim q))$$

$$\equiv p \land ((p \lor z) \lor (\sim q \lor \sim q))$$

$$\equiv p \land ((p \lor z) \lor \sim q)$$

$$\equiv p \land (p \lor (z \lor \sim q))$$

$$\equiv p$$

(Double Negation)

(Associativity)

(Idempotence)

(Associativity)

(Absorption)

Truth Tables vs Proofs of Equivalence

• When we want to show that $\phi(x_1, x_2, \dots x_n) = \psi(x_1, x_2, \dots, x_n)$:

Truth Table		Equivalence Proof	
Pro	Con	Pro	Con
Always works	Requires 2 ⁿ <u>space</u>	Often occupies much less than 2^n space	There are some cases where it will still take 2^n space / time
No "cleverness" needed: just build all rows mechanically	Requires 2^n <u>time</u>	Often spends much less than 2^n time	Requires "cleverness"

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