

**HW 10 CMSC 452. Morally Due April 17**  
**SOLUTIONS**

1. (5 points) What is your name? Write it clearly. Staple the HW.
2. (30 points) Show that the following problem is in NP:  
All  $(A, b)$  where
  - (a)  $A$  is an  $n \times n$  matrix of integers.
  - (b)  $b$  is a vector of  $n$  integers
  - (c) There exists a vector  $x$  of integers between -5 and 5 such that  $Ax = b$
  - (d) (Think about- not to be handed in) If we allow  $x$  to be a rational then is the problem still in NP? If we allow  $x$  to be any integer (so remove the bounds -5 to 5) then is the problem in NP?

**SOLUTION TO PROBLEM TWO**

The witness is just the vector  $x$ . The two things to check:

- Is it easy to check that  $Ax = b$ . This takes  $O(n^2)$  operations, so yes that's easy.
- Is  $x$  short? Yes,  $x$  is  $n$  vectors of 3 bits each, so  $O(n)$  length.

If we let  $x$  be a rational OR if we let  $x$  be any integer than we would need to PROVE that if there is an  $x$  there is a short one. This is true but not obvious.

3. (30 points) If  $G = (V, E)$  is a graph then  $X \subseteq V$  is a *Vertex Cover* if for every  $e \in E$ , there is a  $v \in X$  that is the endpoint of  $e$ .

Show that the following problem is in P:

$$\{G : G \text{ has a vertex cover of size } 17 \}$$

**SOLUTION TO PROBLEM THREE**

- (a) Input  $G$

- (b) Look at all 17-sized sets of vertices  $X$ . There are  $\leq n^{17}$  of these. For each one CHECK if it is a vertex cover. This takes  $O(e)$  time- for every edge check if one of the endpoints is in  $X$ . If ever you get a YES then output YES and stop
- (c) If you got this far then output NO

The algorithm takes  $O(en^{17}) \leq O(n^{19})$ .

4. (40 points) Describe an NFA with  $\leq 500$  states (it will actually be far less than this) for the set

$$\{a^y : (y \neq 999) \wedge (y \neq 1000)\}$$

Prove that it works by showing that if  $a^y$  is rejected then  $y = 999$  or  $y = 1000$ .

HINT- For the big loop use 32 and 33. You may USE the fact that

- For all  $n \geq 992$  there exists  $x, y \in N$  such that  $n = 32x + 33y$ .
- There does not exist  $x, y \in N$  such that  $991 = 32x + 33y$ .

#### SOLUTION TO PROBLEM FOUR

READ the writeup of the small NFA that only rejects  $a^{1000}$  before reading this.

Since 991 is not a linear combo of 32 and 33,

$991 - 33 = 958$ , and  $991 - 32 = 959$  are also not linear combos of 32 and 33.

Hence the big-loop NFA with a loop of size 33 and a shortcut for 32 does not accept  $a^{958}$  or  $a^{959}$

Take that NFA and add 41 states at the beginning. NOW this NFA:

- For all  $y \geq 1001$  accepts  $a^y$ .
- Does not accept  $a^{999}$ .
- Does not accept  $a^{1000}$ .

There are still many strings we want to accept but do not yet. Use e-transitions to add loops that recognize the following sets:

- (a)  $\{a^y : y \not\equiv 999 \pmod{3} \wedge y \not\equiv 1000 \pmod{3}\}$
- (b)  $\{a^y : y \not\equiv 999 \pmod{5} \wedge y \not\equiv 1000 \pmod{5}\}$
- (c)  $\{a^y : y \not\equiv 999 \pmod{7} \wedge y \not\equiv 1000 \pmod{7}\}$
- (d)  $\{a^y : y \not\equiv 999 \pmod{11} \wedge y \not\equiv 1000 \pmod{11}\}$
- (e)  $\{a^y : y \not\equiv 999 \pmod{13} \wedge y \not\equiv 1000 \pmod{13}\}$
- (f)  $\{a^y : y \not\equiv 999 \pmod{17} \wedge y \not\equiv 1000 \pmod{17}\}$
- (g)  $\{a^y : y \not\equiv 999 \pmod{19} \wedge y \not\equiv 1000 \pmod{19}\}$

Let  $M$  be the NFA that had from the start state  $\epsilon$ -transitions to the 33-loop-32-shortcut and to all the prime machines above. The number of states is

$$1 + 41 + 33 + 3 + 5 + 7 + 11 + 13 + 17 + 19 = 150$$

We show that the only strings REJECTED are  $a^{999}$  and  $a^{1000}$ .

Let  $a^y$  be a rejected string. Because of the big loops  $y \leq 1000$ .

Because of the prime machines:

- (a)  $y \equiv 999 \pmod{3} \vee y \equiv 1000 \pmod{3}$
- (b)  $y \equiv 999 \pmod{5} \vee y \equiv 1000 \pmod{5}$
- (c)  $y \equiv 999 \pmod{7} \vee y \equiv 1000 \pmod{7}$
- (d)  $y \equiv 999 \pmod{11} \vee y \equiv 1000 \pmod{11}$
- (e)  $y \equiv 999 \pmod{13} \vee y \equiv 1000 \pmod{13}$
- (f)  $y \equiv 999 \pmod{17} \vee y \equiv 1000 \pmod{17}$
- (g)  $y \equiv 999 \pmod{19} \vee y \equiv 1000 \pmod{19}$

Take the set  $X = \{3, 5, 7, 11, 13, 17, 19\}$ .

For each  $p \in X$  either  $y \equiv 999 \pmod{p}$  OR  $y \equiv 1000 \pmod{p}$ .

Map  $p$  to either 999 or 1000 depending on which is the case.

Since  $X$  has 7 elements, either at least four map to 999 or at least 4 map to 1000.

Let those four be  $p_1, p_2, p_3, p_4$ .

If those four all map to 999 then

$$y \equiv 999 \pmod{p}_1$$

$$y \equiv 999 \pmod{p}_2$$

$$y \equiv 999 \pmod{p}_3$$

$$y \equiv 999 \pmod{p}_4$$

Hence

$$y \equiv 999 \pmod{p_1 p_2 p_3 p_4}.$$

since  $p_1 p_2 p_3 p_4 \geq 3 \times 5 \times 7 \times 11 = 1155 > 999$  it must be that  $y = 999$ .

If those four all map to 1000 then

$$y \equiv 1000 \pmod{p}_1$$

$$y \equiv 1000 \pmod{p}_2$$

$$y \equiv 1000 \pmod{p}_3$$

$$y \equiv 1000 \pmod{p}_4$$

Hence

$$y \equiv 1000 \pmod{p_1 p_2 p_3 p_4}.$$

since  $p_1 p_2 p_3 p_4 \geq 3 \times 5 \times 7 \times 11 = 1155 > 1000$  it must be that  $y = 1000$ .