# Concrete Time Hierarchy Theorem 

Exposition by William Gasarch-U of MD

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## Concrete Time Hierarchy Theorem

Definition Let $A \subseteq\{0,1\}^{*} . A \in \operatorname{DTIME}\left(n^{3}\right)$ is there is a Java Program $J$ such that the following hold.

1. If $x \in A$ then $J(x)$ outputs YES.
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1. There is a Java Program $J$ that, on input $x$, will output YES if $x \in A$, and will output NO if $x \notin A$.
2. $A \notin \operatorname{DTIME}\left(n^{3}\right)$.

## ASCII Table

| Hex | Dec | Char |  | \|Hex | Dec | Char | Hex | Dec | Char | Hex | Dec | Char |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $0 \times 00$ | 0 | NULL | null | $0 \times 20$ | 32 | Space | 0x40 | 64 | e | 0x60 | 96 |  |
| $0 \times 01$ | 1 | SOH | Start of heading | $0 \times 21$ | 33 | 1 | 0x41 | 65 | A | 0x61 | 97 | a |
| $0 \times 02$ | 2 | STX | Start of text | $0 \times 22$ | 34 | " | 0x42 | 66 | B | 0x62 | 98 | b |
| $0 \times 03$ | 3 | ETX | End of text | $0 \times 23$ | 35 | \# | $0 \times 43$ | 67 | C | 0x63 | 99 | c |
| $0 \times 04$ | 4 | EOT | End of transmission | $0 \times 24$ | 36 | \$ | 0x44 | 68 | D | 0x64 | 100 | d |
| $0 \times 05$ | 5 | ENQ | Enquiry | 0×25 | 37 | \% | 0x45 | 69 | E | 0x65 | 101 | e |
| $0 \times 06$ | 6 | ACK | Acknowledge | $0 \times 26$ | 38 | \& | 0x46 | 70 | F | 0x66 | 102 | f |
| $0 \times 07$ | 7 | BELL | Bell | $0 \times 27$ | 39 | , | 0x47 | 71 | G | 0x67 | 103 | $g$ |
| $0 \times 08$ | 8 | BS | Backspace | $0 \times 28$ | 40 | ( | $0 \times 48$ | 72 | H | 0x68 | 104 | h |
| $0 \times 09$ | 9 | TAB | Horizontal tab | 0x29 | 41 | ) | 0x49 | 73 | I | 0x69 | 105 | i |
| $0 \times 0 \mathrm{~A}$ | 10 | LF | New line | 0x2A | 42 | * | 0x4A | 74 | J | 0x6A | 106 | $j$ |
| $0 \times 0 \mathrm{~B}$ | 11 | VT | Vertical tab | $0 \times 2 \mathrm{~B}$ | 43 | + | 0x4B | 75 | K | 0x6B | 107 | k |
| $0 \times 0 \mathrm{C}$ | 12 | FF | Form Feed | 0x2C | 44 | , | 0x4C | 76 | L | 0x6C | 108 | 1 |
| $0 \times 0 \mathrm{D}$ | 13 | CR | Carriage return | 0x2D | 45 | - | 0x4D | 77 | M | 0x6D | 109 | m |
| $0 \times 0 \mathrm{E}$ | 14 | So | Shift out | $0 \times 2 \mathrm{E}$ | 46 | - | 0x4E | 78 | N | 0x6E | 110 | n |
| $0 \times 0 \mathrm{~F}$ | 15 | SI | Shift in | $0 \times 2 \mathrm{~F}$ | 47 | 1 | 0x4F | 79 | 0 | 0x6F | 111 | - |
| $0 \times 10$ | 16 | DLE | Data link escape | 0×30 | 48 | 0 | 0x50 | 80 | P | 0x70 | 112 | P |
| $0 \times 11$ | 17 | DC1 | Device control 1 | 0×31 | 49 | 1 | 0x51 | 81 | Q | 0x71 | 113 | q |
| $0 \times 12$ | 18 | DC2 | Device control 2 | $0 \times 32$ | 50 | 2 | 0x52 | 82 | R | 0x72 | 114 | r |
| $0 \times 13$ | 19 | DC3 | Device control 3 | $0 \times 33$ | 51 | 3 | 0x53 | 83 | S | 0x73 | 115 | $s$ |
| $0 \times 14$ | 20 | DC4 | Device control 4 | $0 \times 34$ | 52 | 4 | 0×54 | 84 | T | 0x74 | 116 | t |
| $0 \times 15$ | 21 | NAK | Negative ack | $0 \times 35$ | 53 | 5 | 0x55 | 85 | U | 0x75 | 117 | u |
| $0 \times 16$ | 22 | SYN | Synchronous idle | $0 \times 36$ | 54 | 6 | 0x56 | 86 | V | 0x76 | 118 | v |
| $0 \times 17$ | 23 | ETB | End transmission block | $0 \times 37$ | 55 | 7 | 0x57 | 87 | W | 0x77 | 119 | w |
| $0 \times 18$ | 24 | CAN | Cancel | $0 \times 38$ | 56 | 8 | 0x58 | 88 | X | 0x78 | 120 | x |
| $0 \times 19$ | 25 | EM | End of medium | 0×39 | 57 | 9 | 0x59 | 89 | Y | 0x79 | 121 | y |
| 0x1A | 26 | SUB | Substitute | 0×3A | 58 | : | 0x5A | 90 | z | 0x7A | 122 | z |
| $0 \times 1 \mathrm{~B}$ | 27 | FSC | Escape | 0x3B | 59 | ; | 0x5B | 91 | [ | 0x7B | 123 | $\{$ |
| $0 \times 1 \mathrm{C}$ | 28 | FS | File separator | 0x3C | 60 | $<$ | 0x5C | 92 | 1 | 0x7C | 124 |  |
| $0 \times 1 \mathrm{D}$ | 29 | GS | Group separator | 0x3D | 61 | = | 0x5D | 93 | ] | 0x7D | 125 | \} |
| $0 \times 1 \mathrm{E}$ | 30 | RS | Record separator | 0×3E | 62 | > | 0x5E | 94 | ^ | 0x7E | 126 | - |
| $0 \times 1 \mathrm{~F}$ | 31 | US | Unit separator | 0×3F | 63 | ? | $0 \times 5 \mathrm{~F}$ | 95 | - | 0x7F | 127 | DEL |

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The ASCII table maps symbols into decimal numbers between 0 and 127 . We include leading 0 's.

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- $0, \ldots, 9$ code to $048, \ldots, 057$.


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- $:, ;,<,=,>, ?$, code to 058 to 064.


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- A,...,Z code to 065,...,090.
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- I won't bother with the rest. See table.


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$x=x+12$
$x$ maps to 120 .
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1 maps to 049
2 maps to 050
So this piece of code maps to $120,061,120,043,049,050$


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4. If $J$ IS NOT a valid Java Program then map $i$ to $\downarrow$.
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Let $J_{i}$ be the Java program that $i$ maps to. So

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$J_{1}, J_{2}, \ldots, \ldots$ is the list of all Java Programs.

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$J_{1}^{\prime}, J_{2}^{\prime}, \ldots, \ldots$ is the list of all $n^{3}$-time Java Programs .
Upshot If $A \in \operatorname{DTIME}\left(n^{3}\right)$ then there exists $i$ such that $J_{i}^{\prime}$ recognizes $A$.

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2. $A \notin \operatorname{DTIME}\left(n^{3}\right)$.

Proof Let $A$ be decided by the following program

1. Input $(x)$. If $x \notin 0^{*}$ output NO and stop. Otherwise $x=0^{n}$.
2. Run $J_{n}^{\prime}\left(0^{n}\right)$.
3. If result is YES then output NO and stop.
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1. There is a Java Program $J$ that, on input $x$, will output YES if $x \in A$, and will output NO if $x \notin A$.
2. $A \notin \operatorname{DTIME}\left(n^{3}\right)$.

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4. If result is NO then output YES and stop.
1) This is clearly a program that recognizes $A$.
2) Proof that $A \notin \operatorname{DTIME}\left(n^{3}\right)$ on next slide.

## $A \notin \operatorname{DTIME}\left(n^{3}\right)$

1. Input $(x)$. If $x \notin 0^{*}$ output NO and stop. Otherwise $x=0^{n}$.
2. Run $J_{n}^{\prime}\left(0^{n}\right)$.
3. If result is YES then output NO and stop.
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Let $A\left(0^{n}\right)$ be YES if $0^{n} \in A$ and NO if $0^{n} \notin A$.

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Let $A\left(0^{n}\right)$ be YES if $0^{n} \in A$ and NO if $0^{n} \notin A$.
$J_{1}^{\prime}$ cannot recognize $A: J_{1}^{\prime}\left(0^{1}\right)$ and $A\left(0^{1}\right)$ DIFFER.

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$J_{2}^{\prime}$ cannot recognize $A: J_{2}^{\prime}\left(0^{2}\right)$ and $A\left(0^{2}\right)$ DIFFER.
$J_{n}^{\prime}$ cannot recognize $A: J_{n}^{\prime}\left(0^{n}\right)$ and $A\left(0^{n}\right)$ DIFFER.

## $A \notin \operatorname{DTIME}\left(n^{3}\right)$

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2. Run $J_{n}^{\prime}\left(0^{n}\right)$.
3. If result is YES then output NO and stop.
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Let $A\left(0^{n}\right)$ be YES if $0^{n} \in A$ and NO if $0^{n} \notin A$.
$J_{1}^{\prime}$ cannot recognize $A$ : $J_{1}^{\prime}\left(0^{1}\right)$ and $A\left(0^{1}\right)$ DIFFER.
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$J_{n}^{\prime}$ cannot recognize $A: J_{n}^{\prime}\left(0^{n}\right)$ and $A\left(0^{n}\right)$ DIFFER.
So NO $J_{n}^{\prime}$ recognizes $A$. Hence $A \notin \operatorname{DTIME}\left(n^{3}\right)$.

