

# BILL, RECORD LECTURE!!!!

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# Gen 2-letter Sub and Matrix Codes

September 28, 2020

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**Def** Gen Sub 2-Cipher with perm  $f$  on  $\{0, \dots, 25\}^2$ .

1. Encrypt via  $xy \rightarrow f(xy)$ .
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Need bijection of  $\{0, \dots, 25\} \times \{0, \dots, 25\}$  that is easy to use.

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Do you recognize the expression  $ad - bc$ ? Determinant!

# Inverse Matrix in $\mathbb{C}$ and in Mods

$$M = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$$

1. Matrix  $M$  over  $\mathbb{C}$  has an inverse iff  $ad - bc \neq 0$ .
2. Matrix  $M$  over Mod  $n$  has an inverse iff  $ad - bc$  is rel prime to  $n$  iff  $ad - bc$  has an inverse in Mod  $n$ .
3. Matrix  $M$  over Mod 26 has an inverse iff  $ad - bc$  is rel prime to 26 iff  $ad - bc$  has no factors of 2 or 13 iff has an inverse in Mod 26.

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So what to do?

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3. Eve cannot use brute force. Key Space is  $\sim 26^{64} \sim 10^{90}$ ,  
Number of protons is  $\sim 10^{79}$ . (the number of non-invertible matrices is very small so  $26^{64}$  is a good approximation).

# Lets Try Brute Force Even if Slow

1. Input  $T$ , a coded text.
2. For EVERY  $8 \times 8$  invertible matrix  $M$  over  $\mathbb{Z}_{26}$ ,
  - 2.1 Decode  $T$  into  $T'$  using  $M$ .
  - 2.2 IF LOOKS-LIKE-ENGLISH( $T'$ )=YES then STOP and output  $T'$ , else goto next matrix  $M$ .

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YES- we can do  $8 \times 26^8$ .

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Let  $Mt_i = m_i$ . Then  $(1, 1, 7) \cdot t_i = m_i^1$  is first letter of  $m_i$ .

$$(m_1^1, m_2^1, m_3^1, \dots, m_N^1)$$

is every third letter. Can do IS-ENGLISH on it.

## Can Crack in $8 \times 26^8$

Eve knows that Alice and Bob decode with  $8 \times 8$  Matrix  $M$ .

Ciphertext is

$$T = t_1 t_2 \cdots t_N \quad t_i = t_i^1 \cdots t_i^8$$

For  $i = 1$  to 8

For all  $r \in \mathbb{Z}_{26}^8$  (guess that  $r$  is  $i$ th row of  $B$ ).

$T' = (r \cdot t_1, \dots, r \cdot t_N)$  (Is every 8th letter.)

IF IS-ENGLISH( $T'$ )=YES then  $r_i = r$  and goto next  $i$ . Else  
goto the next  $r$ .

$M$  is

$$\begin{pmatrix} \cdots & \cdots & \cdots \\ \vdots & \vdots & \vdots \\ r_1 & \cdots & r_n \\ \vdots & \vdots & \vdots \\ \cdots & \cdots & \cdots \end{pmatrix}$$

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The row-by-row method takes  $O(n26^n)$ .

# Important Lesson

**Assume:**  $26^{64}$  time is big enough to thwart Eve.

1. If we think that best Eve can do is  $O(26^{n^2})$  then we take  $n = 8$ , so Eve needs  $O(26^{64})$ .
2. If we think that best Eve can do is  $O(n26^n)$  then we take  $n = 80$ , so Eve needs  $O(80 \times 26^{80})$ .

The  $O(n \times 26^n)$  cracking **does not** show that Matrix Cipher is insecure, but it still is very important: Alice and Bob must increase their parameters. That is already a win since it makes life harder for Alice and Bob.

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Above attack on Matrix Cipher is a microcosm of this history.

Proofs rely on limiting what Eve can do, and hence do not work if Eve does something else.

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3. So this looks like a strong cipher. Is it crackable?

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3. Eve will have old messages and what they decoded to.

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5. Eve knows that  $(3, 9) = M(13, 24)$ .

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4. We will do this in the next slide packet.