BILL RECORD THIS LECTURE

September 16, 2020

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Affine and Quadratic Ciphers

September 16, 2020

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The Affine Ciphers

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Recall: Shift cipher with shift *s*:

- 1. Encrypt via $x \rightarrow x + s \pmod{26}$.
- 2. Decrypt via $x \rightarrow x s \pmod{26}$.

We replace x + s with more elaborate functions.

Def The Affine cipher with *a*, *b*:

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Condition on *a*, *b* so that $x \rightarrow ax + b$ is a bij: *a* rel prime to 26. Condition on *a*, *b* so that *a* has an inv mod 26: *a* rel prime to 26. This is achieved by making *a* relatively prime to 26.

Shift vs Affine

Shift: Key space is size 26.

Affine: Key space is $\{1,3,5,7,9,11,15,17,19,21,23,25\} \times \{0,\ldots,25\}$ which has $12 \times 26 = 312$ elements.

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Both Need: The **Is-English** algorithm. Reading through 312 transcripts to see which one **looks like English** would take A LOT of time!

Key Length of Shift and Affine Ciphers

Let's look at the keys for Shift and Affine.

- 1. Shift cipher key in $\{0, \ldots, 25\}$. 5 bits.
- 2. Affine cipher Key in

 $\{1,3,5,7,9,11,15,17,19,21,23,25\}\times\{0,\ldots,25\}.$ 312 keys, need 9 bits.

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- 2. Bob must compute the inverse of $a \mod m$ in order to decode.
- 3. If Alice wants to also get messages and decode them, she also has to compute the inverse of *a* mod *m* in order to decode.

If $\Sigma = \{a, \ldots, z\}$ (size 26) then, as we saw, the set is

 $\{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$ 12 possibilities

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If given *m*, want to know how many elements in $\{1, \ldots, m-1\}$ are relatively prime to *m*.

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Finding Inverse Mod n

September 16, 2020

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1. Affine cipher over alphabet of size *n*, need to know if *a* has an inverse, and if so, what it is.

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- 3. (Later) Implementing RSA.

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- 4. (Later) Cracking RSA.

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- 5. (Later) Factoring Algorithms.

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- 2. (Later) Cracking psuedo-random ciphers.
- 3. (Later) Implementing RSA.
- 4. (Later) Cracking RSA.
- 5. (Later) Factoring Algorithms.
- 6. Many Many Others!

Greatest Common Divisor (GCD)

We first need to look at GCD. GCD(m, n) is the largest number that divides m AND n. **Examples** GCD(10, 15) =

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Idea: Keep subtracting smaller from larger: GCD(404, 192) =

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Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192) = GCD(212 - 192, 192) = GCD(20, 192).Could keep going, but will be subtracting 20's for a while.

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```

d is largest divisor of **both** 404 and 192 IFF *d* is largest divisor of 404 and 404 - 192 = 212.

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192) = GCD(212 - 192, 192) = GCD(20, 192).Could keep going, but will be subtracting 20's for a while.

Idea: Subtract LOTS of 20's.

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Idea: Subtract LOTS of 20's. Largest $x:192 - 20x \ge 0$, x = 9.

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Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192) = GCD(212 - 192, 192) = GCD(20, 192).Could keep going, but will be subtracting 20's for a while.

Idea: Subtract LOTS of 20's. Largest $x:192 - 20x \ge 0$, x = 9. = GCD(20, 192 - 20 × 9 = 12) = GCD(20 - 12, 12) = GCD(8, 12) = GCD(8, 12 - 8 = 4) = GCD(8 - 2 × 4, 4) = GCD(0, 4) = 4.

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 $404 = 2 \times 192 + 20$

 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \end{array}$



 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \end{array}$



 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \end{array}$



 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$

 $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD.

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 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. Can use this to write 4 as a combination of 404 and 192

 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \\ 8 = 4 \times 2 + 0 \text{ STOP HERE and go back one: 4 is the GCD.} \\ \hline \textbf{Can use this to write 4 as a combination of 404 and 192} \\ \text{Write 4 as a combo of 12's and 8's:} \\ 4 = 12 - 1 \times 8 \end{array}$

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GCD(404,192) The Short Way and More Info

 $404 = 2 \times 192 + 20$

 $\begin{array}{l} 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \\ 8 = 4 \times 2 + 0 \text{ STOP HERE and go back one: 4 is the GCD.} \\ \hline \textbf{Can use this to write 4 as a combination of 404 and 192} \\ Write 4 as a combo of 12's and 8's: \\ 4 = 12 - 1 \times 8 \\ Write 8 as a combo of 20's and 12's: \\ 4 = 12 - 1 \times (20 - 12) = 2 \times 12 - 1 \times 20 \end{array}$

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GCD(404,192) The Short Way and More Info

 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. Can use this to write 4 as a combination of 404 and 192 Write 4 as a combo of 12's and 8's: $4 = 12 - 1 \times 8$ Write 8 as a combo of 20's and 12's: $4 = 12 - 1 \times (20 - 12) = 2 \times 12 - 1 \times 20$ Write 12 as combo of 192's and 20's: $4 = 2 \times (192 - 9 \times 20) - 1 \times 20 = 2 \times 192 - 19 \times 20$

GCD(404,192) The Short Way and More Info

 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. Can use this to write 4 as a combination of 404 and 192 Write 4 as a combo of 12's and 8's: $4 = 12 - 1 \times 8$ Write 8 as a combo of 20's and 12's: $4 = 12 - 1 \times (20 - 12) = 2 \times 12 - 1 \times 20$ Write 12 as combo of 192's and 20's: $4 = 2 \times (192 - 9 \times 20) - 1 \times 20 = 2 \times 192 - 19 \times 20$ Write 20 as a combo of 404 and 192: $4 = 2 \times 192 - 19 \times (404 - 2 \times 192) = 39 \times 192 - 19 \times 404$ **Upshot:** GCD(m, n) is a combo of m and n

 $101 = 2 \times 38 + 25$

 $101 = 2 \times 38 + 25$ $38 = 1 \times 25 + 13$



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101 = 2 \times 38 + 25 
38 = 1 \times 25 + 13 
25 = 1 \times 13 + 12
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101 = 2 \times 38 + 25

38 = 1 \times 25 + 13

25 = 1 \times 13 + 12

13 = 12 + 1
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\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \ \mbox{Go back one: } 1 \ \mbox{is the GCD.} \end{array}
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 $1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25$

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$$\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \end{array}$$
 Go back one: 1 is the GCD.

$$\begin{array}{l} 1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25 \\ 1 = 2(38 - 25) - 25 = 2 \times 38 - 3 \times 25 \end{array}$$

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$$1 = 2(38 - 25) - 25 = 2 \times 38 - 3 \times 25$$

$$1 = 2 \times 38 - 3 \times (101 - 2 \times 38) = 8 \times 38 - 3 \times 101$$

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$$\begin{split} &101 = 2 \times 38 + 25 \\ &38 = 1 \times 25 + 13 \\ &25 = 1 \times 13 + 12 \\ &13 = 12 + 1 \\ &12 = 12 \times 1 + 0. \end{split}$$
 Go back one: 1 is the GCD.

$$1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25$$

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$$1 = 2 \times 38 - 3 \times (101 - 2 \times 38) = 8 \times 38 - 3 \times 101$$

$$1 = 8 \times 38 - 3 \times 101$$

Why is this interesting? Hint: What was our original goal?

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Take both sides mod 101
 $1 \equiv 8 \times 38 \pmod{101}$

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8 is the inverse of 38 mod 101

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Given m, n with m < n we want to know

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Is there an inverse of m mod n

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- Is there an inverse of m mod n
- If so then find it

Given m, n with m < n we want to know

- Is there an inverse of m mod n
- If so then find it
- 1. Find GCD(m, n). If it is NOT 1 then NO inverse.

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Given m, n with m < n we want to know

- Is there an inverse of m mod n
- If so then find it
- 1. Find GCD(m, n). If it is NOT 1 then NO inverse.
- 2. If it IS 1 then use the work you did to find $\operatorname{GCD}(m, n)$ to find $a, b \in \mathbb{Z}$

$$am + bn = 1$$

$$am\equiv 1\pmod{n}$$

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3. a is the inverse of $m \mod n$.

Given m, n with m < n we want to know

- Is there an inverse of m mod n
- If so then find it
- 1. Find GCD(m, n). If it is NOT 1 then NO inverse.
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a is the inverse of m mod n. Not quite: (1) a might be negative (2) a might be > n. That won't do!

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 a is the inverse of m mod n. Not quite: (1) a might be negative (2) a might be > n. That won't do! Take a (mod n).

September 16, 2020

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Def The Quadratic cipher with a, b, c: Encrypt via $x \rightarrow ax^2 + bx + c \pmod{26}$.

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Def The Quadratic cipher with a, b, c: Encrypt via $x \rightarrow ax^2 + bx + c \pmod{26}$.

Does this work? Vote YES or NO.

Def The Quadratic cipher with a, b, c: Encrypt via $x \rightarrow ax^2 + bx + c \pmod{26}$.

Does this work? Vote YES or NO. Answer: NO

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Does this work? Vote YES or NO. Answer: NO

No easy test for Invertibility (depends on def of easy). How Easy?: Given a quadratic f(x) one could compute $f(0), \ldots, f(25)$ all mod 26 and see if all are different.

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Does this work? Vote YES or NO. Answer: NO

No easy test for Invertibility (depends on def of easy). How Easy?: Given a quadratic f(x) one could compute $f(0), \ldots, f(25)$ all mod 26 and see if all are different.

- 1. This takes to long.
- The security is not good enough to justify taking this long setting it up.

History of the Quadratic Cipher

The first place The Quadratic Cipher appeared was

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History of the Quadratic Cipher

The first place The Quadratic Cipher appeared was

my 3-week course on crypto for High School Students in 2010.

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History of the Quadratic Cipher

The first place The Quadratic Cipher appeared was

my 3-week course on crypto for High School Students in 2010.

So, as the kids say, it's not a thing.

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When looking at a cipher one usually asks:

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Is the cipher secure?



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- That is a good question.

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Is the cipher easy to use?

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When looking at a cipher one usually asks:

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But there is another important one:

Is the cipher easy to use?

Quadratic Cipher fails the ease of use test.
The Point of Presenting the Quadratic Cipher

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When looking at a cipher one usually asks:

- Is the cipher secure?
- That is a good question.

But there is another important one:

Is the cipher easy to use?

Quadratic Cipher fails the ease of use test.

It is also insecure.

Shift and Affine:



Shift and Affine:

Some math is used to encrypt and decrypt.

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Shift and Affine:

- Some math is used to encrypt and decrypt.
- The math makes it easy to use. Short Key!

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Shift and Affine:

- Some math is used to encrypt and decrypt.
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▶ The math makes it insecure. Few Keys!

Shift and Affine:

- Some math is used to encrypt and decrypt.
- The math makes it easy to use. Short Key!
- ▶ The math makes it insecure. Few Keys!

Next slide packet: We present a cipher with **less** math so **more secure** in next slide packet.

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BILL STOP RECORDING THIS LECTURE

September 16, 2020

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