

BILL RECORD THIS LECTURE

September 20, 2020

Gen Sub Cipher and Random-Looking Ciphers

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General Substitution Cipher

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We present the **General Substitution Cipher** which:

- ▶ Has a large keyspace.
- ▶ Does not use any math.

General Substitution Cipher

Def Gen Sub Cipher with perm f on $\{0, \dots, 25\}$.

1. Encrypt via $x \rightarrow f(x)$.
2. Decrypt via $x \rightarrow f^{-1}(x)$.

General Substitution Cipher: Example

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<i>d</i>	<i>i</i>	<i>a</i>	<i>b</i>	<i>e</i>	<i>g</i>	<i>f</i>	<i>c</i>	<i>h</i>

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Decrypt Using:

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<i>c</i>	<i>d</i>	<i>h</i>	<i>a</i>	<i>e</i>	<i>g</i>	<i>f</i>	<i>i</i>	<i>b</i>

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<i>c</i>	<i>d</i>	<i>h</i>	<i>a</i>	<i>e</i>	<i>g</i>	<i>f</i>	<i>i</i>	<i>b</i>

If the message is **FBI** it will encrypt to **GIH**.

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Yes: Eve can use Freq Analysis

Freq Analysis

Alice sends Bob a LONG text encrypted by Gen Sub Cipher.
Eve finds freq of letters, pairs, triples,

Text in English.

1. Can use known freq: *e* is most common letter, *th* is most common pair.
2. If Alice is telling Bob about Mid East Politics than may need to adjust: *q* is more common (Iraq, Qatar) and some words more common.

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4. Watch Jeopardy! Alex Trebek's fun TV quiz game.

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We assume long normal texts!

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4. Spoiler Alert: David Zhen has a program that cracks the gen sub cipher.

Random-Looking Ciphers

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Alternatives to Gen Sub (History)

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1. We show two such methods.
2. These methods are primitive examples of **psuedo-random generators** which take a short string and make a **random-looking** much longer string. These are important in crypto. We will encounter them again.

Keyword-Shift Cipher. Key is (Word,Shift)

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1. List out the key word and then the remaining letters:

| j | a | c | k | b | d | e | f | g | h | i |

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2. Now do Shift 4 on this:

| f | g | h | i | j | a | c | k | b | d | e |

This is where a, b, c, \dots go, so:

| a | b | c | d | e | f | g | h | i | j | k |
| f | g | h | i | j | a | c | k | b | d | e |

Keyword-Shift Cipher. Key is (Word,Shift) (cont)

To encrypt use:

<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>	<i>g</i>	<i>h</i>	<i>i</i>	<i>j</i>	<i>k</i>
<i>f</i>	<i>g</i>	<i>h</i>	<i>i</i>	<i>j</i>	<i>a</i>	<i>c</i>	<i>k</i>	<i>b</i>	<i>d</i>	<i>e</i>

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<i>f</i>	<i>g</i>	<i>h</i>	<i>i</i>	<i>j</i>	<i>a</i>	<i>c</i>	<i>k</i>	<i>b</i>	<i>d</i>	<i>e</i>

To decrypt you invert the table:

<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>	<i>g</i>	<i>h</i>	<i>i</i>	<i>j</i>	<i>k</i>
<i>f</i>	<i>i</i>	<i>g</i>	<i>j</i>	<i>k</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>h</i>

From Short Key Got Rand-Looking Perm(?)

From (jack,4) (which is short) we got

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Does this cipher look like it was generated randomly? Discuss.

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5. Keyword-Shift Cipher, 4-let keywords, not that rand looking.

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I suspect this would not leave a tell-tale sign of not being random.

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1. Write **jack**. Under it write the rest of Σ in blocks of size **|jack|**.

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2. Write down these letters column by column:

| *j* | *b* | *g* | *a* | *d* | *h* | *c* | *e* | *i* | *k* | *f* |

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2. Write down these letters column by column:

| *j* | *b* | *g* | *a* | *d* | *h* | *c* | *e* | *i* | *k* | *f* |

3. Put the letters in order under it:

<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>	<i>g</i>	<i>h</i>	<i>i</i>	<i>j</i>	<i>k</i>
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Today's lecture will **support** her viewpoint.

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If T then Alice gen rand perm of Σ and sends to Eve.

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How well can Eve do? Discuss.

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2. Eve is limited computationally. We will clarify later.

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Alice gives Eve perm τ . If τ is one of the σ_i then Eve says KS, otherwise Eve says RP.

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Prob Eve right is $1 - \frac{1}{10^{17}} = 0.9999999999999999 = L$.

Prob Eve wins is

$$\Pr(KS) \times 1 + \Pr(RP) \times L = \frac{1}{2} \times 1 + \frac{1}{2} \times L = \frac{1}{2}(1 + L) = L'$$

which is very close to 1.

Upshot Unlimited Eve wins most of the time.

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Note this is not really rigorous since we are thinking of $|\Sigma|$ as 26, but the idea is sound.

Strategy for Comp Limited Eve (Motivation)

We will do a few more Keyword-Shift Ciphers and see if there are hints that they **are** Keyword-Shift.

Recall Keyword-Shift Cipher with (jack,4)

$\Sigma = \{a, \dots, k\}$. **Key:** (jack,4).

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--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--

2. Now do Shift 4 on this:

	<i>f</i>		<i>g</i>		<i>h</i>		<i>i</i>		<i>j</i>		<i>a</i>		<i>c</i>		<i>k</i>		<i>b</i>		<i>d</i>		<i>e</i>	
--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--	----------	--

3. This is where a, b, c, \dots , goes to. Put the table in order to get how to encode.

	<i>a</i>		<i>b</i>		<i>c</i>		<i>d</i>		<i>e</i>		<i>f</i>		<i>g</i>		<i>h</i>		<i>i</i>		<i>j</i>		<i>k</i>	
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Much of the bottom row is in alpha-order.

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| g | c | j | b | d | e | f | h | i | k | a |

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---	---	---	---	---	---	---	---	---	---	---

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a	b	c	d	e	f	g	h	i	j	k
g	c	j	b	d	e	f	h	i	k	a

Again we get three letters in a row!

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2. If τ has 3 consecutive letters (e.g., p, q, r) then say KS, else say RP. (We do not count wrap around.)

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Prob that Eve wins is $\geq 1 - 0.04 = 0.96$.

Prob Eve wins is $\frac{1}{2} \times 1 + \frac{1}{2} \times 0.096 = 0.98$

Definition of **Random-Looking**

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C generates perms that **look random** if when Eve plays the game the prob that she wins is $\leq \frac{1}{2}$.

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This is an example of a **pseudo-random generator**.
We will visit that concept later and use a similar game.

Which is Better Keyword-Mixed or Keyword-Shift?

1. Have given you a way to find out for yourself.
2. Might make into a HW.