Some Solutions to HW01 Problems

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 $x^2 + 17x + 16 \equiv 0 \pmod{100}$

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Two ways to solve. 1) Write a program that goes through all $x \in \{0, ..., 99\}$. 2) By hand and cleverness on next slide.

$$x^2 + 17x + 16 = (x + 16)(x + 1)$$

Lemma $(x + 1)(x + 16) \equiv 0 \pmod{100} \implies x + 1 \equiv 0 \pmod{5}.$

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Proof
$$x + 1 \not\equiv 0 \pmod{5} \implies x + 16 \not\equiv 0 \pmod{5} \implies$$

 $(x + 1)(x + 16) \not\equiv 0 \pmod{5} \implies (x + 1)(x + 16) \not\equiv 0$
(mod 100).

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Upshot Only need to look x such that $x + 1 \equiv 0 \pmod{5}$. **Upshot** Only need to look at $x \equiv 0 \pmod{5}$.

Lemma $(x+1)(x+16) \equiv 0 \implies x+1 \not\equiv 2 \pmod{4}$.



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Lemma $(x+1)(x+16) \equiv 0 \pmod{100} \implies x+1 \not\equiv 3 \pmod{4}$.

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Upshot Only need to look at x such that $x + 1 \equiv 0, 1 \pmod{4}$. **Upshot** Only need to look at $x \equiv 0, 3 \pmod{4}$.

Problem 2. Clever Sol Cont.

1) $x \equiv 4 \pmod{5}$ and $x \equiv 0 \pmod{4}$ implies $x \equiv 4 \pmod{20}$.

x	(x+1)(x+16)	$\equiv 0 \pmod{100}?$
4	100	Y
24	1000	Y
44	2700	Y
64	5200	Y
84	8400	Y

2) $x \equiv 4 \pmod{5}$ and $x \equiv 3 \pmod{4}$ implies $x \equiv 19 \pmod{20}$.

x	(x+1)(x+16)	$\equiv 0$	(mod 100)?
19	700		Y
39	2200		Y
59	4500		Y
79	7600		Y
99	8400		Y

SO there are 10 solutions.

Problem 2: The Point

Point of the Problem Mod 100 is very different than \mathbb{N} or \mathbb{Z} or even Mod 7 since you can have *d*th degree poly with MORE THAN *d* roots.

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Theorem If the domain is \mathbb{Z} or \mathbb{R} or \mathbb{C} (the complex numbers) then every poly of degree d has $\leq d$ roots.

The proof of this theorem used that in these domains

$$ab = 0 \implies (a = 0) \lor (b = 0)$$

How many $a, b \in \{0, \dots, 29\}$ are cool relative to 30.

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How many $a, b \in \{0, \ldots, 29\}$ are cool relative to 30.

The numbers rel prime to 30 are $\{1, 7, 11, 13, 17, 19, 23, 29\}$. Hence there are 8 of these.

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The number of b's is ALL of them: 30.

How many $a, b \in \{0, \ldots, 29\}$ are cool relative to 30.

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The number of b's is ALL of them: 30.

Hence there are $8 \times 30 = 240$ cool pairs.

Problem 4b

A student picks an $a, b \in \{0..., 29\}$ at random. What is the probability that (a, b) is cool relative to 30?

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Problem 4b

A student picks an $a, b \in \{0..., 29\}$ at random. What is the probability that (a, b) is cool relative to 30?

$$\frac{240}{30 \times 30} = \frac{8 \times 30}{30 \times 30} = \frac{8}{30} = \frac{4}{15} \sim 0.2667$$

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How many (a, b) are cool relative to 31?



How many (a, b) are cool relative to 31?

The numbers rel prime to 31 are $\{1, \ldots, 30\}$. Hence there are 30 of these.

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The number of *b*'s is ALL of them: 31.

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The number of *b*'s is ALL of them: 31.

Hence there are $30 \times 31 = 930$ cool pairs.

Problem 4d

A student picks an $a, b \in \{0..., 30\}$ at random. What is the probability that (a, b) is cool rel to 31? Give the answer to four decimal places.

$$\frac{930}{31\times31} = \frac{30\times31}{31\times31} = \frac{30}{31} = \sim 0.9677$$

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Problem 4e

What types of numbers n are such that the prob of picking an (a, b) that is cool rel to n is close to 1? Give an example of a number between 1000 and 1200 where the prob is close to 1. What is the prob? Give it to 4 places.

Problem 4e

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We want *n* to be PRIME. WE take n = 1001 which is prime. The prob of picking a cool pair is

 $\frac{1000 \times 1001}{10001 \times 1001} = \frac{1000}{1001} = 0.999.$

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Problem 4f

What types of numbers n are such that the prob of picking an (a, b) that is cool rel to n is far from 1? Give an example of a number between 1000 and 1200 where the prob is far from 1.

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Problem 4f

What types of numbers n are such that the prob of picking an (a, b) that is cool rel to n is far from 1? Give an example of a number between 1000 and 1200 where the prob is far from 1.

A number with LOTS of prime factors. We give two examples but leave it to you to work out the answer $n = 1024 = 2^{10}$. $n = 4 \times 3 \times 5 \times 17$

List all *a*, *b* so that the encode-key and the decode-key for affine are the same. All math is mod 26. Need $(\forall x)[a(ax + b) + b \equiv x]$, so $(\forall x)[a^2x + (ab + b) \equiv 1x + 0]$. We match coefficients

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The first equation yields $a \equiv 1$ or $a \equiv 25$. **Case 1** $a \equiv 1$, so $b + b \equiv 0$, $b \equiv 0$ or $b \equiv 13$. Pairs: (1,0), (1,13). **Case 2** $a \equiv 25$, so $25b + b \equiv 0$, $26b \equiv 0$. ALL *b* satisfy this. All Pairs: {(25,0), (25,1), ..., (25,25)}.

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Pairs: (1,0) (1,13), (25,0), (25,1), ..., (25,25).

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If Eve knows Alice and Bob are doing this, the key space goes from 312 to 4. So much easier for Eve to crack the code.