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Affine and Quadratic Ciphers

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The Affine Ciphers

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Recall: Shift cipher with shift $s \in \{0, \ldots, 25\}$.

- 1. Encrypt via $x \rightarrow x + s \pmod{26}$.
- 2. Decrypt via $x \rightarrow x s \pmod{26}$.

We replace x + s with more elaborate functions.

Def The Affine cipher with $a, b \in \{0, \ldots, 25\}$:

- 1. Encrypt via $x \rightarrow ax + b \pmod{26}$.
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Condition on a, b so that $x \to ax + b$ is a bij: a rel prime to 26. Condition on a, b so that a has an inv mod 26: a rel prime to 26.

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Condition on *a*, *b* so that $x \to ax + b$ is a bij: *a* rel prime to 26. Condition on *a*, *b* so that *a* has an inv mod 26: *a* rel prime to 26. This is achieved by making *a* relatively prime to 26. Note Also $a \in \{1, ..., 25\}$ and $b \in \{0, ..., 25\}$. We will not mention this again.

Shift vs Affine

Shift: Key space is size 26.

Affine: Key space is $\{1,3,5,7,9,11,15,17,19,21,23,25\}\times\{0,\ldots,25\}$ which has $12\times 26=312$ elements.

In an Earlier Era Affine would be harder to crack than Shift.

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Both Need: The **Is-English** algorithm. Reading through 312 transcripts to see which one **looks like English** would take A LOT of time!

Key Length of Shift and Affine Ciphers

Let's look at the keys for Shift and Affine.

- 1. Shift cipher key in $\{0, \ldots, 25\}$. 5 bits.
- 2. Affine cipher Key in

 $\{1,3,5,7,9,11,15,17,19,21,23,25\}\times\{0,\ldots,25\}.$ 312 keys, need 9 bits.

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If Alice and Bob use the Affine Cipher with alphabet of size m:

- 1. Alice picks a, b and must make sure that a is rel prime to m.
- 2. Bob must compute the inverse of $a \mod m$ in order to decode.

If Alice and Bob use the Affine Cipher with alphabet of size m:

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- 2. Bob must compute the inverse of *a* mod *m* in order to decode.
- 3. If Alice wants to also get messages and decode them, she also has to compute the inverse of *a* mod *m* in order to decode.

If $\Sigma = \{a, \ldots, z\}$ (size 26) then, as we saw, the set is

 $\{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$ 12 possibilities

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If $\Sigma = \{a, \dots, z\}$ (size 26) then, as we saw, the set is $\{1, 3, 5, 7, 9, 11, 15, 17, 19, 21, 23, 25\}$ 12 possibilities If $\Sigma = \{a, \dots, z, 0, \dots, 9\}$ (size 36) then, as we saw, the set is $\{1, 5, 7, 11, 13, 17, 19, 23, 25, 29, 31, 35\}$ 12 possibilities

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If given *m*, want to know how many elements in $\{1, \ldots, m-1\}$ are relatively prime to *m*.

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Finding Inverse Mod n

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Finding Inverses Given *a*, find $a^{-1} \pmod{n}$.

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Finding Inverses Given *a*, find $a^{-1} \pmod{n}$. There is a fast algorithm for this problem. This algorithm is used a lot:

1. Affine cipher over alphabet of size *n*, need to know if *a* has an inverse, and if so, what it is.

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- 3. (Later) Implementing RSA.

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- 4. (Later) Cracking RSA.

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- 5. (Later) Factoring Algorithms.

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- 2. (Later) Cracking psuedo-random ciphers.
- 3. (Later) Implementing RSA.
- 4. (Later) Cracking RSA.
- 5. (Later) Factoring Algorithms.
- 6. Many Many Others!

Greatest Common Divisor (GCD)

We first need to look at GCD. GCD(m, n) is the largest number that divides m AND n. **Examples** GCD(10, 15) =

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Examples
GCD(10, 15) = 5
GCD(11, 15) = 1
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GCD(14, 15) = 1
GCD(15, 15) =15
GCD(15, 24) = 3
GCD(15, 25) =5
GCD(15, 30) =15
GCD(15,0) = 15 (we will discuss GCD(a,0) = a later)
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d is largest divisor of ${\color{blue}{both}}$ 404 and 192 IFF

d is largest divisor of 192 and 404 - 192 = 212.

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d is largest divisor of 192 and 404 - 192 = 212. Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).

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d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

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d is largest divisor of 192 and 404 - 192 = 212. Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

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Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

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Idea: Keep subtracting smaller from larger: GCD(404, 192) =

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Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
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Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

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Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) =

```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192)=

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```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192)= GCD(212 - 192, 192) =

```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192)= GCD(212 - 192, 192) = GCD(20, 192).

```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192) = GCD(212 - 192, 192) = GCD(20, 192).Could keep going, but will be subtracting 20's for a while.

```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20.

Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192) = GCD(212 - 192, 192) = GCD(20, 192).Could keep going, but will be subtracting 20's for a while.

Idea: Subtract LOTS of 20's.

```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20. Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192)= GCD(212 - 192, 192) = GCD(20, 192).

Could keep going, but will be subtracting 20's for a while.

Idea: Subtract LOTS of 20's. Largest x: $192 - 20x \ge 0$, x = 9.

```
d is largest divisor of both 404 and 192
IFF
d is largest divisor of 192 and 404 - 192 = 212.
Hence GCD(404,192)=GCD(192,404-192)=GCD(192,212).
```

d is largest divisor of **both** 212 and 192 IFF

d is largest divisor of 212 and 212 - 192 = 20. Hence GCD(212,192)=GCD(212-192,192)=GCD(20,192).

Idea: Keep subtracting smaller from larger: GCD(404, 192) = GCD(404 - 192, 192) = GCD(212, 192) = GCD(212 - 192, 192) = GCD(20, 192).Could keep going, but will be subtracting 20's for a while.

Idea: Subtract LOTS of 20's. Largest $x : 192 - 20x \ge 0, x = 9$. = GCD(20, 192 - 20 × 9 = 12) = GCD(20 - 12, 12) = GCD(8, 12) = GCD(8, 12 - 8 = 4) = GCD(8 - 2 × 4, 4) = GCD(0, 4) = 4.

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 $404=2\times192+20$

 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \end{array}$



 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \end{array}$



 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \end{array}$



 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$

 $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD.

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 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. Can use this to write 4 as a combination of 404 and 192

 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. **Can use this to write 4 as a combination of 404 and 192** Write 4 as a combo of 12's and 8's: $4 = 12 - 1 \times 8$

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 $404 = 2 \times 192 + 20$

 $\begin{array}{l} 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \\ 8 = 4 \times 2 + 0 \text{ STOP HERE and go back one: 4 is the GCD.} \\ \hline \textbf{Can use this to write 4 as a combination of 404 and 192} \\ Write 4 as a combo of 12's and 8's: \\ 4 = 12 - 1 \times 8 \\ Write 8 as a combo of 20's and 12's: \\ 4 = 12 - 1 \times (20 - 12) = 2 \times 12 - 1 \times 20 \end{array}$

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 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. Can use this to write 4 as a combination of 404 and 192 Write 4 as a combo of 12's and 8's: $4 = 12 - 1 \times 8$ Write 8 as a combo of 20's and 12's: $4 = 12 - 1 \times (20 - 12) = 2 \times 12 - 1 \times 20$ Write 12 as combo of 192's and 20's: $4 = 2 \times (192 - 9 \times 20) - 1 \times 20 = 2 \times 192 - 19 \times 20$

 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP HERE and go back one: 4 is the GCD. Can use this to write 4 as a combination of 404 and 192 Write 4 as a combo of 12's and 8's: $4 = 12 - 1 \times 8$ Write 8 as a combo of 20's and 12's: $4 = 12 - 1 \times (20 - 12) = 2 \times 12 - 1 \times 20$ Write 12 as combo of 192's and 20's: $4 = 2 \times (192 - 9 \times 20) - 1 \times 20 = 2 \times 192 - 19 \times 20$ Write 20 as a combo of 404 and 192: $4 = 2 \times 192 - 19 \times (404 - 2 \times 192) = 40 \times 192 - 19 \times 404$ **Upshot:** GCD(m, n) is a combo of m and n

 $101 = 2 \times 38 + 25$

 $101 = 2 \times 38 + 25$ $38 = 1 \times 25 + 13$



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101 = 2 \times 38 + 25 
38 = 1 \times 25 + 13 
25 = 1 \times 13 + 12
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```
101 = 2 \times 38 + 25

38 = 1 \times 25 + 13

25 = 1 \times 13 + 12

13 = 12 + 1
```

```
\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \ \mbox{Go back one: } 1 \ \mbox{is the GCD.} \end{array}
```

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 $\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \end{array} \text{ Go back one: } 1 \text{ is the GCD.} \end{array}$

 $1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25$

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$$\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \end{array}$$
 Go back one: 1 is the GCD.

$$\begin{array}{l} 1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25 \\ 1 = 2(38 - 25) - 25 = 2 \times 38 - 3 \times 25 \end{array}$$

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$$\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \end{array} \text{ Go back one: } 1 \text{ is the GCD.} \end{array}$$

$$1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25$$

$$1 = 2(38 - 25) - 25 = 2 \times 38 - 3 \times 25$$

$$1 = 2 \times 38 - 3 \times (101 - 2 \times 38) = 8 \times 38 - 3 \times 101$$

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$$\begin{array}{l} 101 = 2 \times 38 + 25 \\ 38 = 1 \times 25 + 13 \\ 25 = 1 \times 13 + 12 \\ 13 = 12 + 1 \\ 12 = 12 \times 1 + 0. \end{array} \text{ Go back one: } 1 \text{ is the GCD.} \end{array}$$

$$1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25$$

$$1 = 2(38 - 25) - 25 = 2 \times 38 - 3 \times 25$$

$$1 = 2 \times 38 - 3 \times (101 - 2 \times 38) = 8 \times 38 - 3 \times 101$$

$$1 = 8 \times 38 - 3 \times 101$$

Why is this interesting? Hint: What was our original goal?

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$$\begin{split} 101 &= 2 \times 38 + 25 \\ 38 &= 1 \times 25 + 13 \\ 25 &= 1 \times 13 + 12 \\ 13 &= 12 + 1 \\ 12 &= 12 \times 1 + 0. \end{split}$$
 Go back one: 1 is the GCD.

$$\begin{array}{l} 1 = 13 - 12 = 13 - (25 - 13) = 2 \times 13 - 25 \\ 1 = 2(38 - 25) - 25 = 2 \times 38 - 3 \times 25 \\ 1 = 2 \times 38 - 3 \times (101 - 2 \times 38) = 8 \times 38 - 3 \times 101 \\ 1 = 8 \times 38 - 3 \times 101 \\ \end{array}$$

Why is this interesting? Hint: What was our original goal?
Take both sides mod 101
 $1 \equiv 8 \times 38 \pmod{101}$

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$$\begin{split} 101 &= 2 \times 38 + 25 \\ 38 &= 1 \times 25 + 13 \\ 25 &= 1 \times 13 + 12 \\ 13 &= 12 + 1 \\ 12 &= 12 \times 1 + 0. \end{split}$$
 Go back one: 1 is the GCD.

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Why is this interesting? Hint: What was our original goal?
Take both sides mod 101
$$1 \equiv 8 \times 38 \pmod{101}$$

8 is the inverse of 38 mod 101

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$\operatorname{GCD}(x, \mathbf{0})$

Two things about GCD I want to clarify.

• Why is GCD(x, 0) = x for $x \ge 1$?

When does the algorithm stop?

 $404 = 2 \times 192 + 20$



 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \end{array}$



 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \end{array}$



```
404 = 2 \times 192 + 20

192 = 9 \times 20 + 12

20 = 1 \times 12 + 8

12 = 1 \times 8 + 4
```



 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP WHEN GET 0. Go back one: 4 is GCD.

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```
404 = 2 \times 192 + 20

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8 = 4 \times 2 + 0 STOP WHEN GET 0. Go back one: 4 is GCD.

Lets look at what the algorithm actually does:
```

 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP WHEN GET 0. Go back one: 4 is GCD. Lets look at what the algorithm actually does:

 $GCD(404, 192) = GCD(404 - 2 \times 192, 192) = GCD(20, 192) =$

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 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP WHEN GET 0. Go back one: 4 is GCD. Lets look at what the algorithm actually does:

 $GCD(404, 192) = GCD(404 - 2 \times 192, 192) = GCD(20, 192) =$ $GCD(20, 192 - 9 \times 20) = GCD(20, 12) = GCD(20 - 1 \times 12, 12) =$

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 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \\ 8 = 4 \times 2 + 0 \text{ STOP WHEN GET 0. Go back one: 4 is GCD.} \\ \text{Lets look at what the algorithm actually does:} \\ \text{GCD}(404, 192) = \text{GCD}(404 - 2 \times 192, 192) = \text{GCD}(20, 192) = \\ \text{GCD}(20, 192 - 9 \times 20) = \text{GCD}(20, 12) = \text{GCD}(20 - 1 \times 12, 12) = \\ \end{array}$

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GCD(8, 12) = GCD(8, 12 - 8) = GCD(8, 4) =

 $404 = 2 \times 192 + 20$ $192 = 9 \times 20 + 12$ $20 = 1 \times 12 + 8$ $12 = 1 \times 8 + 4$ $8 = 4 \times 2 + 0$ STOP WHEN GET 0. Go back one: 4 is GCD. Lets look at what the algorithm actually does:

 $\begin{array}{l} \operatorname{GCD}(404,192) = \operatorname{GCD}(404 - 2 \times 192,192) = \operatorname{GCD}(20,192) = \\ \operatorname{GCD}(20,192 - 9 \times 20) = \operatorname{GCD}(20,12) = \operatorname{GCD}(20 - 1 \times 12,12) = \\ \operatorname{GCD}(8,12) = \operatorname{GCD}(8,12 - 8) = \operatorname{GCD}(8,4) = \\ \operatorname{GCD}(8 - 2 \times 4,4) = \operatorname{GCD}(0,4) \end{array}$

 $\begin{array}{l} 404 = 2 \times 192 + 20 \\ 192 = 9 \times 20 + 12 \\ 20 = 1 \times 12 + 8 \\ 12 = 1 \times 8 + 4 \\ 8 = 4 \times 2 + 0 \end{array}$ STOP WHEN GET 0. Go back one: 4 is GCD.

Lets look at what the algorithm actually does: $GCD(404, 192) = GCD(404 - 2 \times 192, 192) = GCD(20, 192) =$ $GCD(20, 192 - 9 \times 20) = GCD(20, 12) = GCD(20 - 1 \times 12, 12) =$ GCD(8, 12) = GCD(8, 12 - 8) = GCD(8, 4) = $GCD(8 - 2 \times 4, 4) = GCD(0, 4)$

To make our formula GCD(x, y) = GCD(x - ky, x) work all the way to 0, we define GCD(0, x) = x.

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Why is $\operatorname{GCD}(0, x) = x$?



Why is GCD(0, x) = x? This is a more interesting question than it appears.

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Why is $\operatorname{GCD}(\mathbf{0}, \mathbf{x}) = \mathbf{x}$?

This is a more interesting question than it appears.

Or I am going to make a point about math inspired by the question.



Why is GCD(0, x) = x? This is a more interesting question than it appears.

Or I am going to make a point about math inspired by the question.

First a short detour: why is $5^{1/2} = \sqrt{5}$?

Why is $5^{1/2} = \sqrt{5}$?

Why is

$$5^{1/2} = \sqrt{5}?$$

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Are we multiplying a number by itself half a time?

Why is $5^{1/2} = \sqrt{5}$?

Why is

$$5^{1/2} = \sqrt{5}?$$

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Are we multiplying a number by itself half a time? Discuss.

Why is $5^{1/2} = \sqrt{5}$?

Why is

$$5^{1/2} = \sqrt{5}?$$

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Are we multiplying a number by itself half a time? Discuss. No.

Why is $5^{1/2} = \sqrt{5}$?

$$5^{1/2} = \sqrt{5}?$$

Are we multiplying a number by itself half a time? Discuss. No. For $a, b \in \mathbb{N}$ we have

$$5^a \times 5^b = 5^{a+b}.$$

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We want this rule to still apply when $a, b \in \mathbb{Q}$.

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We want this rule to still apply when $a, b \in \mathbb{Q}$. So we want

$$5^{1/2} \times 5^{1/2} = 5^{1/2+1/2} = 5$$

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Why is $5^{1/2} = \sqrt{5}$?

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$$5^{1/2} \times 5^{1/2} = 5^{1/2+1/2} = 5$$

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How is 5^{π} defined? Discuss.

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$$5^{3.14159} < 5^{\pi} < 5^{3.141593}.$$

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Need to prove that all choices of sequences yield the same result. We won't do that here

START HERE ON SEPT 7

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 \blacktriangleright 5^{*i*} I leave to you to look up or derive.

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For a story about me, my Dad, and π see https://blog.computationalcomplexity.org/2019/06/ a-proof-that-227-pi-0-and-more.html

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1. Key space is $K = \{(a, b) : 0 \le a, b \le 25, a \text{ is rel prime to } 26\}.$

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Next Slide.

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(i) apply $ax + b$ to T to obtain $T_{a,b}$.

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The Quadratic Ciphers

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- 1. Test takes too long.
- 2. Quad Cipher not secure enough to be worth the time.

History of the Quadratic Cipher

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5. There are other reasons they are wrong.

BILL STOP RECORDING THIS LECTURE

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