BILL RECORD THIS LECTURE

September 9, 2021

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Gen Sub Cipher: How to Really Crack

September 9, 2021

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General Substitution Cipher

Def Gen Sub Cipher with perm f on $\{0, \ldots, 25\}$.

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- 1. Encrypt via $x \to f(x)$.
- 2. Decrypt via $x \to f^{-1}(x)$.

Terminology: 1-Gram, 2-Gram, 3-Gram

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- 3. The **3**-grams of T you can guess. Also called trigrams.
- 4. One usually talks about the freq of *n*-grams.

Let the text be:

Ever notice how sometimes people use math words incorrectly?

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- The following 1-gram occurs 9 times: e.
- No 1-gram occurs \geq 10 times.

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The following 2-grams occur 2 times: me, or. The following 2-grams occur 1 time: ev, ve, er, rn, no, ot, ti, ic, eh, ho, ow, ws, so, et, ti, im, es, sp, pe, eo, op, pl, le, eu, us, se, em, ma, at, th, hw, wo, ds, in, nc, co, rr, re, ec, ct, tl, ly.

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1. 1-grams: $f_{E,1} \cdot f_{E,1} \sim 0.065$.

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2. 2-grams: f_{E,2} · f_{E,2} ~ 0.0067.

- 1. 1-grams: $f_{E,1} \cdot f_{E,1} \sim 0.065$.
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- 3. 3-grams: $f_{E,3} \cdot f_{E,3} \sim 0.0011$.

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- 4. 4-grams: $f_{E,4} \cdot f_{E,4} \sim 0.00023$.

Contrast Shift to Gen Sub

To crack shift went through all 26 shifts σ :

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3. Important Will always be large or small. So we have a gap.

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- 2. If $f_{\sigma(T),1} \cdot f_{E,1}$ is small then σ is incorrect perm. Small. Hmmm?
- We have a problem. If σ only changed a few letters around, then likely f_{E,1} · f_{σ(T),1} will be large. We do not have a gap!
 What to do?

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1. Use *n*-grams instead of 1-grams.

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- 2. If σ is a perm and $n \in \mathbb{N}$ then

$$\operatorname{good}_{\sigma,n} = f_{E,n} \cdot f_{\sigma(T),n}.$$

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3. Rather than view the **Is-English** program as a YES-NO, view it as comparative:

 T_1 looks more like English than T_2 .

Input T. Find Freq of 1-grams and *n*-grams.

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An old question: What came first, the chicken or the egg?

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Our Problem We need parameters I and R so the answer looks like English. But we then need a notion of **Is English** that does not use a gap. Need a program to tell us that it looks like English.

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We find the parameters for texts where we know the answers.

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5. Keep track of how how many iterations suffice and how many redos suffice.

David Zhen Found the Parameters

UMCP ugrad CS major David Zhen worked with me on this over the summer.

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For each text he generated many random perm and ran the algorithm.

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1-grams Nothing worked



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So the winner is 3-grams, with I = 2000 and R = 4.

Can we do better than 2 minutes? Can we do something clever?

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There were issues with his work so I would want to see this redone more carefully. However, I suspect

BILL STOP RECORDING THIS LECTURE

September 9, 2021

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