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September 9, 2021

# Gen Sub Cipher: How to Really Crack 

September 9, 2021

## General Substitution Cipher

Def Gen Sub Cipher with perm $f$ on $\{0, \ldots, 25\}$.

1. Encrypt via $x \rightarrow f(x)$.
2. Decrypt via $x \rightarrow f^{-1}(x)$.

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4. One usually talks about the freq of $n$-grams.

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No 1-gram occurs $\geq 10$ times.

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No 2-gram occurs $\geq 3$ times.

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3. We have a problem. If $\sigma$ only changed a few letters around, then likely $f_{E, 1} \cdot f_{\sigma(T), 1}$ will be large. We do not have a gap!
What to do?

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3. Rather than view the Is-English program as a YES-NO, view it as comparative:
$T_{1}$ looks more like English than $T_{2}$.

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Candidates for $\sigma$ are $\sigma_{1}, \ldots, \sigma_{\mathrm{R}}$
Pick the $\sigma_{r}$ with max $\operatorname{good}_{\sigma_{r}, n}$ or have human look at all $\sigma_{r}(T)$

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Same here.
We find the parameters for texts where we know the answers.

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5. Keep track of how how many iterations suffice and how many redos suffice.

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For each text he generated many random perm and ran the algorithm.

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Can we do better than 2 minutes? Can we do something clever?

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There were issues with his work so I would want to see this redone more carefully. However, I suspect

# BILL <br> STOP RECORDING THIS LECTURE 

September 9, 2021

