Schur's Thm + FLT implies Primes Infinite

May 8, 2025

◆□▶ ◆□▶ ◆□▶ ◆□▶ □ - つくぐ

The following people have used Ramsey Theory to show Primes ∞ .

The following people have used Ramsey Theory to show Primes ∞ .

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

1. Alpoge (2015) used Intermediary NT (INT) and VDW.

The following people have used Ramsey Theory to show Primes $\infty.$

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.

The following people have used Ramsey Theory to show Primes ∞ .

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.
- 3. Elsholtz (2021) used INT and Schur's Thm.

The following people have used Ramsey Theory to show Primes ∞ .

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.
- 3. Elsholtz (2021) used INT and Schur's Thm.
- 4. Goral, Ozcan, Serbas (2022) used HNT and Poly-VDW Thm.

The following people have used Ramsey Theory to show Primes ∞ .

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.
- 3. Elsholtz (2021) used INT and Schur's Thm.
- 4. Goral, Ozcan, Serbas (2022) used HNT and Poly-VDW Thm.

5. Gasarch (2023) used INT and Schur's Thm.

The following people have used Ramsey Theory to show Primes ∞ .

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.
- 3. Elsholtz (2021) used INT and Schur's Thm.
- 4. Goral, Ozcan, Serbas (2022) used HNT and Poly-VDW Thm.

- 5. Gasarch (2023) used INT and Schur's Thm.
- 6. We refer to the proof by Elsholtz and Gasarch as the **EG-proof**.

The following people have used Ramsey Theory to show Primes ∞ .

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.
- 3. Elsholtz (2021) used INT and Schur's Thm.
- 4. Goral, Ozcan, Serbas (2022) used HNT and Poly-VDW Thm.

- 5. Gasarch (2023) used INT and Schur's Thm.
- 6. We refer to the proof by Elsholtz and Gasarch as the **EG-proof**.
- 1. All of these proofs are harder than the usual proof

The following people have used Ramsey Theory to show Primes ∞ .

- 1. Alpoge (2015) used Intermediary NT (INT) and VDW.
- 2. Granville (2017) used INT and VDW.
- 3. Elsholtz (2021) used INT and Schur's Thm.
- 4. Goral, Ozcan, Serbas (2022) used HNT and Poly-VDW Thm.
- 5. Gasarch (2023) used INT and Schur's Thm.
- 6. We refer to the proof by Elsholtz and Gasarch as the **EG-proof**.
- 1. All of these proofs are harder than the usual proof
- 2. All of these proofs have other points to make after they prove primes ∞ .

・ロト・個ト・ヨト・ヨト ヨー りへぐ

1. Present the EG-Proof since its the one I know best.

<□▶ <□▶ < □▶ < □▶ < □▶ < □▶ < □ > ○ < ○

- 1. Present the EG-Proof since its the one I know best.
- Look at what it means to ask the question in domains other than N. (In fact, asking it over N is not quite right).

- 1. Present the EG-Proof since its the one I know best.
- Look at what it means to ask the question in domains other than N. (In fact, asking it over N is not quite right).
- 3. Look at domains where the number of primes is finite and see where the standard proof fails, and where the EG-proof fails.

ション ふゆ アメビア メロア しょうくしゃ

Background Needed For EG-Proof

May 8, 2025

▲ロト ▲園 ト ▲ 臣 ト ▲ 臣 ト ● ○ ○ ○ ○ ○

Notation Let $k, n \in \mathbb{N} - \{0\}$.



Notation Let $k, n \in \mathbb{N} - \{0\}$. Let A be any set (it can even be ∞).

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Notation Let $k, n \in \mathbb{N} - \{0\}$. Let A be any set (it can even be ∞).

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

1.
$$[n] = \{1, 2, \ldots, n\}.$$

Notation Let $k, n \in \mathbb{N} - \{0\}$. Let A be any set (it can even be ∞).

- 1. $[n] = \{1, 2, \ldots, n\}.$
- 2. $\binom{A}{k}$ is the set of all subsets of A of size k.

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists *x*, *y*, *z* monochromatic such that x + y = z.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. Pf We determine S later.

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. **Pf** We determine *S* later. Given COL we define $\text{COL'}\binom{[S]}{2} \rightarrow [c]$ as follows:

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. **Pf** We determine *S* later. Given COL we define $\text{COL'}\binom{[S]}{2} \rightarrow [c]$ as follows:

$$\operatorname{COL}'(x, y) = \operatorname{COL}(|x - y|).$$

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. **Pf** We determine *S* later. Given COL we define $\text{COL'}\binom{[S]}{2} \rightarrow [c]$ as follows:

$$\operatorname{COL}'(x, y) = \operatorname{COL}(|x - y|).$$

There exists a COL'-homog set H of size 3 (thats all we need!). Say its a < b < c

$$\operatorname{COL}'(c, b) = \operatorname{COL}'(b, a) = \operatorname{COL}'(c, a)$$

ション ふぼう メリン メリン しょうくしゃ

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. **Pf** We determine *S* later. Given COL we define $\text{COL'}\binom{[S]}{2} \rightarrow [c]$ as follows:

$$\operatorname{COL}'(x, y) = \operatorname{COL}(|x - y|).$$

There exists a COL'-homog set H of size 3 (thats all we need!). Say its a < b < c

$$\operatorname{COL}'(c, b) = \operatorname{COL}'(b, a) = \operatorname{COL}'(c, a)$$

So

$$\operatorname{COL}(c-b) = \operatorname{COL}(b-a) = \operatorname{COL}(c-a)$$

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. **Pf** We determine *S* later. Given COL we define $\text{COL'}\binom{[S]}{2} \rightarrow [c]$ as follows:

$$\operatorname{COL}'(x, y) = \operatorname{COL}(|x - y|).$$

There exists a COL'-homog set H of size 3 (thats all we need!). Say its a < b < c

$$\operatorname{COL}'(c, b) = \operatorname{COL}'(b, a) = \operatorname{COL}'(c, a)$$

So

$$\operatorname{COL}(c-b) = COL(b-a) = COL(c-a)$$

Let x = c - b, y = b - a, z = c - a.

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists x, y, z monochromatic such that x + y = z. **Pf** We determine *S* later. Given COL we define $\text{COL'}\binom{[S]}{2} \rightarrow [c]$ as follows:

$$\operatorname{COL}'(x, y) = \operatorname{COL}(|x - y|).$$

There exists a COL'-homog set H of size 3 (thats all we need!). Say its a < b < c

$$\operatorname{COL}'(c, b) = \operatorname{COL}'(b, a) = \operatorname{COL}'(c, a)$$

So

$$\operatorname{COL}(c-b) = COL(b-a) = COL(c-a)$$

Let x = c - b, y = b - a, z = c - a. So let S(c) = R(3; c) (homog set 3, colors c).

Fermat's Last Theorem

In 1637 Fermat wrote in the margins of **Arithmetica**, a book on Number Theory by Diophantus, the following (translated from Latin)

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Fermat's Last Theorem

In 1637 Fermat wrote in the margins of **Arithmetica**, a book on Number Theory by Diophantus, the following (translated from Latin)

To divide a cube into two cubes, a fourth power, or in general any power whatever above the second into two powers of the same denomination, is impossible, and I have assuredly found a proof of this, but the margin is too narrow to contain it.

ション ふぼう メリン メリン しょうくしゃ

Fermat's Last Theorem

In 1637 Fermat wrote in the margins of **Arithmetica**, a book on Number Theory by Diophantus, the following (translated from Latin)

To divide a cube into two cubes, a fourth power, or in general any power whatever above the second into two powers of the same denomination, is impossible, and I have assuredly found a proof of this, but the margin is too narrow to contain it. In modern terminology:

$$(\forall n \geq 3)(\forall x, y, z \in \mathbb{N} - \{0\})[x^n + y^n \neq z^n].$$

This has come to be known as Fermat's Last Theorem.

Did Fermat Have a Proof? Arguments Against

- * ロ > * 週 > * 注 > * 注 > ・ 注 - の < @

Did Fermat Have a Proof? Arguments Against

1) He proved the n = 4 case later in his life. He would not have done this if he had earlier proved the full theorem.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

1) He proved the n = 4 case later in his life. He would not have done this if he had earlier proved the full theorem.

2) Andrew Wiles proved FLT in the early 1990s with techniques far beyond what Fermat could have known.

Did Fermat Have a Proof? Arguments For

- イロト イボト イモト - モー のへぐ

Did Fermat Have a Proof? Arguments For

1) The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Did Fermat Have a Proof? Arguments For

The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.
The 11th Dr. Who gave The real proof to a group of geniuses to gain their trust.

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 のへぐ
- The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.
 The 11th Dr. Who gave The real proof to a group of geniuses
- 2) The 11th Dr. Who gave The real proof to a group of geniuses to gain their trust.
 - 1. He later said that it was Fermat's original proof (possible but unlikely),

- 1) The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.
- 2) The 11th Dr. Who gave **The real proof** to a group of geniuses to gain their trust.
 - 1. He later said that it was Fermat's original proof (possible but unlikely),
 - but that Fermat didn't write it down since he died in a duel (not true).

- 1) The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.
- 2) The 11th Dr. Who gave **The real proof** to a group of geniuses to gain their trust.
 - 1. He later said that it was Fermat's original proof (possible but unlikely),
 - but that Fermat didn't write it down since he died in a duel (not true). The writers of the show either

ション ふぼう メリン メリン しょうくしゃ

- 1) The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.
- 2) The 11th Dr. Who gave **The real proof** to a group of geniuses to gain their trust.
 - 1. He later said that it was Fermat's original proof (possible but unlikely),
 - but that Fermat didn't write it down since he died in a duel (not true). The writers of the show either

ション ふぼう メリン メリン しょうくしゃ

2.1 confused Galois with Fermat, or

- 1) The 7th Dr. Who had a 5-line proof that uses Boolean Algebra.
- 2) The 11th Dr. Who gave **The real proof** to a group of geniuses to gain their trust.
 - 1. He later said that it was Fermat's original proof (possible but unlikely),
 - but that Fermat didn't write it down since he died in a duel (not true). The writers of the show either
 - 2.1 confused Galois with Fermat, or
 - 2.2 meant to say that Fermat died in a duel in a dual timeline.

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ - 三 - のへぐ

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN. **Whoops**

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ - 三 - のへぐ

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN. **Whoops**

In **Star Trek: DSN**, the episode **Facets** which aired on June 12, 1995, Dax says that one of her previous hosts, Tobin, had done *the most creative work on Fermat's Last Theorem since Wiles*.

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN. **Whoops**

In **Star Trek: DSN**, the episode **Facets** which aired on June 12, 1995, Dax says that one of her previous hosts, Tobin, had done *the most creative work on Fermat's Last Theorem since Wiles*.

My guess is that Tobin wrote this limerick:

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN. **Whoops**

In **Star Trek: DSN**, the episode **Facets** which aired on June 12, 1995, Dax says that one of her previous hosts, Tobin, had done *the most creative work on Fermat's Last Theorem since Wiles*.

My guess is that Tobin wrote this limerick: A challenge for many long ages Had baffled the savants and sages

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN. **Whoops**

In **Star Trek: DSN**, the episode **Facets** which aired on June 12, 1995, Dax says that one of her previous hosts, Tobin, had done *the most creative work on Fermat's Last Theorem since Wiles*.

My guess is that Tobin wrote this limerick:

A challenge for many long ages Had baffled the savants and sages

Yet at last came the light Seems that Fermat was right

In **Star Trek: TNG**, the episode **The Royale** which aired on March 27, 1989, Captain Picard, in the 24th Century is working on Fermat's Last Theorem, which is still OPEN. **Whoops**

In **Star Trek: DSN**, the episode **Facets** which aired on June 12, 1995, Dax says that one of her previous hosts, Tobin, had done *the most creative work on Fermat's Last Theorem since Wiles*.

My guess is that Tobin wrote this limerick:

A challenge for many long ages Had baffled the savants and sages

Yet at last came the light Seems that Fermat was right

To the margin add 200 pages.

May 8, 2025

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Thm The number of primes is infinite.

▲□▶▲圖▶▲圖▶▲圖▶ 圖 のへで

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L .

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

 $\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

 $\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L})=(a_1\pmod{4},\ldots,a_L\pmod{4})$

By Schur's Thm there exists x, y, z same color with x + y = z.

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L})=(a_1\pmod{4},\ldots,a_L\pmod{4})$$

ション ふぼう メリン メリン しょうくしゃ

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \ldots, a_L \pmod{4})$$

ション ふぼう メリン メリン しょうくしゃ

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= p_1^{4x_1+e_1} \cdots p_L^{4x_L+e_L} \\ y &= p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L} \\ z &= p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L} \end{aligned}$$

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \ldots, a_L \pmod{4})$$

ション ふぼう メリン メリン しょうくしゃ

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

 $x = p_1^{4x_1+e_1} \cdots p_L^{4x_L+e_L}$ $y = p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L}$ $z = p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L}$ x + y = z

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= p_1^{4x_1+e_1} \cdots p_L^{4x_L+e_L} \\ y &= p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L} \\ z &= p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L} \\ x+y &= z \\ p_1^{4x_1+e_1} \cdots p_1^{4x_L+e_L} + p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L} = p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L} \end{aligned}$$

▲□▶▲□▶▲□▶▲□▶ = のへの

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= p_1^{4x_1+e_1} \cdots p_L^{4x_L+e_L} \\ y &= p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L} \\ z &= p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L} \end{aligned}$$

$$\begin{aligned} x + y &= z \\ p_1^{4x_1 + e_1} \cdots p_L^{4x_L + e_L} + p_1^{4y_1 + e_1} \cdots p_L^{4y_L + e_L} &= p_1^{4z_1 + e_1} \cdots p_L^{4z_n + e_L} \\ p_1^{4x_1} \cdots p_L^{4x_L} + p_1^{4y_1} \cdots p_L^{4y_L} &= p_1^{4z_1} \cdots p_L^{4z_n} \end{aligned}$$

・ロト・西ト・ヨト・ヨー うらぐ

Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$x = p_1^{4x_1+e_1} \cdots p_L^{4x_L+e_L}$$

$$y = p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L}$$

$$z = p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L}$$

$$\begin{aligned} x + y &= z \\ p_1^{4x_1+e_1} \cdots p_L^{4x_L+e_L} + p_1^{4y_1+e_1} \cdots p_L^{4y_L+e_L} &= p_1^{4z_1+e_1} \cdots p_L^{4z_n+e_L} \\ p_1^{4x_1} \cdots p_L^{4x_L} + p_1^{4y_1} \cdots p_L^{4y_L} &= p_1^{4z_1} \cdots p_L^{4z_n} \\ (p_1^{x_1} \cdots p_L^{x_L})^4 + (p_1^{y_1} \cdots p_L^{y_L})^4 &= (p_1^{z_1} \cdots p_L^{z_L})^4 \\ \end{aligned}$$
This violates FLT for $n = 4$.

◆□▶ ◆□▶ ◆三▶ ◆三▶ ◆□ ◆ ◆

How to Ask the Question of Primes Infinite

May 8, 2025

*ロト *目 * * * * * * * * * * * * * * *

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへの

1. *D* is closed under + and \times .

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

- 1. *D* is closed under + and \times .
- 2. There is an element $0 \in D$ such that $(\forall x \in D)[x + 0 = x]$.

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

- 1. *D* is closed under + and \times .
- 2. There is an element $0 \in D$ such that $(\forall x \in D)[x + 0 = x]$.
- 3. There is an element $1 \in D$ such that $(\forall x \in D)[x \times 1 = x]$.

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

- 1. *D* is closed under + and \times .
- 2. There is an element $0 \in D$ such that $(\forall x \in D)[x + 0 = x]$.
- 3. There is an element $1 \in D$ such that $(\forall x \in D)[x \times 1 = x]$.

4. + and \times are communicative and associative.

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

- 1. *D* is closed under + and \times .
- 2. There is an element $0 \in D$ such that $(\forall x \in D)[x + 0 = x]$.
- 3. There is an element $1 \in D$ such that $(\forall x \in D)[x \times 1 = x]$.

4. + and \times are communicative and associative.

5. (Key) $(\forall x)(\exists y)[x+y=0].$

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

- 1. *D* is closed under + and \times .
- 2. There is an element $0 \in D$ such that $(\forall x \in D)[x + 0 = x]$.
- 3. There is an element $1 \in D$ such that $(\forall x \in D)[x \times 1 = x]$.

4. + and \times are communicative and associative.

5. (Key)
$$(\forall x)(\exists y)[x + y = 0]$$

6. (Key) If ab = 0 then either a = 0 or b = 0.

Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

- 1. *D* is closed under + and \times .
- 2. There is an element $0 \in D$ such that $(\forall x \in D)[x + 0 = x]$.
- 3. There is an element $1 \in D$ such that $(\forall x \in D)[x \times 1 = x]$.

4. + and \times are communicative and associative.

5. (Key)
$$(\forall x)(\exists y)[x+y=0]$$
.

6. (Key) If ab = 0 then either a = 0 or b = 0.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide. **Integral Domains**

Upshot $+,\times,0,1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

1) Z.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

Integral Domains

1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) Z. CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} .

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.
- 3) Algebraic Numbers

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) $\mathbb{Q},$ $\mathbb{R},$ $\mathbb{C}.$ CAN divide.

3) Algebraic Numbers

 $\mathbb{AN} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x])) [f(a) = 0].$

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

ション ふゆ アメビア メロア しょうくしゃ

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

3) Algebraic Numbers

 $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0]\}.$ CANNOT DIVIDE.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0]\}.$ CANNOT DIVIDE.

Proof that \mathbb{AI} is closed under + and \times is hard.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.
- 3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0]\}.$ CANNOT DIVIDE.

Proof that \mathbb{AI} is closed under + and \times is hard.

5)
$$\left\{\frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2}\right\}.$$

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) \mathbb{Q} , \mathbb{R} , \mathbb{C} . CAN divide.

3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0]\}.$ CANNOT DIVIDE.

Proof that \mathbb{AI} is closed under + and \times is hard.

5) $\left\{\frac{a}{b} : \operatorname{gcd}(a, b) = 1 \land b \equiv 1 \pmod{2}\right\}$. CANNOT DIVIDE. $\neg \exists \frac{1}{2}$.

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) $\mathbb{Q},$ $\mathbb{R},$ $\mathbb{C}.$ CAN divide.

3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0]\}.$ CANNOT DIVIDE.

Proof that \mathbb{AI} is closed under + and \times is hard.

5) $\{\frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2}\}$. CANNOT DIVIDE. $\neg \exists \frac{1}{2}$. 6) $\mathbb{Z}[i] = \{a + bi : a, b \in \mathbb{Z}\}.$

Upshot $+, \times, 0, 1$ act as you expect, you can subtract, you might not be able to divide.

Integral Domains

- 1) \mathbb{Z} . CANNOT divide: $\neg(\exists \frac{1}{3})$.
- 2) $\mathbb{Q},$ $\mathbb{R},$ $\mathbb{C}.$ CAN divide.
- 3) Algebraic Numbers $\mathbb{AN} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x]))[f(a) = 0].$ CAN DIVIDE. Proof that \mathbb{AN} is closed under + and \times is hard.

4) Algebraic integers $\mathbb{AI} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0]\}.$ CANNOT DIVIDE.

Proof that \mathbb{AI} is closed under + and \times is hard.

5) $\{\frac{a}{b} : \operatorname{gcd}(a, b) = 1 \land b \equiv 1 \pmod{2}\}$. CANNOT DIVIDE. $\neg \exists \frac{1}{2}$. 6) $\mathbb{Z}[i] = \{a + bi : a, b \in \mathbb{Z}\}$. CANNOT DIVIDE. $\neg \exists \frac{1}{2}$.

- イロト イ理ト イヨト イヨト ヨー のへぐ

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

1) $\mathbb{Z}_{12}=\{0,\ldots,11\}$ with mod 12 math.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

1) $\mathbb{Z}_{12} = \{0, \dots, 11\}$ with mod 12 math. Note that $3 \times 4 = 0$ but $3 \neq 0$ and $4 \neq 0$.

1) $\mathbb{Z}_{12} = \{0, \ldots, 11\}$ with mod 12 math. Note that $3 \times 4 = 0$ but $3 \neq 0$ and $4 \neq 0$. (Note: \mathbb{Z}_n is an integral domain iff *n* is prime.)

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

1) $\mathbb{Z}_{12} = \{0, \ldots, 11\}$ with mod 12 math. Note that $3 \times 4 = 0$ but $3 \neq 0$ and $4 \neq 0$. (Note: \mathbb{Z}_n is an integral domain iff *n* is prime.)

2) N. There is no -3.

1) $\mathbb{Z}_{12} = \{0, \ldots, 11\}$ with mod 12 math. Note that $3 \times 4 = 0$ but $3 \neq 0$ and $4 \neq 0$. (Note: \mathbb{Z}_n is an integral domain iff *n* is prime.)

2) N. There is no -3.

We will look at which Integral Domains have an infinite number of primes.

1) $\mathbb{Z}_{12} = \{0, \ldots, 11\}$ with mod 12 math. Note that $3 \times 4 = 0$ but $3 \neq 0$ and $4 \neq 0$. (Note: \mathbb{Z}_n is an integral domain iff *n* is prime.)

2)
$$\mathbb{N}$$
. There is no -3 .

We will look at which Integral Domains have an infinite number of primes.

ション ふゆ アメビア メロア しょうくしゃ

Will need to ask the question carefully.

*ロト *昼 * * ミ * ミ * ミ * のへぐ

Def Let \mathbb{D} be an integral domain.

Def Let \mathbb{D} be an integral domain.

1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Def Let \mathbb{D} be an integral domain.

1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- 2. An irreducible is a $p \in \mathbb{D} \mathbb{U}$ such that if p = ab then either $a \in \mathbb{U}$ or $b \in \mathbb{U}$.

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- 2. An irreducible is a $p \in \mathbb{D} \mathbb{U}$ such that if p = ab then either $a \in \mathbb{U}$ or $b \in \mathbb{U}$.

We let ${\mathbb I}$ be the set of irreducibles.

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- An irreducible is a p ∈ D − U such that if p = ab then either a ∈ U or b ∈ U.
 We let I be the set of irreducibles.
- 3. A **prime** is a $p \in \mathbb{D} \mathbb{U}$ such that if p divides ab then either p divides a or p divides b.

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- An irreducible is a p ∈ D − U such that if p = ab then either a ∈ U or b ∈ U.
 We let I be the set of irreducibles.
- 3. A **prime** is a $p \in \mathbb{D} \mathbb{U}$ such that if p divides ab then either p divides a or p divides b.

In any integral domain all primes are irreducible but not all irreducibles are primes. (We will not be getting into that).

ション ふゆ アメビア メロア しょうくしゃ

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- An irreducible is a p ∈ D − U such that if p = ab then either a ∈ U or b ∈ U.
 We let I be the set of irreducibles.
- 3. A **prime** is a $p \in \mathbb{D} \mathbb{U}$ such that if p divides ab then either p divides a or p divides b.

In any integral domain all primes are irreducible but not all irreducibles are primes. (We will not be getting into that).

4. A composite is an $n \in \mathbb{D} - \mathbb{U}$ such that there exists $a, b \in \mathbb{D} - \mathbb{U}$, n = ab.

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- An irreducible is a p ∈ D − U such that if p = ab then either a ∈ U or b ∈ U.
 We let I be the set of irreducibles.
- 3. A **prime** is a $p \in \mathbb{D} \mathbb{U}$ such that if p divides ab then either p divides a or p divides b.

In any integral domain all primes are irreducible but not all irreducibles are primes. (We will not be getting into that).

4. A composite is an $n \in \mathbb{D} - \mathbb{U}$ such that there exists $a, b \in \mathbb{D} - \mathbb{U}$, n = ab.

Units are **not** irreducibles. This is why 1, -1 are not primes.

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- An irreducible is a p ∈ D − U such that if p = ab then either a ∈ U or b ∈ U.
 We let I be the set of irreducibles.
- 3. A **prime** is a $p \in \mathbb{D} \mathbb{U}$ such that if p divides ab then either p divides a or p divides b.

In any integral domain all primes are irreducible but not all irreducibles are primes. (We will not be getting into that).

4. A composite is an $n \in \mathbb{D} - \mathbb{U}$ such that there exists $a, b \in \mathbb{D} - \mathbb{U}$, n = ab.

Units are **not** irreducibles. This is why 1, -1 are not primes. We will be concerned with irreducibles, not primes.

Def Let \mathbb{D} be an integral domain.

- 1. A unit is a $u \in \mathbb{D}$ such that there exists $v \in \mathbb{D}$ with uv = 1. We let \mathbb{U} be the set of units.
- An irreducible is a p ∈ D − U such that if p = ab then either a ∈ U or b ∈ U.
 We let I be the set of irreducibles.
- 3. A **prime** is a $p \in \mathbb{D} \mathbb{U}$ such that if p divides ab then either p divides a or p divides b.

In any integral domain all primes are irreducible but not all irreducibles are primes. (We will not be getting into that).

4. A composite is an $n \in \mathbb{D} - \mathbb{U}$ such that there exists $a, b \in \mathbb{D} - \mathbb{U}$, n = ab.

Units are **not** irreducibles. This is why 1, -1 are not primes. We will be concerned with irreducibles, not primes.

Types of Elts in an ID 0, units, irreducibles, composites.

1) Domain is \mathbb{Z} . Are 7 and -7 DIFFERENT irreducibles? Discuss

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[i]. Are 7, −7, 7i, −7i DIFFERENT irreducibles? Discuss

・ロト・日本・モト・モト・モー うへぐ

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[i]. Are 7, −7, 7i, −7i DIFFERENT irreducibles? Discuss

3) Domain is $\mathbb{CI} = \mathbb{Z}[\{e^{2\pi i k/n} : 0 \le k \le n\}]$. $e^{2\pi i k/n}$'s are all units.

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[i]. Are 7, −7, 7i, −7i DIFFERENT irreducibles? Discuss

3) Domain is $\mathbb{CI} = \mathbb{Z}[\{e^{2\pi ik/n} : 0 \le k \le n\}]$. $e^{2\pi ik/n}$'s are all units. (\mathbb{CI} stands for Cyclotomic Integers.)

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[*i*]. Are 7, −7, 7*i*, −7*i* DIFFERENT irreducibles? Discuss

3) Domain is $\mathbb{CI} = \mathbb{Z}[\{e^{2\pi ik/n} : 0 \le k \le n\}]$. $e^{2\pi ik/n}$'s are all units. (\mathbb{CI} stands for **Cyclotomic Integers**.) Is the following argument valid or mid:

ション ふぼう メリン メリン しょうくしゃ
Which Irreducibles are Different?

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[*i*]. Are 7, −7, 7*i*, −7*i* DIFFERENT irreducibles? Discuss

3) Domain is $\mathbb{CI} = \mathbb{Z}[\{e^{2\pi ik/n} : 0 \le k \le n\}]$. $e^{2\pi ik/n}$'s are all units. (\mathbb{CI} stands for **Cyclotomic Integers**.) Is the following argument valid or mid: \mathbb{CI} has an infinite number of irreducibles:

$$\{7 \times \zeta_i^n : n \in \mathbb{N}, 0 \le i \le n\}.$$

ション ふゆ アメリア メリア しょうくしゃ

Which Irreducibles are Different?

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[i]. Are 7, −7, 7i, −7i DIFFERENT irreducibles? Discuss

3) Domain is $\mathbb{CI} = \mathbb{Z}[\{e^{2\pi ik/n} : 0 \le k \le n\}]$. $e^{2\pi ik/n}$'s are all units. (\mathbb{CI} stands for **Cyclotomic Integers**.) Is the following argument valid or mid: \mathbb{CI} has an infinite number of irreducibles:

 $\{7 \times \zeta_i^n : n \in \mathbb{N}, 0 \le i \le n\}.$

It seems like this is cheating. Even 7 and -7 seem to be the same.

Which Irreducibles are Different?

Domain is Z. Are 7 and −7 DIFFERENT irreducibles? Discuss
 Domain is Z[i]. Are 7, −7, 7i, −7i DIFFERENT irreducibles? Discuss

3) Domain is $\mathbb{CI} = \mathbb{Z}[\{e^{2\pi ik/n} : 0 \le k \le n\}]$. $e^{2\pi ik/n}$'s are all units. (\mathbb{CI} stands for **Cyclotomic Integers**.) Is the following argument valid or mid: \mathbb{CI} has an infinite number of irreducibles:

 $\{7 \times \zeta_i^n : n \in \mathbb{N}, 0 \le i \le n\}.$

It seems like this is cheating. Even 7 and -7 seem to be **the same**. What to do? Discuss

Equivalence Classes of Irreducibles

Convention Let $\mathbb D$ be an Int Dom with Units $\mathbb U,$ Irreds $\mathbb I.$

(ロト (個) (E) (E) (E) (E) のへの

Equivalence Classes of Irreducibles

Convention Let \mathbb{D} be an Int Dom with Units \mathbb{U} , Irreds \mathbb{I} . We define the following equivalence relation on \mathbb{I} :

$$p \equiv q$$
 iff $(\exists u \in \mathbb{U})[p = uq].$

Equivalence Classes of Irreducibles

Convention Let \mathbb{D} be an Int Dom with Units \mathbb{U} , Irreds \mathbb{I} . We define the following equivalence relation on \mathbb{I} :

$$p \equiv q$$
 iff $(\exists u \in \mathbb{U})[p = uq].$

I is infinite up to units if the number of equivalence classes is infinite.

Convention Let \mathbb{D} be an Int Dom with Units \mathbb{U} , Irreds \mathbb{I} . We define the following equivalence relation on \mathbb{I} :

$$p \equiv q$$
 iff $(\exists u \in \mathbb{U})[p = uq].$

I is infinite up to units if the number of equivalence classes is infinite.

New Question Given \mathbb{D} try to show that \mathbb{D} has an infinite number of equiv classes or irreducibles.

ション ふゆ アメリア メリア しょうくしゃ

Convention Let \mathbb{D} be an Int Dom with Units \mathbb{U} , Irreds \mathbb{I} . We define the following equivalence relation on \mathbb{I} :

$$p \equiv q$$
 iff $(\exists u \in \mathbb{U})[p = uq].$

I is infinite up to units if the number of equivalence classes is infinite.

New Question Given \mathbb{D} try to show that \mathbb{D} has an infinite number of equiv classes or irreducibles.

・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・

On theses slides infinite will mean infinite up to units.

The Normal Proof that Primes are Infinite and Where it Falls Apart

May 8, 2025

ション ふゆ アメビア メロア しょうくしゃ

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.)

◆□▶ ◆□▶ ◆三▶ ◆三▶ ・三 ・ つへぐ

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.) Form $N = p_1 \cdots p_n + 1$.

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.) Form $N = p_1 \cdots p_n + 1$. Two Cases.

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.) Form $N = p_1 \cdots p_n + 1$. Two Cases.

- 1. *N* is prime. **Done** since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.
- N is composite. Then N = ab where a, b ∉ {-1,1}. If a and b are composite then break them down until you get to prime p, p divides N. So N = Mp.

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.) Form $N = p_1 \cdots p_n + 1$. Two Cases.

- 1. *N* is prime. **Done** since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.
- 2. *N* is composite. Then N = ab where $a, b \notin \{-1, 1\}$. If a and b are composite then break them down until you get to prime p, p divides N. So N = Mp.

 $Mp = p_1 \cdots p_n + 1$. Take this mod p.

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.) Form $N = p_1 \cdots p_n + 1$. Two Cases.

- 1. *N* is prime. **Done** since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.
- N is composite. Then N = ab where a, b ∉ {-1,1}. If a and b are composite then break them down until you get to prime p, p divides N. So N = Mp. Mp = p₁ ··· p_n + 1. Take this mod p. 0 ≡ p₁ ··· p_n + 1 (mod p).

ション ふゆ アメビア メロア しょうくしゃ

Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.) Form $N = p_1 \cdots p_n + 1$. Two Cases.

- 1. *N* is prime. **Done** since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.
- 2. N is composite. Then N = ab where a, b ∉ {-1,1}. If a and b are composite then break them down until you get to prime p, p divides N. So N = Mp.
 Mp = p₁ ··· p_n + 1. Take this mod p.
 0 ≡ p₁ ··· p_n + 1 (mod p).
 p ∉ {p₁,..., p_n} since if it was then 0 ≡ 1 (mod p).

 ${\mathbb Q}$ has 0, units, NO primes, NO composites.



 $\mathbb Q$ has 0, units, NO primes, NO composites. Where does proof primes ∞ go wrong? Discuss

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

 $\mathbb Q$ has 0, units, NO primes, NO composites. Where does proof primes ∞ go wrong? Discuss See next slide.

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

If p_1, \ldots, p_n are **any** set of rationals then $N = p_1 p_2 \cdots p_n + 1$ is a **a unit**.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

If p_1, \ldots, p_n are **any** set of rationals then $N = p_1 p_2 \cdots p_n + 1$ is a **a unit**.

Note that in the proof we considered two cases: N is prime.

N is composite.

If p_1, \ldots, p_n are **any** set of rationals then $N = p_1 p_2 \cdots p_n + 1$ is a **a unit**.

Note that in the proof we considered two cases:

ション ふゆ アメビア メロア しょうくしゃ

N is prime.

N is composite.

We never considered N is a unit.

If p_1, \ldots, p_n are **any** set of rationals then $N = p_1 p_2 \cdots p_n + 1$ is a **a unit**.

Note that in the proof we considered two cases:

N is prime.

N is composite.

We never considered N is a unit.

Upshot The proof that \mathbb{Z} has an infinite number of primes uses that, for all $p_1 \cdots p_n + 1$ is never a unit.

ション ふゆ アメビア メロア しょうくしゃ

$$\mathbb{Q}_2 = \{ \tfrac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

$$\mathbb{Q}_2 = \{ rac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

 \mathbb{Q}_2 has 0.

▲□▶▲□▶▲□▶▲□▶ ■ りへぐ

$$\mathbb{Q}_2 = \{ rac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

 \mathbb{Q}_2 has 0.
 \mathbb{Q}_2 has units: all $rac{a}{b}$ where $a \equiv 1 \pmod{2}$.

$$\mathbb{Q}_2 = \{ \frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

$$\mathbb{Q}_2 \text{ has 0.}$$

$$\mathbb{Q}_2 \text{ has units: all } \frac{a}{b} \text{ where } a \equiv 1 \pmod{2}.$$

$$\mathbb{Q}_2 \text{ has primes: } 2, \frac{2}{3}, \frac{2}{5}, \frac{2}{7}, \frac{2}{9}, \dots$$

▲□▶▲□▶▲□▶▲□▶ ■ りへぐ

$$\begin{aligned} \mathbb{Q}_2 &= \{ \frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}. \\ \mathbb{Q}_2 \text{ has 0.} \\ \mathbb{Q}_2 \text{ has units: all } \frac{a}{b} \text{ where } a \equiv 1 \pmod{2}. \\ \mathbb{Q}_2 \text{ has primes: } 2, \frac{2}{3}, \frac{2}{5}, \frac{2}{7}, \frac{2}{9}, \dots \\ \text{Are there any more primes?} \end{aligned}$$

▲□▶▲圖▶▲圖▶▲圖▶ 圖 のへで

$$\begin{split} \mathbb{Q}_2 &= \{ \frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}. \\ \mathbb{Q}_2 \text{ has 0.} \\ \mathbb{Q}_2 \text{ has units: all } \frac{a}{b} \text{ where } a \equiv 1 \pmod{2}. \\ \mathbb{Q}_2 \text{ has primes: } 2, \frac{2}{3}, \frac{2}{5}, \frac{2}{7}, \frac{2}{9}, \dots \\ \text{Are there any more primes? No. I leave that for you to prove} \end{split}$$

$$\begin{split} \mathbb{Q}_2 &= \{ \frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}. \\ \mathbb{Q}_2 \text{ has 0.} \\ \mathbb{Q}_2 \text{ has units: all } \frac{a}{b} \text{ where } a \equiv 1 \pmod{2}. \\ \mathbb{Q}_2 \text{ has primes: } 2, \frac{2}{3}, \frac{2}{5}, \frac{2}{7}, \frac{2}{9}, \dots \\ \text{Are there any more primes? No. I leave that for you to prove.} \\ \text{So it looks like } \mathbb{Q}_2 \text{ has an infinite number of primes.} \end{split}$$

▲□▶▲□▶▲□▶▲□▶ ■ りへぐ

$$\begin{aligned} \mathbb{Q}_2 &= \{\frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}. \\ \mathbb{Q}_2 \text{ has 0.} \\ \mathbb{Q}_2 \text{ has units: all } \frac{a}{b} \text{ where } a \equiv 1 \pmod{2}. \\ \mathbb{Q}_2 \text{ has primes: } 2, \frac{2}{3}, \frac{2}{5}, \frac{2}{7}, \frac{2}{9}, \dots \end{aligned}$$
Are there any more primes? No. I leave that for you to prove.
So it looks like \mathbb{Q}_2 has an infinite number of primes.
BUT all of the primes listed are equivalent. So \mathbb{Q}_2 has only one prime.

◆□▶ ◆□▶ ◆三▶ ◆三▶ ◆□▶

$$\mathbb{Q}_2 = \{ \frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

$$\mathbb{Q}_2 \text{ has } 0.$$

 \mathbb{Q}_2 has units: all $\frac{a}{b}$ where $a \equiv 1 \pmod{2}$. \mathbb{Q}_2 has primes: $2, \frac{2}{3}, \frac{2}{5}, \frac{2}{7}, \frac{2}{9}, \ldots$

Are there any more primes? No. I leave that for you to prove.

So it looks like \mathbb{Q}_2 has an infinite number of primes.

BUT all of the primes listed are equivalent. So \mathbb{Q}_2 has only one prime.

So where does the proof that the primes are infinite go wrong? Discuss

ション ふゆ アメビア メロア しょうくしゃ

$$\mathbb{Q}_2 = \{ rac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

 \mathbb{Q}_2 has 0.

 \mathbb{Q}_2 has units: all $\frac{a}{b}$ where $a \equiv 1 \pmod{2}$.

 \mathbb{Q}_2 has primes: $2,\frac{2}{3},\frac{2}{5},\frac{2}{7},\frac{2}{9},\ldots$

Are there any more primes? No. I leave that for you to prove.

So it looks like \mathbb{Q}_2 has an infinite number of primes.

BUT all of the primes listed are equivalent. So \mathbb{Q}_2 has only one prime.

So where does the proof that the primes are infinite go wrong? Discuss

See next slide.

We actually have a list of primes: {2}. N = 2 + 1 = 3 which is a unit. So similar to why the proof fails for \mathbb{Q} .

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

 $\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$

・ロト・日下・日下・日、 日、 りへぐ

AI has a Finite Number of Primes

$$\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$$

The units are $\mathbb{U} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff 1 and constant coeff is 1}) | [p(a) = 0] \}.$

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >
$$\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$$

The units are

 $\mathbb{U} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff 1 and constant coeff is 1}) \\ [p(a) = 0]\}.$ We won't prove or need this.

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

$$\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$$

The units are

 $\mathbb{U} = \{a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff 1 and constant coeff is 1}) \\ [p(a) = 0]\}.$

We won't prove or need this. Units are not the problem this time.

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

$$\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$$

The units are

 $\mathbb{U} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff 1 and constant coeff is 1}) \\ [p(a) = 0] \}.$

We won't prove or need this. Units are not the problem this time.

Give me a number in \mathbb{AI} thats a prime. Discuss.

$$\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$$

The units are

 $\mathbb{U} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff 1 and constant coeff is 1}) \\ [p(a) = 0] \}.$

We won't prove or need this. Units are not the problem this time.

Give me a number in \mathbb{AI} thats a prime. Discuss.

There are no primes. See next slide.

Let $p \in \mathbb{AI}$.



Let $p \in \mathbb{AI}$. We show that p is not prime.

▲□▶ ▲□▶ ▲目▶ ▲目▶ 目 のへの

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$.

・ロト・日本・ヨト・ヨト・日・ つへぐ

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$.

*ロト *昼 * * ミ * ミ * ミ * のへぐ

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ - 三 - のへぐ

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0. $f(x) = x^n + a_{n-1}x^{n-1} + \cdots + a_1x + a_0$.

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0. $f(x) = x^n + a_{n-1}x^{n-1} + \cdots + a_1x + a_0$. $f(p) = p^n + a_{n-1}p^{n-1} + \cdots + a_1p + a_0 = 0$.

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0. $f(x) = x^n + a_{n-1}x^{n-1} + \dots + a_1x + a_0$. $f(p) = p^n + a_{n-1}p^{n-1} + \dots + a_1p + a_0 = 0$. Let $g(x) = x^{2n} + a_{n-1}x^{2(n-1)} + \dots + a_1x^2 + a_0$.

ション ふぼう メリン メリン しょうくしゃ

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0. $f(x) = x^n + a_{n-1}x^{n-1} + \dots + a_1x + a_0$. $f(p) = p^n + a_{n-1}p^{n-1} + \dots + a_1p + a_0 = 0$. Let $g(x) = x^{2n} + a_{n-1}x^{2(n-1)} + \dots + a_1x^2 + a_0$. $g(p^{1/2}) = p^n + a_{n-1}p^{n-1} + a_1p + a_0 = 0$.

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0. $f(x) = x^n + a_{n-1}x^{n-1} + \dots + a_1x + a_0$. $f(p) = p^n + a_{n-1}p^{n-1} + \dots + a_1p + a_0 = 0$. Let $g(x) = x^{2n} + a_{n-1}x^{2(n-1)} + \dots + a_1x^2 + a_0$. $g(p^{1/2}) = p^n + a_{n-1}p^{n-1} + a_1p + a_0 = 0$. So $\sqrt{p} \in \mathbb{AI}$.

Let $p \in \mathbb{AI}$. We show that p is not prime. Note that $p = \sqrt{p} \times \sqrt{p}$. We need to show that $\sqrt{p} \in \mathbb{AI}$. Let f be poly with lead coeff 1 such that f(p) = 0. $f(x) = x^{n} + a_{n-1}x^{n-1} + \dots + a_{1}x + a_{0}$ $f(p) = p^{n} + a_{n-1}p^{n-1} + \dots + a_{1}p + a_{0} = 0.$ l et $g(x) = x^{2n} + a_{n-1}x^{2(n-1)} + \dots + a_1x^2 + a_n$ $g(p^{1/2}) = p^n + a_{n-1}p^{n-1} + a_1p + a_0 = 0.$ So $\sqrt{p} \in \mathbb{AI}$. So there are no primes.

Lets revisit the proof.



Lets revisit the proof. Assume \mathbb{AI} has only a finite number of primes. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{AI} .

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Lets revisit the proof. Assume AI has only a finite number of primes. Let $\{p_1, \ldots, p_n\}$ be all of the primes in AI. Form $N = p_1 \cdots p_n + 1$

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

Lets revisit the proof. Assume AI has only a finite number of primes. Let $\{p_1, \ldots, p_n\}$ be all of the primes in AI. Form $N = p_1 \cdots p_n + 1$

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

Lets revisit the proof.

Assume \mathbb{AI} has only a finite number of primes. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{AI} .

Form $N = p_1 \cdots p_n + 1$

1. *N* prime. done since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.

ション ふぼう メリン メリン しょうくしゃ

Lets revisit the proof.

Assume \mathbb{AI} has only a finite number of primes. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{AI} .

Form $N = p_1 \cdots p_n + 1$

- 1. *N* prime. done since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.
- N is composite. N = ab where a, b ∉ U. If a and b are composite then break them down further until you get prime p, p divides N. So N = Mp.

Lets revisit the proof.

Assume \mathbb{AI} has only a finite number of primes. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{AI} .

Form $N = p_1 \cdots p_n + 1$

- 1. *N* prime. done since, for all $1 \le i \le n$, $p_i < N$ so $p_i \ne N$. *N* is a prime but not in $\{p_1, \ldots, p_n\}$. Contradiction.
- N is composite. N = ab where a, b ∉ U. If a and b are composite then break them down further until you get prime p, p divides N. So N = Mp.

This is where the proof breaks down! In AI you can keep going down and never get to a prime.

Example

2

◆□▶ ◆□▶ ◆三▶ ◆三▶ ・三 ・ のへで

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Example

 $\begin{array}{l} 2 \\ = 2^{1/2} \times 2^{1/2} \end{array}$

Example 2 = $2^{1/2} \times 2^{1/2}$ = $2^{1/4} \times 2^{1/4} \times 2^{1/4} \times 2^{1/4}$ = $2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8}$

Example 2 = $2^{1/2} \times 2^{1/2}$ = $2^{1/4} \times 2^{1/4} \times 2^{1/4} \times 2^{1/4}$ = $2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8}$ etc.

Example 2 $= 2^{1/2} \times 2^{1/2}$ $= 2^{1/4} \times 2^{1/4} \times 2^{1/4} \times 2^{1/4}$ $= 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8}$ etc.

So what property of \mathbb{Z} was used to avoid this problem?

Example 2 $= 2^{1/2} \times 2^{1/2}$ $= 2^{1/4} \times 2^{1/4} \times 2^{1/4} \times 2^{1/4}$ $= 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8} \times 2^{1/8}$ etc.

So what property of $\ensuremath{\mathbb{Z}}$ was used to avoid this problem? See next slide.

Def An **Atomic Integral Domain** is an integral domain such that every element of $\mathbb{D} - (\mathbb{U} \cup \{0\})$ can be written (not necessarily uniquely) as $up_1^{x_1} \cdots p_m^{x_m}$ where u is a unit and all of the p_i 's are irreducible.

Def An **Atomic Integral Domain** is an integral domain such that every element of $\mathbb{D} - (\mathbb{U} \cup \{0\})$ can be written (not necessarily uniquely) as $up_1^{x_1} \cdots p_m^{x_m}$ where u is a unit and all of the p_i 's are irreducible.

ション ふぼう メリン メリン しょうくしゃ

Examples

Z. Key is that f(x) = |x| is such that f(a) < f(ab) and f(ab) < f(b). So when you factor you end up with smaller numbers.

Def An **Atomic Integral Domain** is an integral domain such that every element of $\mathbb{D} - (\mathbb{U} \cup \{0\})$ can be written (not necessarily uniquely) as $up_1^{x_1} \cdots p_m^{x_m}$ where u is a unit and all of the p_i 's are irreducible.

Examples

Z. Key is that f(x) = |x| is such that f(a) < f(ab) and f(ab) < f(b). So when you factor you end up with smaller numbers.

 $\mathbb{Z}[i]$. Key is that $f(x + iy) = x^2 + y^2$ decreases when you factor.

Def An **Atomic Integral Domain** is an integral domain such that every element of $\mathbb{D} - (\mathbb{U} \cup \{0\})$ can be written (not necessarily uniquely) as $up_1^{x_1} \cdots p_m^{x_m}$ where u is a unit and all of the p_i 's are irreducible.

Examples

Z. Key is that f(x) = |x| is such that f(a) < f(ab) and f(ab) < f(b). So when you factor you end up with smaller numbers.

 $\mathbb{Z}[i]$. Key is that $f(x + iy) = x^2 + y^2$ decreases when you factor. AI has no such function (called a Norm).

Def An **Atomic Integral Domain** is an integral domain such that every element of $\mathbb{D} - (\mathbb{U} \cup \{0\})$ can be written (not necessarily uniquely) as $up_1^{x_1} \cdots p_m^{x_m}$ where u is a unit and all of the p_i 's are irreducible.

Examples

Z. Key is that f(x) = |x| is such that f(a) < f(ab) and f(ab) < f(b). So when you factor you end up with smaller numbers.

 $\mathbb{Z}[i]$. Key is that $f(x + iy) = x^2 + y^2$ decreases when you factor. AI has no such function (called a Norm).

Upshot The proof that $\mathbb Z$ has an infinite number of primes used that $\mathbb Z$ is atomic.

The EG-Proof that Primes are Infinite and Where it Falls Apart

May 8, 2025

ション ふぼう メリン メリン しょうくしゃ

Where Does EG-Proof Fail for \mathbb{Q} ?

Thm The number of primes in \mathbb{Q} is infinite (attempt).

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Where Does EG-Proof Fail for Q?

Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L .
Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . We define a coloring on $N \subseteq \mathbb{Q}$ as follows.

Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . We define a coloring on $N \subseteq \mathbb{Q}$ as follows.

Let $\text{COL} \colon \mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

ション ふゆ アメビア メロア しょうくしゃ

Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . We define a coloring on $N \subseteq \mathbb{Q}$ as follows.

Let $\text{COL} \colon \mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L})=(a_1 \pmod{4},\ldots,a_L \pmod{4})$$

ション ふゆ アメビア メロア しょうくしゃ

Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . We define a coloring on $N \subseteq \mathbb{Q}$ as follows.

Let $\operatorname{COL}: \mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

Two Issues

1) Factoring elements of \mathbb{N} into primes in \mathbb{N} every number is of the form $p_1^{a_1} \cdots p_L^{a_L}$. No issue with units since the only units is 1. We are factoring elements of \mathbb{N} into primes in \mathbb{Q} so units may be needed.

Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . We define a coloring on $N \subseteq \mathbb{Q}$ as follows.

Let $\operatorname{COL}: \mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

Two Issues

1) Factoring elements of \mathbb{N} into primes in \mathbb{N} every number is of the form $p_1^{a_1} \cdots p_L^{a_L}$. No issue with units since the only units is 1. We are factoring elements of \mathbb{N} into primes in \mathbb{Q} so units may be needed.

2) In our proof we used mod 4. Lets keep it n for now and try to pick some n that will work.

Thm The number of primes in \mathbb{Q} is infinite (attempt). Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . We define a coloring on $N \subseteq \mathbb{Q}$ as follows.

Let $\operatorname{COL}: \mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

Two Issues

1) Factoring elements of \mathbb{N} into primes in \mathbb{N} every number is of the form $p_1^{a_1} \cdots p_L^{a_L}$. No issue with units since the only units is 1. We are factoring elements of \mathbb{N} into primes in \mathbb{Q} so units may be needed.

2) In our proof we used mod 4. Lets keep it n for now and try to pick some n that will work.

We define the coloring as follows:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

Let COL: $\mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

 $\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L})=(a_1\pmod{n},\ldots,a_L\pmod{n})$

Let $\text{COL} \colon \mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

 $\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \ldots, a_L \pmod{n})$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

By Schur's Thm there exists x, y, z same color with x + y = z.

Let $\text{COL} \colon \mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

 $\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

Let $\text{COL}: \mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L})=(a_1 \pmod{n},\ldots,a_L \pmod{n})$$

ション ふぼう メリン メリン しょうくしゃ

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

 $\begin{aligned} x &= u_x p_1^{nx_1 + e_1} \cdots p_L^{nx_L + e_L} \\ y &= u_y p_1^{ny_1 + e_1} \cdots p_L^{ny_L + e_L} \\ z &= u_z p_1^{nz_1 + e_1} \cdots p_L^{nz_n + e_L} \end{aligned}$

Let $\text{COL}: \mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

ション ふぼう メリン メリン しょうくしゃ

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

 $x = u_x p_1^{nx_1+e_1} \cdots p_L^{nx_L+e_L}$ $y = u_y p_1^{ny_1+e_1} \cdots p_L^{ny_L+e_L}$ $z = u_z p_1^{nz_1+e_1} \cdots p_L^{nz_n+e_L}$ x + y = z

Let COL: $\mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= u_x p_1^{nx_1+e_1} \cdots p_L^{nx_L+e_L} \\ y &= u_y p_1^{ny_1+e_1} \cdots p_L^{ny_L+e_L} \\ z &= u_z p_1^{nz_1+e_1} \cdots p_L^{nz_n+e_L} \\ x + y &= z \\ u_x p_1^{nx_1+e_1} \cdots p_L^{nx_L+e_L} + p_1^{ny_1+e_1} \cdots p_L^{ny_L+e_L} = u_x p_1^{nz_1+e_1} \cdots p_L^{nz_n+e_L} \end{aligned}$$

・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・
・

Let COL: $\mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= u_{x} p_{1}^{nx_{1}+e_{1}} \cdots p_{L}^{nx_{L}+e_{L}} \\ y &= u_{y} p_{1}^{ny_{1}+e_{1}} \cdots p_{L}^{ny_{L}+e_{L}} \\ z &= u_{z} p_{1}^{nz_{1}+e_{1}} \cdots p_{L}^{nz_{n}+e_{L}} \\ x + y &= z \\ u_{x} p_{1}^{nx_{1}+e_{1}} \cdots p_{L}^{nx_{L}+e_{L}} + p_{1}^{ny_{1}+e_{1}} \cdots p_{L}^{ny_{L}+e_{L}} = u_{x} p_{1}^{nz_{1}+e_{1}} \cdots p_{L}^{nz_{n}+e_{L}} \\ u_{y} p_{1}^{nx_{1}} \cdots p_{L}^{nx_{L}} + p_{1}^{ny_{1}} \cdots p_{L}^{ny_{L}} = u_{y} p_{1}^{nz_{1}} \cdots p_{L}^{nz_{n}} \end{aligned}$$

Let COL: $\mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= u_{x} p_{1}^{nx_{1}+e_{1}} \cdots p_{L}^{nx_{L}+e_{L}} \\ y &= u_{y} p_{1}^{ny_{1}+e_{1}} \cdots p_{L}^{ny_{L}+e_{L}} \\ z &= u_{z} p_{1}^{nz_{1}+e_{1}} \cdots p_{L}^{nz_{n}+e_{L}} \\ x + y &= z \\ u_{x} p_{1}^{nx_{1}+e_{1}} \cdots p_{L}^{nx_{L}+e_{L}} + p_{1}^{ny_{1}+e_{1}} \cdots p_{L}^{ny_{L}+e_{L}} = u_{x} p_{1}^{nz_{1}+e_{1}} \cdots p_{L}^{nz_{n}+e_{L}} \\ u_{y} p_{1}^{nx_{1}} \cdots p_{L}^{nx_{L}} + p_{1}^{ny_{1}} \cdots p_{L}^{ny_{L}} = u_{y} p_{1}^{nz_{1}} \cdots p_{L}^{nz_{n}} \\ u_{x} (p_{1}^{x_{1}} \cdots p_{L}^{x_{L}})^{n} + u_{y} (p_{1}^{y_{1}} \cdots p_{L}^{y_{L}})^{n} = u_{z} (p_{1}^{z_{1}} \cdots p_{L}^{z_{L}})^{n} \end{aligned}$$

・ロト・西ト・ヨト・ヨト・ 日・ のへぐ

Let COL: $\mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

By Schur's Thm there exists x, y, z same color with x + y = z. Assume the color is (e_1, \ldots, e_L) .

$$\begin{aligned} x &= u_{x} p_{1}^{nx_{1}+e_{1}} \cdots p_{L}^{nx_{L}+e_{L}} \\ y &= u_{y} p_{1}^{ny_{1}+e_{1}} \cdots p_{L}^{ny_{L}+e_{L}} \\ z &= u_{z} p_{1}^{nz_{1}+e_{1}} \cdots p_{L}^{nz_{n}+e_{L}} \\ x + y &= z \\ u_{x} p_{1}^{nx_{1}+e_{1}} \cdots p_{L}^{nx_{L}+e_{L}} + p_{1}^{ny_{1}+e_{1}} \cdots p_{L}^{ny_{L}+e_{L}} = u_{x} p_{1}^{nz_{1}+e_{1}} \cdots p_{L}^{nz_{n}+e_{L}} \\ u_{y} p_{1}^{nx_{1}} \cdots p_{L}^{nx_{L}} + p_{1}^{ny_{1}} \cdots p_{L}^{ny_{L}} = u_{y} p_{1}^{nz_{1}} \cdots p_{L}^{nz_{n}} \\ u_{x} (p_{1}^{x_{1}} \cdots p_{L}^{x_{L}})^{n} + u_{y} (p_{1}^{y_{1}} \cdots p_{L}^{y_{L}})^{n} = u_{z} (p_{1}^{z_{1}} \cdots p_{L}^{z_{L}})^{n} \\ u_{x} \chi^{n} + u_{y} Y^{n} = u_{z} Z^{n}. \end{aligned}$$

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

We need that there is an n such that

We need that there is an *n* such that for all $u_x, u_y, u_z \in U$ and $a, b, c \in \mathbb{Q}$

$$u_x X^n + u_y Y^n = u_z Z^n$$

has no solution.

We need that there is an *n* such that for all $u_x, u_y, u_z \in U$ and $a, b, c \in \mathbb{Q}$

$$u_x X^n + u_y Y^n = u_z Z^n$$

has no solution.

Not True Fix *n*. Let $u_x = u_y = \frac{1}{2}$, $u_z = 1$, X = Y = Z = 1.

$$u_x X^n + u_y Y^n = u_z Z^n$$

Becomes

$$\frac{1}{2}1^{n} + \frac{1}{2}1^{n} = 1 \times 1^{n}$$
$$\frac{1}{2} + \frac{1}{2} = 1$$

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

The EG-proof that there are an infinite number of primes (in $\mathbb N)$ did not transfer to $\mathbb Q$ because

▲□▶ ▲□▶ ▲目▶ ▲目▶ 二目 - のへで

Where Does EG-Proof Fail for Q?: Upshot

The EG-proof that there are an infinite number of primes (in \mathbb{N}) did not transfer to \mathbb{Q} because Its NOT that **FLT** is false over \mathbb{Q} . Indeed—FLT is true over \mathbb{Q} (follows form FLT being true over \mathbb{Z}).

Where Does EG-Proof Fail for Q?: Upshot

The EG-proof that there are an infinite number of primes (in \mathbb{N}) did not transfer to \mathbb{Q} because Its NOT that **FLT** is false over \mathbb{Q} . Indeed—FLT is true over \mathbb{Q} (follows form FLT being true over \mathbb{Z}).

Its because the following variant of FLT is false for \mathbb{Q} :

Where Does EG-Proof Fail for Q?: Upshot

The EG-proof that there are an infinite number of primes (in \mathbb{N}) did not transfer to \mathbb{Q} because Its NOT that **FLT** is false over \mathbb{Q} . Indeed—FLT is true over \mathbb{Q} (follows form FLT being true over \mathbb{Z}).

Its because the following **variant** of FLT is false for \mathbb{Q} : There exists $n \in \mathbb{N}$ such that the following has no solution:

$$u_x X^n + u_y Y^n = u_z Z^n$$

ション ふぼう メリン メリン しょうくしゃ

where $u_x, u_y, y_z \in \mathbb{U}$ and $X, Y, Z \in \mathbb{Q}$.

May 8, 2025

- イロト イボト イモト - モー のへぐ

1. Read the Gasarch paper. Note that its initial proof was a generalization of what was presented here.

▲□▶ ▲□▶ ▲目▶ ▲目▶ 二目 - のへで

- 1. Read the Gasarch paper. Note that its initial proof was a generalization of what was presented here.
- 2. Read in Gasarch's paper the **Sanity Check** which has more domains with a finite number of primes.

- 1. Read the Gasarch paper. Note that its initial proof was a generalization of what was presented here.
- 2. Read in Gasarch's paper the **Sanity Check** which has more domains with a finite number of primes.
- Read the other papers on the website of Ramsey-Primes paper. Some of the papers are difficult so try to just figure out the proof for Z or N, and then see where it fails for Q and Q₂. (I think they all fail for AI because AI is not atomic, though check that.)