# Van Der Warden's (VDW) Thm

#### **Exposition by William Gasarch**

July 19, 2024

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#### VDW's Thm

**Def** Let  $W, k, c \in \mathbb{N}$ . Let COL:  $[W] \rightarrow [c]$ . A mono k-**AP** is an arithmetic progression of length k where every elements has the same color. We often say

 $a, a + d, \ldots, a + (k - 1)d$  are all he same color

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W(2, c) = c + 1. By Pigeon Hole Principle.

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W(3,2) =Hmmm, this is the first non-trivial one.

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- 1. Within one block.
- 2. If we take enough blocks, how they relate.

**Def**: a, a + d, a + 2d is an **almost mono 3AP** if  $COL(a) = COL(a+d) \neq COL(a+2d)$ . The color of an almost **mono 3AP** is COL(a) = COL(a+d).

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Look at the first three elements of a block of 5:

- 1. RRR or BBB. 1-2-3 is mono 3AP.
- 2. RBR or BRB. 1-3-5 is mono 3AP or almost mono 3AP.
- 3. RBB or BRR. 2-3-4 is mono 3AP or almost mono 3AP.

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- 4. BBR or RRB. 1-2-3 is almost mono 3AP.
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So always get a mono 3AP or an almost mono 3AP. Can assume its almost mono 3AP and its R.

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We can get by with LESS blocks- we will consider this point after the proof.

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#### Let $COL \colon [W] \to [2]$ .

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Break [W] into 65 blocks of size 5 which we think of as being 32-colored.

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If ? is B then get B 3-AP.

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If ? is B then get B 3-AP. If ? is R then get R 3-AP.

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If ? is B then get B 3-AP. If ? is R then get R 3-AP. Done!

#### Side Note: Can Get By With Less Blocks

Warning This Slide is NOT important.



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So we don't really have to look at 32 colorings.

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How many colorings of a block already have a mono 3AP.

#### Side Note: Can Get By With Less Blocks (cont)

```
RRRXY with X, Y \in \{R, B\}. 4 colorings.

BBBXY with X, Y \in \{R, B\}. 4 colorings.

RBRRR

RBRBR

BRBBB

BRBBB

RBBBX with X \in \{R, B\}. 2 colorings.

BRRRX with X \in \{R, B\}. 2 colorings.

RRBBB

BBRRR
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There are 16 blocks which already have a mono 3AP. Hence can use 32 - 16 = 16 blocks.

### Side Note: Can Get By With Less Blocks (cont)

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I really do not care.

Is W(3,2) = 365?

No What is W(3,2)?



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One can work out by hand that

W(3,2) = 9.

We will later say which VDW numbers are know and how they compare to the bounds given by the proof of VDW's Thm.

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No What is W(3,2)?

One can work out by hand that

W(3,2)=9.

We will later say which VDW numbers are know and how they compare to the bounds given by the proof of VDW's Thm.

**Spoiler Alert** The few known VDW numbers are **much smaller** than the bounds given by the proof of VDW's Thm.

#### $\mathrm{COL}\colon [W]\to [3].$

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COL:  $[W] \rightarrow [3]$ . How big should the blocks be?



COL:  $[W] \rightarrow [3]$ . How big should the blocks be? 7.



 $\begin{aligned} & \mathrm{COL}\colon [\mathcal{W}] \to [3]. \\ & \text{How big should the blocks be? 7.} \\ & \text{Then } \forall \text{ 3-coloring of block } \exists \text{ mono 3AP or almost mono 3AP.} \end{aligned}$ 

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Darn. Now what? Discuss

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Darn. Now what? Discuss We have 2 almost mono 3APs of diff colors that same last element.

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Let *W* be LOTS of blocks of size  $7 \times 2 \times (3^7 + 1)$ .

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# $W(\mathbf{3}, c)$

From what you have seen:

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- The Hales-Jewitt Thm is a general theorem from which VDW is a corollary. We won't be doing that.

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$$\begin{split} & \mathcal{W}(2,c) = c+1 \text{ is just PHP.} \\ & \mathcal{W}(2,2^5) \implies \mathcal{W}(3,2) \\ & \mathcal{W}(2,3^{2\times 3^7}+1) \implies \mathcal{W}(3,3). \\ & \mathcal{W}(2,X) \implies \mathcal{W}(3,4) \text{ where } X \text{ is very large.} \end{split}$$

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Note that we **do not** do  $W(3,2) \implies W(3,3).$ 



#### $\mathrm{COL}\colon [W]\to [4].$

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COL:  $[W] \rightarrow [4]$ . Key Take blocks of size 2W(3, 2).

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 $\begin{array}{l} {\rm COL}\colon [W] \to [4]. \\ \\ \hbox{Key Take blocks of size $2W(3,2)$.} \\ \\ {\rm Within a block is mono $4AP$ or almost mono $4AP$.} \end{array}$ 

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If ? is B get mono 4AP. If ? is R get mono 4AP. Done!

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in which case just know that the proof given gives bounds that are NOT prim rec.)