There Is a 2-Coloring Of the Plane Without a mono Red 3-Stick or a mono Blue Big-Stick

Exposition by William Gasarch-U of MD

Credit Where Credit is Due

The main result in these slides is due to Conlon and Wu (2022).

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Main Theorem

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Thm There exists $\operatorname{COL} \colon \mathbb{R}^n \to [2]$ such that there is no a $\mathbb{R} \ \ell_3$, and there is no B ℓ_m where m will be determined later. m will be around 10^{50} . The proof for \mathbb{R}^n and \mathbb{R}^2 are identical.

Thm There exists ${\rm COL}\colon \mathbb{R}^n \to [2]$ such that there is no a $\mathbb{R}\ \ell_3$, and there is no B ℓ_m where m will be determined later. m will be around 10^{50} . The proof for \mathbb{R}^n and \mathbb{R}^2 are identical. Open Find an easier proof of \mathbb{R}^2 case.

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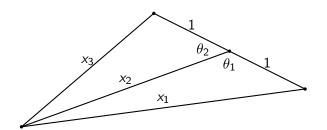
Let x_1 = d(\vec{0}, \vec{a}_1), x_2 = d(\vec{0}, \vec{a}_2), x_3 = d(\vec{0}, \vec{a}_3)
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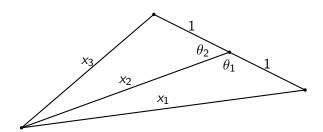
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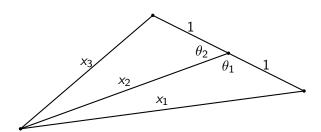
Let x_1 = d(\vec{0}, \vec{a}_1), x_2 = d(\vec{0}, \vec{a}_2), x_3 = d(\vec{0}, \vec{a}_3)

And we know 1 = d(\vec{a}_1, \vec{a}_2), 1 = d(\vec{a}_2, \vec{a}_3),
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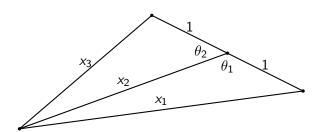


Bottom Triangle:



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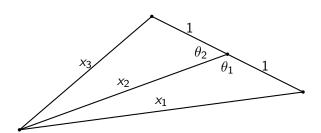
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Top Triangle:

Law of cosines: $x_3^2 = x_2^2 + 1 - 2x_2 \cos(\theta_2)$.

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. Hence $\cos(\theta_2) = -\cos(\theta_1)$.

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$$heta_2=\pi- heta_2$$
. Hence $\cos(heta_2)=-\cos(heta_1)$. Law of cosines: $x_1^2=x_2^2+1-2x_2\cos(heta_1)$. Law of cosines: $x_3^2=x_2^2+1-2x_2\cos(heta_2)=x_2^2+1+2x_2\cos(heta_1)$. Add to get
$$x_1^2+x_2^2=2x_2^2+2.$$

First Plan On How to Avoid R ℓ_3

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Upshot on R ℓ_3

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We will also have a condition on COL' that will make $\mathrm{COL}(\vec{a}) = \mathrm{COL}'(d(\vec{0},\vec{a}))$ not have any $\mathbf{B} \ \ell_m$

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For all $1 \le i \le m-1$ we know $1 = d(\vec{a_i}, \vec{a_{i+1}})$.

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$$x_{i-1}^2 + x_{i+1}^2 = 2x_i^2 + 2.$$

Real Plan

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We will define COL'': $\mathbb{Z}_q \to [2]$.

We will then define $COL' \colon \mathbb{R} \to [2]$ by

$$COL'(y) = COL''(\lfloor y \rfloor \pmod{q}).$$

An Example of A Coloring with q = 5

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COL''(0) = R

COL''(1) = B

COL''(2) = B

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-5
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The next slide recaps where we are and says why COL'' helps us.

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Let COL: $\mathbb{R}^2 \to [2]$ be $\mathrm{COL}(\vec{a}) = \mathrm{COL}'(d(0, \vec{a}))$. Did show 1) COL has no \mathbb{R} ℓ_3 (in \mathbb{R}^2).

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- 2) Has no ${\bf B}$ solution (in ${\mathbb Z}$) to

For all
$$2 \le i \le m-1$$
, $y_{i-1} + y_{i+1} = 2y_i + 2$

Let COL: $\mathbb{R}^2 \to [2]$ be $COL(\vec{a}) = COL'(d(0, \vec{a}))$. Did show

- 1) COL has no $\mathbb{R} \ell_3$ (in \mathbb{R}^2).
- 2) COL has no $\mathbf{B} \ell_m$ (in \mathbb{R}^2).

We Define COL"

TO define $\mathrm{COL}^{\prime\prime}$ we'll need some hard math. Or will we?

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BILL: We need to find a coloring. This requires hard math.

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The End



Pick a Coloring Randomly

We will pick $\mathrm{COL} \colon \mathbb{Z}_q \to [2]$ randomly.

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Lemmas and a Theorem of Independent Interest

What does $p(x) = x^2 + \pi x + e \pmod{13}$ mean?

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Lets say p(10) = 134.1325 (thats not true but its a good approx). 134.1324 (mod 13) = 4.1324.

So it makes sense to consider $p(x) \pmod{q}$ where $p(x) \in \mathbb{R}[x]$.

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Lemma Let $p(x) = x^2 + \alpha x + \beta$ where $\alpha, \beta \in \mathbb{R}$.

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Then

X hits at least q/6 of the intervals [0,1), [1,2), ..., [q-1,q).



Consider $\alpha \pmod{q}$, $2\alpha \pmod{q}$, ..., $q^2\alpha \pmod{q}$.

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So there exists i,j such that $|i\alpha \pmod q - j\alpha \pmod q| \le \frac1q$.

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We will consider two cases:

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We will consider two cases:

Case 1 $k \not\equiv 0 \pmod{q}$.

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Map each one to which interval [0,1), ..., [q-1,q) that it is in.

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We will consider two cases:

Case 1 $k \not\equiv 0 \pmod{q}$.

Case 2 $k \equiv 0 \pmod{q}$.

Recap There is a $k \not\equiv 0 \pmod{q}$ such that $|k\alpha \mod q| \leq \frac{1}{q}$.

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1) Show $x^2 + \beta \pmod{q}$ hits $\geq (q+1)/2$ intervals.

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Plan

- 1) Show $x^2 + \beta \pmod{q}$ hits $\geq (q+1)/2$ intervals.
- 2) Show that adding αx has a small effect since $|k\alpha \pmod{q}| \leq \frac{1}{a}$.

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We consider several sets and see how many intervals they hit.

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$$\begin{split} &\mathrm{SQ}_q = \{1^2 \pmod q, 2^2 \pmod q, \ldots, q^2 \pmod q\}\}. \\ &q \text{ is a prime so squaring is 2-to-1}. \ \ \mathsf{Hence} \ |\mathrm{SQ}_q| = (q+1)/2. \end{split}$$

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 $\mathrm{SQ}_q=\{1^2\pmod q,2^2\pmod q,\ldots,q^2\pmod q\}.$ q is a prime so squaring is 2-to-1. Hence $|\mathrm{SQ}_q|=(q+1)/2.$ Since every element in SQ_q is an integer, hits (q+1)/2 intervals.

We consider $f_1(x) = x^2 + \beta \pmod{q}$.

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Since X is the squares all shifted by β , $|X_1| = (q+1)/2$.

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Since $k \not\equiv 0 \pmod{q}, \{k, 2k, \dots, qk\} = \{1, 2, \dots, q\}$. Hence

Since $k \neq 0 \pmod{q}$, $\{k, 2k, \ldots, qk\} = \{1, 2, \ldots, q\}$. Hence X = Y.

Why
$$m = q^3$$
?

We have shown that

$$\{f_1(k), f_1(2k), \ldots, f_1(qk)\}.$$

hits (q+1)/2 intervals. Note that $qk \le q^3 = m$. This is why we needed $m = q^3$ in the hypothesis.

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We need to show that $Z = \{f(1), f(2), \dots, f(q^3)\}$ hits $\geq q/6$ intervals.

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We will do this on the next slide.

 $\{f_1(k), f_1(2k), \dots, f_1(qk)\}$ hits (q+1)/2 intervals.

```
\{f_1(k), f_1(2k), \ldots, f_1(qk)\} hits (q+1)/2 intervals. We show that \{f(1), \ldots, f(q^3)\} hits \geq q/6 intervals by just looking at the subset \{f(k), f(2k), \ldots, f(qk)\}.
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```

```
 \{f_1(k), f_1(2k), \ldots, f_1(qk)\} \text{ hits } (q+1)/2 \text{ intervals.}  We show that \{f(1), \ldots, f(q^3)\} hits \geq q/6 intervals by just looking at the subset \{f(k), f(2k), \ldots, f(qk)\}.  \{f(k), f(2k), \ldots, f(qk)\}:  f(k) = f_1(k) + k\alpha. Key Recall k\alpha \pmod q | \leq \frac1q \leq 1.  f(2k) = f_1(2k) + 2k\alpha. Key Recall 2k\alpha \pmod q | \leq \frac2q \leq 1.  \vdots   f(qk) = f_1(2k) + qk\alpha. Key Recall qk\alpha \pmod q | \leq \frac q \leq 1.
```

Recap The set $Y = \{f_1(k), \dots, f_1(qk)\}$ hits (q+1)/2 intervals of length 1.

 $Z = \{f(k), \dots, f(qk)\}$ can be viewed as taking every element in Y and adding or subtracting ≤ 1 to it. It is easy to show that Z hits $\geq q/6$ intervals.

Case 2: $k \equiv 0 \pmod{q}$

OMITTED FOR NOW.

Another Lemma Of Independent Interest

The Sign Function and Other Notation

Def if $a \in \mathbb{R}$ then

The Sign Function and Other Notation

Def if $a \in \mathbb{R}$ then

$$sign(a) = \begin{cases} -1 & \text{if } a < 0 \\ 0 & \text{if } a = 0 \\ 1 & \text{if } a > 0 \end{cases}$$
 (1)

The Sign Function and Other Notation

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 (1)

Notation If $\eta \in \{-1,0,1\}^*$ then $\eta(i)$ is the *i*th character in η .

$$p_1(x, y) = x + 2y - 3$$
 $p_2(x, y) = -2x + 3y - 7$
 $p_3(x, y) = 4x - y$

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 $p_2(x,y) = -2x + 3y - 7$
 $p_3(x,y) = 4x - y$
We care about $(sign(p_1(x,y)), sign(p_2(x,y)), sign(p_3(x,y)))$.

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We care about $(sign(p_1(x,y)), sign(p_2(x,y)), sign(p_3(x,y)))$.

(x,y)	$(p_1(x,y), p_2(x,y), p_3(x,y))$	sign pattern
(0,0)	(-3, -7, 0)	(-, -, 0)
(10,0)	(7, -27, 40)	(+,-,+)
(0, 10)	(17, 23, -10)	(+,+,-)
(1,1)	(0, -6, 3)	(0, -, +)
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There are $3^3 = 27$ sign patterns. (p_1, p_2, p_3) has at least 5. I doubt it has anywhere near 27.

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Note Obvious bound on number of sign patterns: 3^M **Question** Is there a better bound? Yes!

Lemma Let $M \in \mathbb{N}$. Let $p_1, \ldots, p_M \in \mathbb{Z}[x, y]$.

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Lemma also a corollary of a theorem by Olenik-Petrovsky-Thom-Milnor.

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- 2) Assume the reals are 2-colored. Then a **mono image-solution** of F is an image-solution $(a, d; y_1, \ldots, y_m)$ where y_1, \ldots, y_m are all the same color. Note that a, d need not be that color.

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I am skipping some stuff of interest to get to some stuff that is of more interest.

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prob there is a **blue image-solution** to p_1, \ldots, p_m is $< 2500 m^6 b^{m/6}$.

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Seems like there are $(2m^2)^m$ possibilities.



Since the coloring is mod q we can assume $a, d \in [0, q)$.

By premise, $0 \le p_i(a, d) \le 2m^2$.

We now ask about the intervals (not mod q).

1) Of the ints [0,1), [1,2), ..., $[2m^2-1,2m^2)$ which one has $p_1(a,d)$? There are $2m^2$ possibilities.

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Number of ways $p_1(a, d), \dots, p_m(a, d)$ are in the ints is $\leq 2500 m^6$.



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Hence there exists a 2-coloring with no $\mathbb{R}\ell_3$ or $\mathbb{B}\ell_m$ for large enough m.

