Schur's Thm + FLT implies Primes Infinite

May 10, 2025

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- 2. All of these proofs have other points to make after they prove primes ∞ .

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- 3. Look at domains where the number of primes is finite and see where the standard proof fails, and where the EG-proof fails.

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Background Needed For EG-Proof

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- 1. $[n] = \{1, 2, \ldots, n\}.$
- 2. $\binom{A}{k}$ is the set of all subsets of A of size k.

Thm $(\forall c)(\exists S = S(c))$ st for all *c*-colorings COL: $[S] \rightarrow [c]$ there exists *x*, *y*, *z* monochromatic such that x + y = z.

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Let x = c - b, y = b - a, z = c - a. So let S(c) = R(3; c) (homog set 3, colors c).

Fermat's Last Theorem

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To divide a cube into two cubes, a fourth power, or in general any power whatever above the second into two powers of the same denomination, is impossible, and I have assuredly found a proof of this, but the margin is too narrow to contain it.

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$$(\forall n \geq 3)(\forall x, y, z \in \mathbb{N} - \{0\})[x^n + y^n \neq z^n].$$

This has come to be known as Fermat's Last Theorem.

Did Fermat Have a Proof? Arguments Against

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2) Andrew Wiles and Richard Taylor proved FLT in the early 1990s with techniques far beyond what Fermat could have known.

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 - 2.2 meant to say that Fermat died in a duel in a dual timeline.

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To the margin add 200 pages.

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$$\begin{aligned} x + y &= z \\ p_1^{4x_1 + e_1} \cdots p_L^{4x_L + e_L} + p_1^{4y_1 + e_1} \cdots p_L^{4y_L + e_L} &= p_1^{4z_1 + e_1} \cdots p_L^{4z_n + e_L} \\ p_1^{4x_1} \cdots p_L^{4x_L} + p_1^{4y_1} \cdots p_L^{4y_L} &= p_1^{4z_1} \cdots p_L^{4z_n} \end{aligned}$$

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Thm The number of primes is infinite. Assume, BWOC, that the primes are finite. p_1, \ldots, p_L . Let COL: $\mathbb{N} \to \{0, 1, 2, 3\}^L$ be the following coloring:

$$\operatorname{COL}(p_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{4}, \dots, a_L \pmod{4})$$

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How to Ask the Question of Primes Infinite

May 10, 2025

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Def An **Integral Domain** is a set *D* together with operations +, \times such that the following hold

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1) $\mathbb{Z}_{12}=\{0,\ldots,11\}$ with mod 12 math.

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Types of Elts in an ID 0, units, irreducibles, composites.

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Equivalence Classes of Irreducibles

Convention Let $\mathbb D$ be an Int Dom with Units $\mathbb U,$ Irreds $\mathbb I.$

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On theses slides infinite will mean infinite up to units.

The Normal Proof that Primes are Infinite and Where it Falls Apart

May 10, 2025

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Thm The set of primes in \mathbb{Z} is infinite. Assume not. Let $\{p_1, \ldots, p_n\}$ be all of the primes in \mathbb{Z} . (Note- if p and -p both appear, we just take p.)

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 p ∉ {p₁,..., p_n} since if it was then 0 ≡ 1 (mod p).

 ${\mathbb Q}$ has 0, units, NO primes, NO composites.



 $\mathbb Q$ has 0, units, NO primes, NO composites. Where does proof primes ∞ go wrong? Discuss

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If p_1, \ldots, p_n are **any** set of rationals then $N = p_1 p_2 \cdots p_n + 1$ is a **a unit**.

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Upshot The proof that \mathbb{Z} has an infinite number of primes uses that, for all $p_1 \cdots p_n + 1$ is never a unit.

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$$\mathbb{Q}_2 = \{ \frac{a}{b} : \gcd(a, b) = 1 \land b \equiv 1 \pmod{2} \}.$$

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See next slide.

We actually have a list of primes: {2}. N = 2 + 1 = 3 which is a unit. So similar to why the proof fails for \mathbb{Q} .

 $\mathbb{AI} = \{ a \in \mathbb{C} : (\exists f(x) \in \mathbb{Z}[x] \text{ lead coeff } 1))[p(a) = 0] \}.$

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There are no primes. See next slide.

Let $p \in \mathbb{AI}$.



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Lets revisit the proof.



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Form $N = p_1 \cdots p_n + 1$

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This is where the proof breaks down! In AI you can keep going down and never get to a prime.

Example

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 $\begin{array}{l} 2 \\ = 2^{1/2} \times 2^{1/2} \end{array}$

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So what property of \mathbb{Z} was used to avoid this problem?

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So what property of $\ensuremath{\mathbb{Z}}$ was used to avoid this problem? See next slide.

Def An **Atomic Integral Domain** is an integral domain such that every element of $\mathbb{D} - (\mathbb{U} \cup \{0\})$ can be written (not necessarily uniquely) as $up_1^{x_1} \cdots p_m^{x_m}$ where u is a unit and all of the p_i 's are irreducible.

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Upshot The proof that $\mathbb Z$ has an infinite number of primes used that $\mathbb Z$ is atomic.

The EG-Proof that Primes are Infinite and Where it Falls Apart

May 10, 2025

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Where Does EG-Proof Fail for \mathbb{Q} ?

Thm The number of primes in \mathbb{Q} is infinite (attempt).

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Where Does EG-Proof Fail for Q?

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 $\begin{aligned} x &= u_x p_1^{nx_1 + e_1} \cdots p_L^{nx_L + e_L} \\ y &= u_y p_1^{ny_1 + e_1} \cdots p_L^{ny_L + e_L} \\ z &= u_z p_1^{nz_1 + e_1} \cdots p_L^{nz_n + e_L} \end{aligned}$

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 $x = u_x p_1^{nx_1+e_1} \cdots p_L^{nx_L+e_L}$ $y = u_y p_1^{ny_1+e_1} \cdots p_L^{ny_L+e_L}$ $z = u_z p_1^{nz_1+e_1} \cdots p_L^{nz_n+e_L}$ x + y = z

Let COL: $\mathbb{N} \to \{0, \dots, n-1\}^L$ be the following coloring:

$$\operatorname{COL}(up_1^{a_1}\cdots p_L^{a_L}) = (a_1 \pmod{n}, \dots, a_L \pmod{n})$$

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Not True Fix *n*. Let $u_x = u_y = \frac{1}{2}$, $u_z = 1$, X = Y = Z = 1.

$$u_x X^n + u_y Y^n = u_z Z^n$$

Becomes

$$\frac{1}{2}1^{n} + \frac{1}{2}1^{n} = 1 \times 1^{n}$$
$$\frac{1}{2} + \frac{1}{2} = 1$$

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The EG-proof that there are an infinite number of primes (in $\mathbb N)$ did not transfer to $\mathbb Q$ because

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The EG-proof that there are an infinite number of primes (in \mathbb{N}) did not transfer to \mathbb{Q} because Its NOT that **FLT** is false over \mathbb{Q} . Indeed—FLT is true over \mathbb{Q} (follows form FLT being true over \mathbb{Z}).

Its because the following **variant** of FLT is false for \mathbb{Q} : There exists $n \in \mathbb{N}$ such that the following has no solution:

$$u_x X^n + u_y Y^n = u_z Z^n$$

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where $u_x, u_y, y_z \in \mathbb{U}$ and $X, Y, Z \in \mathbb{Q}$.

May 10, 2025

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1. Read the Gasarch paper. Note that its initial proof was a generalization of what was presented here.

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- 1. Read the Gasarch paper. Note that its initial proof was a generalization of what was presented here.
- 2. Read in Gasarch's paper the **Sanity Check** which has more domains with a finite number of primes.
- Read the other papers on the website of Ramsey-Primes paper. Some of the papers are difficult so try to just figure out the proof for Z or N, and then see where it fails for Q and Q₂. (I think they all fail for AI because AI is not atomic, though check that.)