

An Exposition of Two Proofs of Can VDW

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July 23, 2025

1 Introduction

We need the following terminology.

Notation 1.1 Let $W \in \mathbb{N} \cup \omega$. Let $k, c \in \mathbb{N}$.

1. A *monochromatic k -AP* is an arithmetic sequence of length k where all of the elements of it are the same color. Henceforth we use the abbreviation *mono k -AP*.
2. A *rainbow k -AP* is an arithmetic sequence of length k where all of the elements have different colors. Henceforth we use the abbreviation *RB k -AP*.

Recall VDW's Theorem:

Theorem 1.2 *For all k, c there exists $W = W(k, c)$ such that for all $\text{COL}: [W] \rightarrow [c]$ there is a mono k -AP.*

What if there is no bound on the number of colors? Clearly there is a coloring of ω without a mono k -AP: $\text{COL}(x) = x$. However, in that coloring, you have a k -RB. This counterexample motivates the Can VDW Theorem:

Theorem 1.3 *For all k there exists $C = W(k)$ such that for all $\text{COL}: [C] \rightarrow \mathbf{N}$ there is either a mono k -AP or a RB k -AP.*

We present several proofs of the Can VDW theorem and discuss their pros and cons.

2 Proof that Does Not Use Gallai-Witt

2.1 The Set Up

We will prove a generalization of the Can VDW theorem which effectively loads the induction hypothesis. There are several changes we make:

1) We replace the range \mathbf{N} with the range \mathbf{N}^m . This may seem trivial; however, the next item will use the new range in an interesting way.

Notation 2.1 Let $C, m \in \mathbf{N}$. Let $\text{COL}: [C] \rightarrow \mathbf{N}^m$. Then $\text{COL}_i: [C] \rightarrow \mathbf{N}^m$ is the function that, on input x , outputs the i th element of $\text{COL}(x)$.

2) We will **redefine** mono k -AP and RB k -AP.

Def 2.2 Let $k, m, C \in \mathbf{N}$. Let $\text{COL}: [C] \rightarrow \mathbf{N}^m$.

1. A *mono k -AP* is a sequence $a, a + d, \dots, a + (k - 1)d$ such that

$$(\exists 1 \leq i \leq m)[\text{COL}_i(a) = \text{COL}_i(a + d) = \dots = \text{COL}_i(a + (k - 1)d)].$$

Note that this new definition of mono k -AP is weaker than the usual definition of mono k -AP which would insist that

$$(\forall 1 \leq i \leq m)[\text{COL}_i(a) = \text{COL}_i(a + d) = \dots = \text{COL}_i(a + (k - 1)d)].$$

Also note that if $m = 1$ then the old and new definition of mono k -AP are the same.

2. A *RB k -AP* is a sequence $a, a + d, \dots, a + (k - 1)d$ such that

$$(\forall 0 \leq x < y \leq k - 1)(\forall 1 \leq i < j \leq m)[\text{COL}_i(a + xd) \neq \text{COL}_j(a + yd)].$$

For example, we have.

$$\{\text{COL}_1(a + 2d), \text{COL}_2(a + 2d), \dots, \text{COL}_2(a + (k - 1)d)\} \bigcap$$

$$\{\text{COL}_1(a + 7d), \text{COL}_2(a + 7d), \dots, \text{COL}_7(a + (k - 1)d)\} = \emptyset$$

Note that this new definition of RB k -AP is stronger than the usual definition of RB k -AP which would only requires

$$(\forall 0 \leq x < y \leq k - 1)(\exists 1 \leq i \leq m)[\text{COL}_i(a + xd) \neq \text{COL}_i(a + yd)].$$

Also note that if $m = 1$ then the old and new definition of RB k -AP are the same.

3) We will prove the following asymmetric version of the Can VDW: For all $k, t, m \in \mathbb{N}$ there exists C such that, for all $\text{COL}: [C] \rightarrow [\mathbb{N}]^m$, one of the following occurs:

- There is a mono k -AP.
- There is a RB t -AP.

This theorem implies the Can VDW theorem by setting $t = k$ and $m = 1$.

2.2 The Proof

Theorem 2.3 *For all $k, t, m \in \mathbb{N}$ there exists C such that, for all $\text{COL}: [C] \rightarrow [\mathbb{N}]^m$, one of the following occurs:*

- *There is a mono k -AP.*
- *There is a RB t -AP.*

Proof:

We prove this by induction on t .

Base Case: $t = 1$. For any k, m can take $C(k, 1, m) = 1$.

The $t = 1$ case is so trivial that we do the $t = 2$ case.

Base Case: $t = 2$. Let $\text{COL}: [C] \rightarrow \mathbb{N}^m$ where we determine C later.

If there is a RB 2-AP then we are done. So assume there is no RB 2-AP. Let

$$\text{COL}(1) = (c_1, \dots, c_m).$$

For all $x \geq 2$, there has to be an i, j such that $\text{COL}_i(x) = c_j$. We define a coloring using this fact.

$\text{COL}': [C] - \{1\} \rightarrow [m] \times [m]$ is defined by

$$\text{COL}'(x) = (i, j) \text{ where } \text{COL}_i(x) = c_j.$$

We want to use VDW's theorem on COL' . Let $C = W(k, m) + 1$. Then there exists a, d, i, j such that

$$\text{COL}'(a) = \text{COL}'(a + d) = \dots = \text{COL}'(a + (k - 1)d) = (i, j).$$

Note that

$$\text{COL}_i(a) = \text{COL}_i(a + d) = \dots = \text{COL}_i(a + (k - 1)d) = c_j.$$

Hence we get a mono k -AP.

Case that contains most of the ideas: $t = 3, m = 1$ We can assume that, for all k, m , $C(k, 2, m)$ exists.

Let $\text{COL}: [C] \rightarrow \mathbb{N}$ where we determine C later.

Break $[C]$ into (suggestively named, TBD) C' blocks of size (suggestively named, TBD)

m . Call the blocks $B_1, B_2, \dots, B_{C'}$.

We define a coloring

$\text{COL}': [C'] \rightarrow \mathbb{N}^m$ by

$\text{COL}'(x) = (\text{COL}_1(B_x), \text{COL}_2(B_x), \dots, \text{COL}_m(B_x))$

We take $C' \geq C(k, 2, m)$.

Inductively we have two cases.

Case 1: There is a mono k -AP, in which case we are done (overdone really since we would have a mono k -AP under the old definition)

Case 2: There is a RB 2-AP x, y . Let z be such that x, y, z is a 3-AP. In order for z to be in the domain of COL' . we need $C' \geq 2C(k, 2, m)$.

Let A, D, M be such that

- $B_x = (A, A + 1, \dots, M - 2, M - 1, M)$
- $B_y = (A + D, A + D + 1, \dots, M + D - 2, M + D - 1, M + D)$
- $B_z = (A + 2D, A + 2D + 1, \dots, M + 2D - 2, M + 2D - 1, M + 2D)$

Since x, y is a RB 2-AP we have

$$(1) \quad \{\text{COL}(A), \text{COL}(A+1), \text{COL}(A+2), \dots, \text{COL}(M)\} \cap \{\text{COL}(A+D), \text{COL}(A+D+1), \text{COL}(A+2D), \dots, \text{COL}(M+2D)\} \neq \emptyset$$

We consider $\text{COL}(M + 2d)$. Look at each 3-AP (written backwards)

$$\begin{aligned}
&M + 2D, M + D, M \\
&M + 2D, M + D - 1, M - 2 \\
&M + 2D, M + D - 2, M - 4 \\
&M + 2D, M + D - 3, M - 6 \\
&\vdots \\
&M + 2D, M + D - x, M - 2x \\
&\vdots \\
&M + 2D, M + D - \frac{M-A}{2}, A
\end{aligned}$$

Lets look at $M + 2D, M + D - x, M - 2x$. We know that $\text{COL}(M + D - x) \neq \text{COL}(M - 2x)$ by (1). If $\text{COL}(M + 2D) \notin \text{COL}(M + d - x) \neq \text{COL}(M - 2x)$ then we would have a RB 3-AP. This observation motivates the following coloring.

$$\text{COL}'': B_y \rightarrow [2]$$

$$\text{COL}''(w) = \begin{cases} 1 & \text{if } \text{COL}''(w) = \text{COL}(M + 2D) \\ 2 & \text{if } \text{COL}''(w) \neq \text{COL}(M + 2D) \end{cases} \quad (1)$$

We will take M (the size of the blocks) to be at least $W(2, k)$. Hence we apply VDWs Theorem to COL' and get two cases.

Case 2.1: There is a mono k -AP with color 1. The mono k -AP actually a mono k -AP with color COL.

Case 2.2: There is a mono k -AP with color 2. So we have

$\text{COL}(M + 2D) \neq \text{COL}(M + D)$ and $\text{COL}(M + D) \neq \text{COL}(M)$. If $\text{COL}(M + 2D) \neq \text{COL}(M)$ then we have a RB 3-AP hence we can assume $\text{COL}(M + 2D) = \text{COL}(M)$.

Similarly $\text{COL}(M + 2D) = \text{COL}(M - 2)$.

\vdots

Similarly, for $1 \leq x \leq 0$, $\text{COL}(M + 2D) = \text{COL}(M + D - x)$.

Hence we have the following mono k -AP.

$$A, A + 1, \dots, M$$

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