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A RESTRICTED VERSION OF HALES—JEWETT'S THEOREM

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1.

A well known theorem of van der Waerden states that for every pair  $\delta, k$  of positive integers there exists a positive integer  $n$  with the property that for every partition of  $\{1, \dots, n\}$  into  $\delta$  many classes there exists a  $k$ -term arithmetic progression contained in one class. Thus in order to obtain a  $k$ -term arithmetic progression within one class of the partition a much richer structure (viz an  $n$ -term progression) is partitioned. Erdős [1] conjectured that for every pair  $\delta, k$  of positive integers there exists a set  $A$  of positive integers which contains no  $(k + 1)$ -term progression and still has the property that for every partition of  $A$  into  $\delta$  many classes at least one of the classes contains a  $k$ -term progression. That such a restricted version of van der Waerden's theorem is valid was shown independently by Spencer [8] with the aid of Hales — Jewett's theorem [4] and Nešetřil — Rödl [6] by a direct construction. In this paper we shall give a restricted version of Hales — Jewett's theorem for partitioning  $0$  parameter sets.

2.

Let  $A$  be a finite alphabet.  $A^n$  is the set of words of length  $n$  over  $A$ .  $A^n$  may also be viewed as an  $n$ -dimensional cube over  $A$ . A  $k$ -dimensional subcube  $C$  of  $A^n$  has a parametric representation of the following type

$$C = \{(a_0, \dots, \lambda_0, \dots, \lambda_1, \dots, \lambda_{k-1}, \dots, a_{n-1}) \mid \lambda_i \in A, i < k\}.$$

This representation is given by the word

$$f = (a_0, \dots, \lambda_0, \dots, \lambda_1, \dots, \lambda_{k-1}, \dots, a_{n-1}).$$

This suggests the following definition:

**Definition 2.1.** Let  $A$  be a finite alphabet and  $n, k$  be nonnegative integers. The set  $[A] \binom{n}{k}$  of  $k$ -parameter words of length  $n$  over  $A$  is the set of all mappings  $f: n \rightarrow A \cup \{\lambda_i \mid i < k\}$  where  $A \cap \{\lambda_i \mid i < k\} = \emptyset$  satisfying

2.1.1. For every  $j < k$  there exists  $i < n$  such that  $f(i) = \lambda_j$ .

2.1.2.  $\min f^{-1}(\lambda_i) < \min f^{-1}(\lambda_j)$  whenever  $i < j < k$ .

2.1.1 guarantees that  $k$ -parameters really occur in  $f$  and 2.1.2 gives a natural ordering of the first occurrences of the parameters. With these conditions we obtain a bijection between the  $k$ -parameter words of length  $n$  over  $A$  and the  $k$ -dimensional subcubes of  $A^n$ . Next we define the composition of parameter words.

**Definition 2.2.** Let  $f \in [A] \binom{n}{m}$  and  $g \in [A] \binom{m}{k}$ . The  $k$ -parameter word  $f \cdot g \in [A] \binom{n}{k}$  is defined by

$$f \cdot g(i) = \begin{cases} f(i) & \text{if } f(i) \in A \\ g(j) & \text{if } f(i) = \lambda_j. \end{cases}$$

In geometric terms  $f \cdot g$  is the  $k$ -dimensional subcube of the  $m$ -dimensional subcube  $f$  which has parametric representation  $g$  in  $A^m$ .

**Theorem (Hales – Jewett [4]).** *Let  $A$  be a finite alphabet. For every pair  $m, \delta$  of positive integers there exists a positive integer  $n$  such that for every coloring  $\Delta: [A] \binom{n}{0} \rightarrow \delta$  of the points of the  $n$ -dimensional cube over  $A$  there exists a monochromatic  $f \in [A] \binom{n}{m}$ , i.e. the coloring  $\Delta_f: [A] \binom{m}{0} \rightarrow \delta$  defined by  $\Delta_f(g) = \Delta(f \cdot g)$  is constant.*

3.

Here we prove a restricted version of Hales–Jewett’s theorem: For positive integers  $\delta, m$  there exists a positive integer  $n$  and a subset  $S$  of  $[A] \binom{n}{0}$  such that  $S$  does not contain an  $(m + 1)$ -dimensional subcube, i.e. for every  $f \in [A] \binom{n}{m+1}$  there exists  $g \in [A] \binom{m+1}{0}$  such that  $f \cdot g \notin S$ , and still for every coloring of the points of  $S$  there exists a monochromatic  $m$ -dimensional subcube in  $S$ .

**Notation 3.1.** Let  $S \subseteq [A] \binom{n}{0}$  and  $k < n$ .

$$\begin{aligned} \mathcal{F}_k(S) &:= \{f \in [A] \binom{n}{k} \mid f \cdot [A] \binom{k}{0} \subseteq S\} = \\ &= \{f \in [A] \binom{n}{k} \mid f \cdot g \in S \text{ for every } g \in [A] \binom{k}{0}\}. \end{aligned}$$

$\mathcal{F}_k(S)$  is the set of all  $k$ -dimensional subcubes of  $A^n$  which are contained in  $S$ . In particular  $\mathcal{F}_0(S) = S$ .

We are ready to state the main theorem.

**Main theorem.** *Let  $A$  be a finite alphabet and  $\delta, m$  be positive integers. Then there exists a positive integer  $n$  and  $S \subseteq [A] \binom{n}{0}$  such that*

- ( $\alpha$ )  $\mathcal{F}_{m+1}(S) = \phi$ , i.e.  $S$  does not contain  $(m + 1)$ -dimensional subcubes.
- ( $\beta$ ) For every coloring  $\Delta: S \rightarrow \delta$  of the points of  $S$  with  $\delta$  many colors there exists  $f \in \mathcal{F}_m(S)$  such that the coloring  $\Delta_f: [A] \binom{m}{0} \rightarrow \delta$  defined by  $\Delta_f(g) = \Delta(f \cdot g)$  is constant.

**Notation 3.2.** A convenient abbreviation for  $(\alpha)$  and  $(\beta)$  is " $S \xrightarrow{A} (m)_\delta$ ".

The crucial part of the proof of the main theorem consists in showing its validity for  $m = 1$ . This case is stated in

**Lemma 3.3.** *Let  $A$  be a finite alphabet and  $\delta$  be a positive integer. There exists a positive integer  $n$  and a subset  $S$  of  $[A] \binom{n}{0}$  such that  $S \xrightarrow{A} (1)_\delta$ .*

The main theorem follows from this lemma by concatenation of various sets  $S$ .

**Definition 3.4.** Let  $f \in [A] \binom{m}{0}$  and  $g \in [A] \binom{n}{0}$ . The concatenation  $f \otimes g \in [A] \binom{m+n}{0}$  is defined by

$$f \otimes g(i) = \begin{cases} f(i) & \text{for } i < m \\ g(i-m) & \text{for } m \leq i < m+n. \end{cases}$$

Thus  $f \otimes g$  is the usual concatenation of words.

Next we derive the Main theorem from Lemma 3.3. Proceed by induction on  $m$ .  $m = 1$  is settled by the lemma. For  $m + 1$  consider  $S_0 \subseteq [A] \binom{n_0}{0}$  with  $S_0 \xrightarrow{A} (m)_\delta$ , which exists by induction hypothesis. Lemma 3.3 guarantees the existence of a positive integer  $n_1$  and  $S_1 \subseteq [A] \binom{n_1}{0}$  with  $S_1 \xrightarrow{A} (1)_{\delta'}$ , where  $\delta' = \delta^{|S_0|}$ .

**Claim.**  $S_0 \otimes S_1 \xrightarrow{A} (m+1)_\delta$  where  $S_0 \otimes S_1 = \{g \otimes h \mid g \in S_0, h \in S_1\}$ .

We have to check conditions  $(\alpha)$  and  $(\beta)$  of the Main theorem:

For  $(\alpha)$ : Let  $f \in [A] \binom{n_0+n_1}{m+2}$ .

In case  $\min f^{-1}(\lambda_m) < n_0$ , consider  $f^* \in [A] \binom{n_0}{m+1}$  defined by

$$f^*(i) = \begin{cases} f(i) & \text{for } i < n_0 \text{ and } f(i) \neq \lambda_{m+1} \\ \lambda_0 & \text{for } i < n_0 \text{ and } f(i) = \lambda_{m+1}. \end{cases}$$

Since  $f^* \notin \mathcal{F}_{m+1}(S_0) = \phi$  there exists  $g \in [A] \binom{m+1}{0}$  such that  $f^* \cdot g \notin S_0$ . Therefore  $f \notin \mathcal{F}_{m+2}(S_0 \otimes S_1)$ .

In case  $\min f^{-1}(\lambda_m) \geq n_0$  consider  $f^* \in [A] \binom{n_1}{2}$  defined by

$$f^*(i) = \begin{cases} f(i + n_0) & \text{if } f(i + n_0) \in A \\ \lambda_0 & \text{if } f(i + n_0) \in \{\lambda_0, \dots, \lambda_m\} \\ \lambda_1 & \text{if } f(i + n_0) = \lambda_{m+1}. \end{cases}$$

Again since  $f^* \notin \mathcal{F}_2(S_1) = \phi$ , there exists  $g \in [A] \binom{2}{0}$  such that  $f^* \cdot g \notin S_1$  and therefore  $f \notin \mathcal{F}_{m+2}(S_0 \otimes S_1)$ .

For ( $\beta$ ): Let  $\Delta: S_0 \otimes S_1 \rightarrow \delta$  be a coloring. Consider first the coloring  $\Delta_1: S_1 \rightarrow \delta^{|S_0|}$  defined by  $\Delta_1(\xi) = \Delta(g \otimes \xi | g \in S_0)$ . By choice of  $S_1$  there exists  $h \in \mathcal{F}_1(S_1)$  such that  $\Delta_1$  is constant on  $h$ . Consider next  $\Delta_0: S_0 \rightarrow \delta$  defined by  $\Delta_0(\eta) = \Delta(\eta \otimes ha)$ , where  $a \in [A] \binom{1}{0}$  may be chosen arbitrarily. By choice of  $S_0$  there exists a  $g \in \mathcal{F}_m(S_0)$  such that  $\Delta_0$  is constant on  $g$ . Thus  $\Delta$  is constant on  $g \tilde{\otimes} h \in \mathcal{F}_{m+1}(S_0 \otimes S_1)$  where

$$g \tilde{\otimes} h(i) = \begin{cases} g(i) & \text{for } i < n_0 \\ \lambda_m & \text{for } i \geq n_0 \text{ and } h(i - n_0) = \lambda_0 \\ h(i - n_0) & \text{for } i \geq n_0 \text{ and } h(i - n_0) \in A. \blacksquare \end{cases}$$

**Proof of Lemma 3.3.** Proceed by induction on  $|A|$ . For  $|A| = 1$  the lemma is obvious. For  $|A| = 2$ , e.g.  $A = \{0, 1\}$ , choose  $n = \delta$  and consider the following set  $S \subseteq \{[0, 1]\} \binom{n}{0}$ :

$$S := \left\{ \begin{array}{l} (0 \dots 0), \\ (0 \dots 0 \ 1), \\ (0 \dots 0 \ 1 \ 1), \\ \vdots \\ (1 \dots 1 \ 1), \end{array} \right\}$$

Obviously  $S$  does not contain a 2-dimensional subcube because each element of  $S$  has the property that all entries after the first occurrence of the letter 1 are 1, i.e. no 0 occurs. Thus for every  $f \in \{0, 1\}^{\binom{n}{2}}$  we have that  $f(1, 0) \notin S$ . Furthermore, any two elements of  $S$  form a 1-dimensional subcube. As  $|S| = n + 1 = \delta + 1$  the pigeon hole principle guarantees the existence of two elements of  $S$  having the same color for any coloring  $\Delta: S \rightarrow \delta$ .

The general step may be done by a somewhat tricky iteration of the ideas just mentioned. Consider the alphabet  $A = \{0, \dots, t-1, t\}$ . By induction hypothesis applied to  $A^* = \{0, \dots, t-1\}$ , for every  $\delta^*$  there exists  $n$  and  $S \subseteq [A] \binom{n}{0}$  satisfying  $S \xrightarrow{A^*} (1)_{\delta^*}$ . This guarantees the existence of positive integers  $n_0, \dots, n_{\delta-1}$  and sets  $S_i \subseteq [A^*] \binom{n_i}{0}$  ( $i < \delta$ ) satisfying

$$S_i \xrightarrow{A^*} (1)_{\delta_i}$$

where  $\delta_0 = \delta^\delta$  and  $\delta_{i+1} = \delta^{(\delta-1-i) \circ |S_i \times S_{i-1} \times \dots \times S_0|}$  for  $i < \delta - 1$ .

Let  $n = n_0 + \dots + n_{\delta-1}$ . In order to define  $S \subseteq [A] \binom{n}{0}$  consider the following scheme:

$$T = \begin{array}{l} A^* \times \dots \times A^* \cup \\ A^* \times \dots \times A^* \times \{t\} \cup \\ A^* \times \dots \times \{t\} \times \{t\} \cup \\ \vdots \\ \underbrace{\{t\} \times \dots \times \{t\}}_{\delta} \end{array}$$

Let

$$S = (\mathcal{F}_1(S_0) \otimes \dots \otimes \mathcal{F}_1(S_{\delta-1})) \cdot T$$

which is defined as follows

$$\begin{aligned} S = & \{f_0(a_0) \otimes \dots \otimes f_{\delta-1}(a_{\delta-1}) \mid f_i \in \mathcal{F}_1(S_i), a_i \in A^*, i < \delta\} \cup \\ & \cup \{f_0(a_0) \otimes \dots \otimes f_{\delta-2}(a_{\delta-2}) \otimes f_{\delta-1}(t) \mid f_i \in \mathcal{F}_1(S_i), \\ & \qquad \qquad \qquad a_i \in A^*, i < \delta\} \cup \\ & \cup \{f_0(a_0) \otimes \dots \otimes f_{\delta-3}(a_{\delta-3}) \otimes f_{\delta-2}(t) \otimes f_{\delta-1}(t) \mid f_i \in \mathcal{F}_1(S_i), \\ & \qquad \qquad \qquad a_i \in A^*, i < \delta\} \cup \\ & \vdots \\ & \cup \{f_0(t) \otimes \dots \otimes f_{\delta-1}(t) \mid f_i \in \mathcal{F}_1(S_i), i < \delta\}. \end{aligned}$$

A typical element of  $S$  consists of  $\delta$  blocks, where the  $i$ -th block is obtained by substituting a letter from  $A$  into  $f_i \in \mathcal{F}_1(S_i)$ . Moreover if the  $i$ -th block is obtained by substituting the letter  $t$  then in all following blocks the letter  $t$  is substituted. As the parameter words  $f_i \in \mathcal{F}_1(S_i) \subseteq [A] \binom{n_i}{1}$  are defined on the smaller alphabet  $A^*$  and do not contain the letter  $t$ , the only way that the letter  $t$  occurs in an element of  $S$  is by substitution of  $t$  for the parameter in a parameter word  $f_i$ . In order to show that  $S \xrightarrow{A} (1)_\delta$  we check conditions  $(\alpha)$ ,  $(\beta)$ .

For  $(\alpha)$ : In order to show that  $\mathcal{F}_2(S) = \phi$  assume that there exists  $h \in \mathcal{F}_2(S)$ . Decompose  $h$  into its blocks, i.e.  $h = h_0 \otimes h_1 \otimes \dots \otimes h_{\delta-1}$ , where the length of the block  $h_i$  is  $n_i$  ( $i < \delta$ ).

**Fact 1.** No block is a 2-parameter word. For otherwise it would be a 2-parameter word over the alphabet  $A^*$ , contradicting the choice of  $S_i$ .

**Fact 2.** No block contains simultaneously a parameter and the letter  $t$ . For otherwise such a block is a 2-parameter word over the alphabet  $A^*$  in which  $t$  acts as a parameter, contradicting the choice of  $S_i$ .

**Fact 3.** Facts 1, 2 imply that  $h(t, a) \notin S$  for all  $a \in A^*$ . For



otherwise  $\min h^{-1}(\lambda_0) < \min h^{-1}(\lambda_1)$  implies that  $h(t, a)$  contains a block without letter  $t$  following blocks containing  $t$ . This contradicts the definition of  $S$  via the scheme  $T$ .

These facts show that  $\mathcal{F}_2(S) = \phi$ .

For  $(\beta)$ : Let  $\Delta: S \rightarrow \delta$  be a coloring. Consider first the coloring

$$\Delta_{\delta-1}: S_{\delta-1} \rightarrow \delta_{\delta-1} = \delta^{|S_{\delta-2} \times \dots \times S_0|}$$

defined by

$$\Delta_{\delta-1}(g) = (\Delta(g_0 \otimes \dots \otimes g_{\delta-2} \otimes g) \mid g_i \in S_i, i < \delta - 1),$$

where for convenience  $\Delta(g_0 \otimes \dots \otimes g_{\delta-2} \otimes g)$  shall be 0 if  $g_0 \otimes \dots \otimes g_{\delta-2} \otimes g$  does not belong to  $S$ .

By choice of  $S_{\delta-1}$  there exists  $f_{\delta-1} \in \mathcal{F}_1(S_{\delta-1})$  monochromatic for  $\Delta_{\delta-1}$ . Consider next the coloring

$$\Delta_{\delta-2}: S_{\delta-2} \rightarrow \delta_{\delta-2} = \delta^{2|S_{\delta-3} \times \dots \times S_0|}$$

defined by

$$\Delta_{\delta-2}(g) = (\Delta(g_0 \otimes \dots \otimes g_{\delta-3} \otimes g \otimes f_{\delta-1}(\alpha)) \mid g_i \in S_i$$

for  $i < \delta - 2$  and  $\alpha \in A$ ).

By choice of  $f_{\delta-1}$ :

$$\begin{aligned} \Delta(g_0 \otimes \dots \otimes g_{\delta-3} \otimes g \otimes f_{\delta-1}(\alpha)) &= \\ &= \Delta(g_0 \otimes \dots \otimes g_{\delta-3} \otimes g \otimes f_{\delta-1}(\beta)) \end{aligned}$$

for all  $\alpha, \beta \in A^*$ . Thus the only relevant distinction whether  $\alpha \in A^*$  or  $\alpha = t$  contributes a factor 2 in the exponent of the number of colors of  $\Delta_{\delta-2}$ . By choice of  $S_{\delta-2}$  there exists  $f_{\delta-2} \in \mathcal{F}_1(S_{\delta-2})$  monochromatic for  $\Delta_{\delta-2}$ . Proceed iteratively and consider for  $i = 1, \dots, \delta$

$$\Delta_{\delta-i}: S_{\delta-i} \rightarrow \delta_{\delta-i} = \delta^{i|S_{\delta-i-1} \times \dots \times S_0|}$$

defined by

$$\Delta_{\delta-i}(g) = (\Delta(g_0 \otimes \dots \otimes g_{\delta-i-1} \otimes g \otimes f_{\delta-i+1}(\alpha_1) \otimes \dots \\ \dots \otimes f_{\delta-i+j}(\alpha_j) \otimes f_{\delta-i+j+1}(t) \otimes \dots \otimes f_{\delta-1}(t)) | \\ g_0 \in S_0, \dots, g_{\delta-i-1} \in S_{\delta-i-1}, \alpha_1 \dots \alpha_j \in A^*, 0 \leq j < i).$$

Observe again that in the definition of  $\Delta_{\delta-1}$  the particular choice of the letters  $\alpha$  is not relevant, i.e. for  $\alpha_1 \dots \alpha_j, \beta_1 \dots \beta_j \in A^*$

$$\Delta(g_0 \otimes \dots \otimes g_{\delta-i-1} \otimes g \otimes f_{\delta-i+1}(\alpha_1) \otimes \dots \\ \dots \otimes f_{\delta-i+j}(\alpha_j) \otimes f_{\delta-i+j+1}(t) \otimes \dots \otimes f_{\delta-1}(t)) = \\ = \Delta(g_0 \otimes \dots \otimes g_{\delta-i-1} \otimes g \otimes f_{\delta-i+1}(\beta_1) \otimes \dots \\ \dots \otimes f_{\delta-i+j}(\beta_j) \otimes f_{\delta-i+j+1}(t) \otimes \dots \otimes f_{\delta-1}(t)).$$

By choice of  $S_{\delta-i}$  find  $f_{\delta-i} \in \mathcal{F}_1(S_{\delta-i})$  monochromatic for  $\Delta_{\delta-i}$ .

Finally consider the  $\delta$ -parameter word

$$f = f_0 \otimes \dots \otimes f_{\delta-1} \in \mathcal{F}_1(S_0) \otimes \dots \otimes \mathcal{F}_1(S_{\delta-1}).$$

The induced coloring  $\Delta_f$  acts constantly on the rows of scheme  $T$ . Thus the  $\delta + 1$  many rows are colored with  $\delta$  colors. Two of them have the same color, thus defining a 1-parameter word which is monochromatic for  $\Delta$ . ■

#### 4.

An immediate corollary of Hales–Jewett’s theorem is a partition theorem for affine points:

**Theorem 4.1 [3].** *Let  $F$  be a finite field and let  $\delta, m$  be positive integers. Then there exists a positive integer  $n$  such that for every coloring of the affine points in the  $n$ -dimensional affine space over  $F$  with  $\delta$ -many colors, i.e. for every coloring  $\Delta: F^n \rightarrow \delta$ , there exists a monochromatic  $m$ -dimensional affine subspace.*

Luckily enough the configuration  $S$  which has been constructed in the proof of the Main theorem also yields the following restricted version of the above theorem, viz.

**Theorem 4.2.** Let  $F$  be a finite field and let  $\delta, m$  be positive integers. Then there exist a positive integer  $n$  and a set  $S \subseteq F^n$  of affine points such that

( $\alpha$ )  $S$  does not contain an  $(m + 1)$ -dimensional affine space,

( $\beta$ ) For every coloring of the points of  $S$  with  $\delta$  many colors there exists an affine  $m$ -dimensional monochromatic subspace which is contained in  $S$ .

**Proof.** Again the crucial case is  $m = 1$ , which may be proved as follows: Let  $F = \{0, 1, \dots, t\}$  be the finite field, where  $0$  is the zero element of  $F$ . Let  $S \subseteq F^n = [F] \binom{n}{0}$  be the set constructed in the proof of Lemma 3.3.

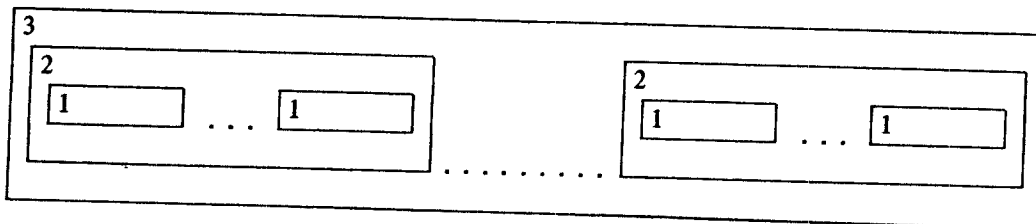
We claim that  $S$  has the desired properties. As in particular each  $f \in [F] \binom{n}{m}$  represents an  $m$ -dimensional affine subspace of  $F^n$  (compare e.g. [5]) property ( $\beta$ ) follows immediately. In order to show ( $\alpha$ ) we have to be more careful because there exist certain  $m$ -dimensional affine subspaces which are not represented by any  $f \in [F] \binom{n}{m}$ .

As  $S$  is defined recursively we may think of members of  $S$  being arranged in blocks of blocks of ... of blocks. There are  $t$  "levels" of blocks. The blocks of the first level are always of the form

$$0 \dots 0 \ a \dots a \ 1 \dots 1,$$

where  $a \in F$ .

The following diagram depicts the structure of  $S$  for  $GF(4)$ , where the small numbers indicate the levels of the corresponding blocks:



Let  $L = \{\bar{c} + \alpha\bar{x} \mid \alpha \in F\}$  be an affine line and assume that  $L$  is contained in  $S$ .

We examine the implications of the assumption " $L \subseteq S$ " on the vectors  $\bar{c}$  and  $\bar{x}$ . In particular we look at the structure of  $\bar{c}$  and  $\bar{x}$  in the blocks of level 1.

**Claim.** Let  $B$  be a block of level 1. Then the parts of  $\bar{c}$  and  $\bar{x}$  belonging to this block  $B$  – which will be denoted by  $\bar{c}_B$ , resp.  $\bar{x}_B$  – are of the form

$$\bar{c}_B = 0 \dots 0 \ b \dots b \ 1 \dots 1$$

$$\bar{x}_B = 0 \dots 0 \ x \dots x \ 0 \dots 0,$$

where  $b, x \in \mathbf{F}$ .

**Proof of claim.** Let  $B$  be any block of level 1 and assume that  $\bar{x}_B$  is not the zero vector. Let  $\bar{c}_B = c_1 \dots c_n$  and  $\bar{x}_B = x_1 \dots x_n$ .

1. Assume that  $x_i \neq 0$  and  $c_j = 0$ , where  $i < j$ , then  $x_i = x_j$ . For otherwise consider  $\bar{c} + \left(\frac{1}{x_i}\right)\bar{x}$  and find a block of level 1 with an entry 1 preceding an entry different from 1, again contradicting the structure of blocks of level 1.

2. Assume that  $x_i \neq 0$ ,  $c_i = 0$  and  $c_j \neq 0$  for some  $j > i$ , then  $x_j = 0$ . For otherwise consider  $\bar{c} - \left(\frac{c_j}{x_j}\right)\bar{x}$  and find a block of level 1 with a non-zero entry preceding a zero entry, again a contradiction.

3. Assume that  $c_i = 0$  and  $c_j \neq 0$ ,  $c_j \neq 1$  for some  $i < j$ , then  $x_i = 0$ . For otherwise 1. and 2. imply that

$$\bar{c}_B = 0 \dots 0 \ 0 \dots 0 \ a \dots a \ 1 \dots 1,$$

where  $a \neq 0, 1$  and

$$\bar{x}_B = 0 \dots 0 \ x \dots x \ 0 \dots 0 \ 0 \dots 0.$$

Consider  $\bar{c} + \left(\frac{1}{x}\right)\bar{x}$  and find a block of level 1 with an entry 1 preceding an entry different from 1, a contradiction.

Now the claim may be proved as follows: let  $i$  be the minimal index

such that  $x_i \neq 0$  and let  $\alpha \in \mathbf{F}$  be such that for  $\bar{d} = \bar{c} + \alpha\bar{x}$ , where  $\bar{d}_B = d_1 \dots d_n$ , it follows that  $d_i = 0$ . By 1., 2. and 3. above then  $\bar{d}_B$  and  $\bar{x}_B$  have the following structure:

$$\bar{d}_B = 0 \dots 0 \ 0 \dots 0 \ 1 \dots 1$$

$$\bar{x}_B = 0 \dots 0 \ x \dots x \ 0 \dots 0,$$

but then  $\bar{c} = \bar{d} - \alpha\bar{x}$  and thus

$$\bar{c}_B = 0 \dots 0 \ b \dots b \ 1 \dots 1$$

$$\bar{x}_B = 0 \dots 0 \ x \dots x \ 0 \dots 0,$$

where  $b = -\alpha x$  gives the desired result. ■

Let  $P = \{\bar{c} + \alpha\bar{x} + \beta\bar{y} \mid \alpha, \beta \in \mathbf{F}\}$  be a configuration of points which is contained in  $S$ . We show that  $\bar{x}$  and  $\bar{y}$  are linearly dependent, thus  $S$  does not contain a 2-dimensional affine subspace.

As an immediate corollary from the claim one obtains that for every block  $B$  of level 1 the vectors  $\bar{c}_B, \bar{x}_B$  and  $\bar{y}_B$  have the following structure:

$$\bar{c}_B = 0 \dots 0 \ b \dots b \ 1 \dots 1$$

$$\bar{x}_B = 0 \dots 0 \ x \dots x \ 0 \dots 0$$

$$\bar{y}_B = 0 \dots 0 \ y \dots y \ 0 \dots 0.$$

The proof of Theorem 4.2 will be finished by showing that for every  $k = 1, \dots, t$  and every block  $B$  of level  $k$  the following holds:

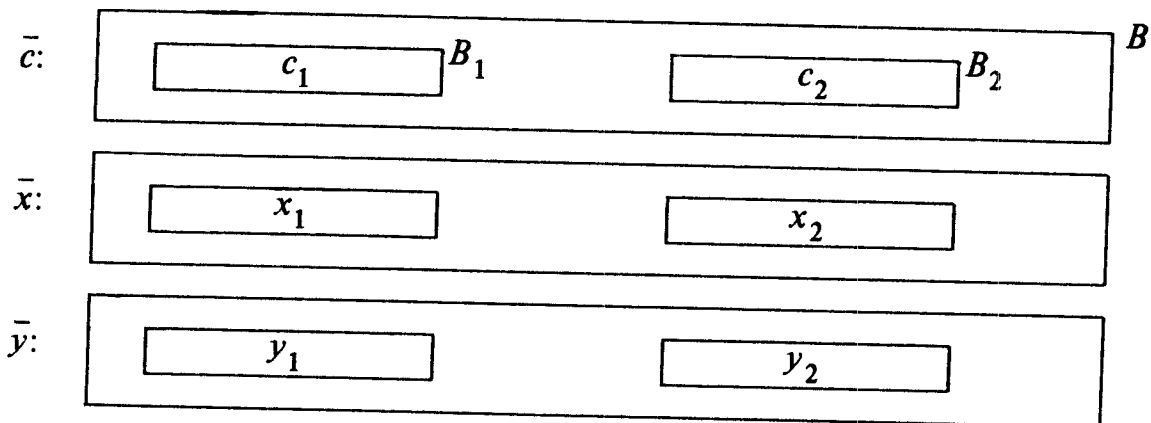
(\*)  $\bar{x}_B$  and  $\bar{y}_B$  are linearly dependent

(\*\*) if  $k < t$  and if for some  $\alpha, \alpha', \beta, \beta' \in \mathbf{F}$  both  $(\bar{c} + \alpha\bar{x} + \beta\bar{y})_B$  and  $(\bar{c} + \alpha'\bar{x} + \beta'\bar{y})_B$  contain at least one entry of  $t$  then all entries of  $t$  in  $(\bar{c} + \alpha\bar{x} + \beta\bar{y})_B$  and  $(\bar{c} + \alpha'\bar{x} + \beta'\bar{y})_B$  occur exactly at the same positions.

The assertions (\*) and (\*\*) are proved by induction on  $k$ . The case  $k = 1$  follows from the claim. So assume that (\*) and (\*\*) hold for  $k - 1$

and consider a block  $B$  of level  $k > 1$ . Suppose  $(*)$  does not hold on  $B$ . Then there must be two  $(k - 1)$ -level blocks  $B_1, B_2$  in  $B$  on the entries of the union of which  $\bar{x}$  and  $\bar{y}$  are linearly independent.

By induction  $\bar{x}_{B_1}$  and  $\bar{y}_{B_1}$  are linearly dependent and  $\bar{x}_{B_2}$  and  $\bar{y}_{B_2}$  are linearly dependent. Let  $c_1, x_1, y_1$  be the entries in the first position of  $B_1$  where not both  $x_1, y_1$  are equal to zero and let  $c_2, x_2, y_2$  be similarly defined for  $B_2$ . The situation is as follows:



Then  $x_1 y_2 \neq x_2 y_1$  by choice of  $B_1$  and  $B_2$  and the induction assumption on  $B_1, B_2$ . Thus we can find  $\alpha, \beta$  so that  $c_1 + \alpha x_1 + \beta y_1 = t$  but  $c_2 + \alpha x_2 + \beta y_2 \neq t$ . From property  $(**)$  we see that  $(\bar{c} + \alpha \bar{x} + \beta \bar{y})_{B_2}$  does not contain any  $t$ . This violates the conditions for  $S$ . Hence  $(*)$  holds on  $B$ .

Now let  $k < t$ . We show that  $(**)$  holds for  $B$ . Assume that it does not. By  $(*)$  we need only consider vectors of form  $\bar{c} + \alpha \bar{x}$  on  $B$ . Then for some  $\alpha, \alpha'$  we must have  $(\bar{c} + \alpha \bar{x})_B$  and  $(\bar{c} + \alpha' \bar{x})_B$  with different entries of  $t$ . In particular there exists a  $(k - 1)$ -level block  $B_1$  in  $B$  such that  $(\bar{c} + \alpha \bar{x})_{B_1}$  and  $(\bar{c} + \alpha' \bar{x})_{B_1}$  have different entries of  $t$ . By induction on  $(**)$  then one of these blocks, say  $(\bar{c} + \alpha \bar{x})_{B_1}$ , contains some entries of  $t$  while  $(\bar{c} + \alpha' \bar{x})_{B_1}$  does not contain any  $t$ . In particular there exists another  $(k - 1)$ -level block  $B_2$  such that  $(\bar{c} + \alpha' \bar{x})_{B_2}$  contains some entries of  $t$ . By the rules for  $S$  the block  $B_1$  precedes the block  $B_2$  and thus also  $(\bar{c} + \alpha \bar{x})_{B_2}$  contains some entries of  $t$ . Again by

induction on (\*\*) then all entries of  $t$  in  $(\bar{c} + \alpha\bar{x})_{B_2}$  and  $(\bar{c} + \alpha'\bar{x})_{B_2}$  occur exactly at the same positions. This means that for a suitable  $\alpha''$ , the vector  $\bar{c} + \alpha''\bar{x}$  has entries of  $t$  in  $B_2$  and entries of  $t - 1$  in  $B_1$ . This violates the rules for  $S$ , which exclude any entry of  $t - 1$  preceding an entry of  $t$  in any block of level  $t - 1$  or smaller. Thus (\*\*) holds. The induction is now complete and letting  $k = 1$  we see that (\*) implies that  $P = \{\bar{c} + \alpha\bar{x} + \beta\bar{y} \mid \alpha, \beta \in \mathbb{F}\}$  is just a line at best and not a plane. ■

Now much is known about restricted versions of Graham and Rothschild's partition theorem for  $k$ -parameter words [3] with  $k > 0$ . The only result in this direction is due to Nešetřil and Rödl [7] who announced the case of 2-parameter words over the empty alphabet:

**Theorem 4.1** [7]. *Let  $m, \delta$  be positive integers. There exist a positive integer  $n$  and a subset  $S$  of  $[\phi] \binom{n}{2}$  not containing an  $(m + 1)$ -dimensional subcube (i.e. for all  $g \in [\phi] \binom{n}{m+1}$  there exists an  $h \in [\phi] \binom{m+1}{2}$  with  $g \cdot h \notin S$ ) such that for every coloring  $\Delta: S \rightarrow \delta$  there exists an  $f \in [\phi] \binom{n}{m}$  which is contained  $S$  and monochromatic with respect to  $\Delta$ .*

Even for the empty alphabet the general case remains unsettled. Also nothing is known about restricted versions of the partition theorems for finite vector spaces [2].

Finally it could be worthwhile to note that in case of colorings of 0-parameter words (i.e.  $k = 0$ ) and requiring a monochromatic 1-parameter word not only a restricted but also simultaneously a restricted and induced version may be established, viz.

**Theorem 4.3.** *Let  $A$  be a finite set and let  $I \subseteq A$  be a subset of  $A$ . Then for every positive integer  $\delta$  there exists a positive integer  $n$  and a set  $S \subseteq [A] \binom{n}{0}$  such that*

( $\alpha$ )  $S$  does not contain a 2-parameter word,

( $\beta$ ) for every coloring  $\Delta: S \rightarrow \delta$  there exists a 1-parameter word

$f \in [A] \binom{n}{1}$  such that

- (i) all elements  $f \cdot a$ ,  $a \in I$ , are colored the same, and thus particularly  $f \cdot a \in S$  for  $a \in I$ ;
- (ii)  $f \cdot a \notin S$  for  $a \notin I$ .

This strengthens a result of [9], where an induced version of Hales–Jewett’s theorem has been established. Obviously the same strengthening applies to Theorem 4.2 as well. We do not know whether in general a restricted and simultaneously induced version of Hales–Jewett’s theorem is valid.

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