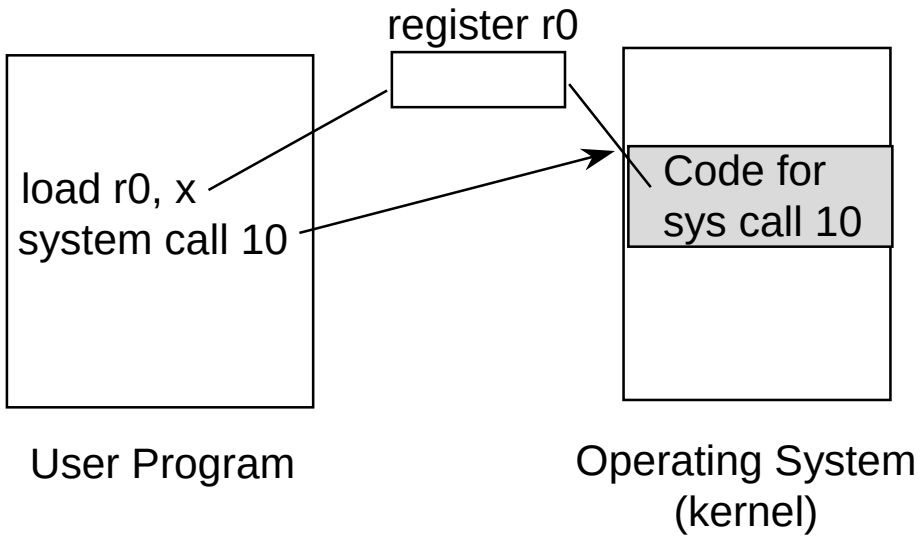
Announcements

- Program #1
  - Is on the web
- Reading
  - Chapter 3
  - Chapter 4 (for Thursday)

System Calls

- Provide the interface between application programs and the kernel
- Are like procedure calls
  - take parameters
  - calling routine waits for response
- Permit application programs to access protected resources



System Call Mechanism

- Use numbers to indicate what call is made
- Parameters are passed in registers or on the stack
- Why do we use indirection of system call numbers rather than directly calling a kernel subroutine?
  - provides protection since the only routines available are those that are export
  - permits changing the size and location of system call implementations without having to re-link application programs

Types of System Calls

- File Related
  - open, create
  - read, write
  - close, delete
  - get or set file attributes
- Information
  - get time
  - set system data (OS parameters)
  - get process information (id, time used)
- Communication
  - establish a connection
  - send, receive messages
  - terminate a connection
- Process control
  - create/terminate a process (including self)

System Structure

- Simple Structure (or no structure)
  - any part of the system may use the functionality of the rest of the system
  - MS-DOS (user programs can call low level I/O routines)
- Layered Structure
  - layer n can only see the functionality that layer n-1 exports
  - provides good abstraction from the lower level details
    - new hardware can be added if it provides the interface required of a particular layer
  - system call interface is an example of layering
  - can be slow if there are too many layers
- Hybrid Approach
  - most real systems fall somewhere in the middle

Policy vs. Mechanism

- Policy - what to do
  - users should not be able to read other users files
- Mechanism- how to accomplish the goal
  - file protection properties are checked on open system call
- Want to be able to change policy without having to change mechanism
  - change default file protection
- Extreme examples of each:
  - micro-kernel OS - all mechanism, no policy
  - MACOS - policy and mechanism are bound together

Processes

- What is a process?
  - a program in execution
  - “An execution stream in the context of a particular state”
  - a piece of code along with all the things the code can affect or be affected by.
    - this is a bit too general. It includes all files and transitively all other processes
  - only one thing happens at a time within a process
- What’s not a process?
  - program on a disk - a process is an active object, but a program is just a file

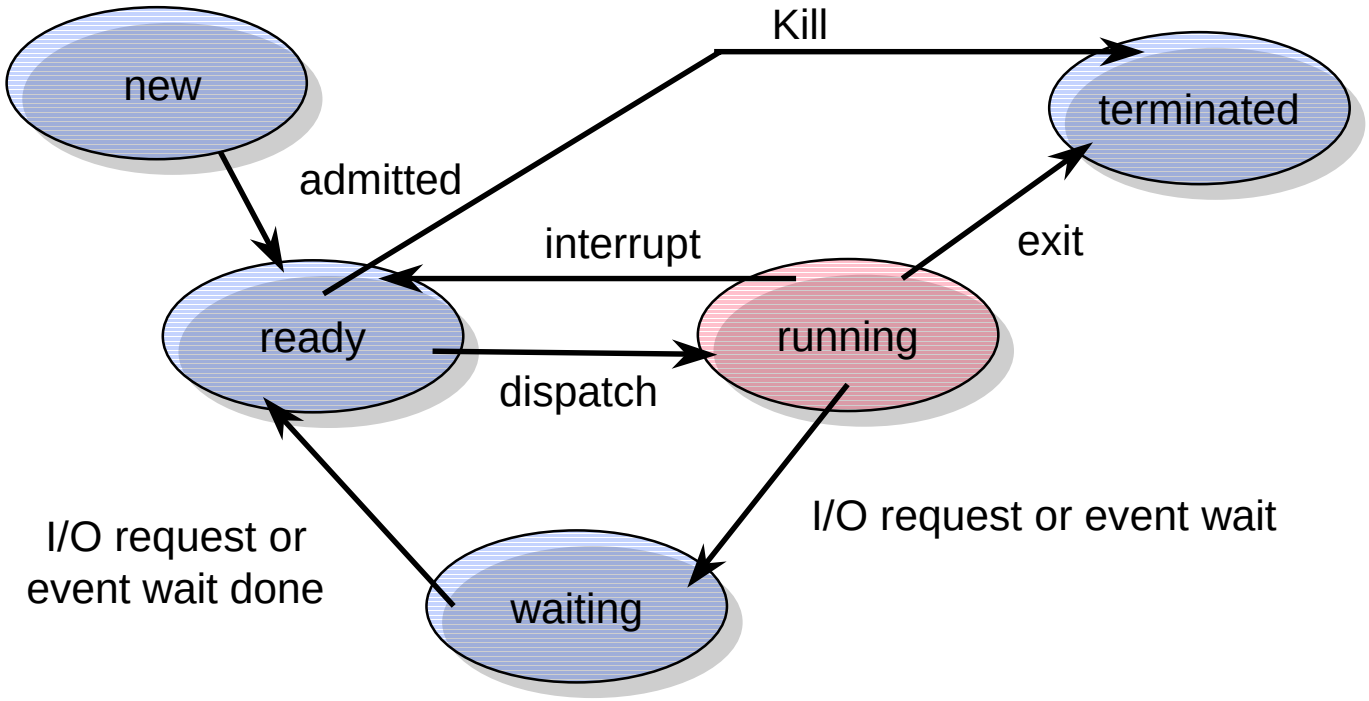
Multi-programming

- Systems that permit more than one process at once
  - virtually all computers today
- Permits more efficient use of resources
  - while one process is waiting another can run
- Provides natural abstraction of different activities
  - windowing system
  - editor
  - mail daemon
- Preemptive vs. non-preemptive muti-programming
  - preemptive means that a process can be forced off the processor by the OS
  - provides processor protection

Process State

- Processes switch between different states based on internal and external events
- Each process is in exactly one state at a time
- Typical States of Processes (varies with OS)
  - New: The process is just being created
  - Running: Instructions are being executed
    - only one process per processor may be running
  - Waiting: The process is waiting for an event to occur
    - examples: I/O events, signals
  - Ready: The process is waiting to be assigned to a processor
  - Terminated: The process has finished execution

Process State Transitions



Components of a Process

- Memory Segments
  - Program - often called the text segment
  - Data - global variables
  - Stack - contains activation records
- Processor Registers
  - program counter - next instruction to execute
  - general purpose CPU registers
  - processor status word
    - results of compare operations
  - floating point registers