

Announcements

- Program #1
 - See additional tips in hot news
 - Due in one week
- Reading
 - Chapter 7 (this entire week)

Interprocess Communication

- **Communicating processes establish a link**
 - can more than two processes use a link?
 - are links one way or two way?
 - how to establish a link
 - how do processes name other processes to talk to
 - use the process id (signals work this way)
 - use a name in the filesystem (UNIX domain sockets)
 - indirectly via mailboxes (a separate object)
- **Use send/receive functions to communicate**
 - send(dest, message)
 - receive(dest, message)

Producer-consumer pair

- producer creates data and sends it to the consumer
- consumer read the data and uses it
- examples: compiler and assembler can be used as a producer consumer pair
- Buffering
 - processes may not produce and consume items one by one
 - need a place to store produced items for the consumer
 - called a buffer
 - could be fixed size (bounded buffer) or unlimited (unbounded buffer)

Message Passing

- What happens when a message is sent?
 - sender blocks waiting for receiver to receive
 - sender blocks until the message is on the wire
 - sender blocks until the OS has a copy of the message
 - sender blocks until the receiver responds to the message
 - sort of like a procedure call
 - could be expanded into a remote procedure call (RPC) system
- Error cases
 - a process terminates:
 - receiver could wait forever
 - sender could wait or continue (depending on semantics)
 - a message is lost in transit
 - who detects this? could be OS or the applications
- Special case: if 2 messages are buffered, drop the older one
 - useful for real-time info systems

Signals (UNIX)

- provide a way to convey one bit of information between two processes (or OS and a process)
- types of signals:
 - change in the system: window size
 - time has elapsed: alarms
 - error events: segmentation fault
 - I/O events: data ready
- are like interrupts
 - a processes is stopped and a special handler function is called
- a fixed set of signals is normally available

Producer-consumer: shared memory

- Consider the following code for a producer

```
repeat
    ....
    produce an item into nextp
    ...
    while counter == n;
    buffer[in] = nextp;
    in = (in+1) % n;
    counter++;
until false;
```

- Now consider the consumer

```
repeat
    while counter == 0;
    nextc = buffer[out];
    out = (out + 1) % n;
    counter--;
    consume the item in nextc
until false;
```

- Does it work? **Answer: NO!**

Problems with the Producer-Consumer Shared Memory Solution

- Consider the three address code for the counter

Counter Increment

$\text{reg}_1 = \text{counter}$

$\text{reg}_1 = \text{reg}_1 + 1$

$\text{counter} = \text{reg}_1$

Counter Decrement

$\text{reg}_2 = \text{counter}$

$\text{reg}_2 = \text{reg}_2 - 1$

$\text{counter} = \text{reg}_2$

- Now consider an ordering of these instructions

T_0 producer $\text{reg}_1 = \text{counter}$ { $\text{reg}_1 = 5$ }

T_1 producer $\text{reg}_1 = \text{reg}_1 + 1$ { $\text{reg}_1 = 6$ }

T_2 consumer $\text{reg}_2 = \text{counter}$ { $\text{reg}_2 = 5$ }

T_3 consumer $\text{reg}_2 = \text{reg}_2 - 1$ { $\text{reg}_2 = 4$ }

T_4 producer $\text{counter} = \text{reg}_1$ { $\text{counter} = 6$ }

T_5 consumer $\text{counter} = \text{reg}_2$ { $\text{counter} = 4$ }

 This should be 5!