

More on comparison-based sorting

Note Title

11/7/2007

Is it possible to sort in linear time?

What is the worst-case lower bound
for comparison-based sorting?

$$\Omega(n \log n)$$

Do any comparison-based
sorting algorithms have
a linear best case?

Insertion Sort is $\Omega(n)$

Some features of comparison-based
sorting:

- insertion sort does very well when input almost sorted
- also practically good when n is small
- quicksort often fast in practice but it depends on the input and the pivots chosen

In addition to comparisons,

we also care about

- memory consumed
- # of reads
- # of writes

Heapsort and quicksort can be
executed in place.

Can hybrid algorithms get us improved performance?

Example: Insertion Sort

- in iteration i , insert i th element into sorted sublist
- use binary search to find the place to put i th value in list

$O(\log n)$ comparisons to insert $\rightarrow O(n)$ shifting
 $O(n)$ insertions
 $O(n \log n)$ in worst case? ?



Complications

- ① duplicate elements
- ② non-primitive keys

First Name
Last Name
Test Score

} what if you
sort by
Last Name, First
Name?

- ③ Data Stability

Say you have list

(Ada Smith 75, Bob Jones 80, Clara Barton 70,
George Mason 90, Shia LeBouf 65,
Clara Barton 95)

Notice "Clara Barton" is in the list twice.

DATA STABILITY means that items with the same keys stay in the same relative positions.

input: (Ada Smith 75, Bob Jones 80, Clara Barton 70,
George Mason 90, Shia LeBouf 65,
Clara Barton 95)

STABLE
OUTPUT

Clara Barton 70
Clara Barton 95
Bob Jones 80
Shia LeBouf 65
George Mason 90
Ada Smith 75

not stable
output

Clara Barton 95
Clara Barton 70
Bob Jones 80
Shia LeBouf 65
George Mason 90
Ada Smith 75

← not in
order
input
arrived
in

Why is data stability useful?

Lower Bound of Comparison-Based Sorting is $n \log n$

- means worst-case of any algorithm that could ever be invented is $\Omega(n \log n)$.

Insertion Sort is $O(n^2)$
 $\Omega(n)$.

When you use $\Omega(f(n))$ to describe an algorithm,

you are saying on all inputs running time is at least $f(n)$.

→ best case is at least $f(n)$.

Worst-case of Insertion Sort is $\Omega(n \log n)$, in fact it is $\Omega(n^2)$.

