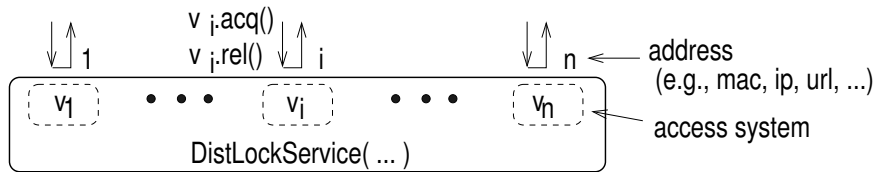


Distributed Lock Service

Shankar

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Overview: Distributed Lock Service



- Parameter: set of addresses
- Access system at each address
 - sid returned at instantiation
- Functions at address j , assuming access system v_j
 - input $v_j.acq()$: acquire lock
 - input $v_j.rel()$: release lock
- Termination function(s): can be added

- Main

- $ic \{ADDR \text{ not empty } \}$
- $eating \leftarrow null$ // user with lock if not null
- $users_j \leftarrow []$ // users at addr j
- $v_j \leftarrow sid()$ // sid of access system at j
- $return \{v_j\}$ // map of sids

- atomicity assumption: input and output parts

- progress assumptions

- $thread\ u\ in\ v_j.rel \text{ leads-to } not\ u\ in\ v_j.rel$
- $(eating \neq null \text{ leads-to } eating = null) \Rightarrow$
 $(u\ in\ users_j \text{ leads-to } u = eating)$

- `v[j].acq()`
 - `ic {eating \neq mytid}`
add `mytid` to `usersj`
 - `oc {eating = null}`
`eating` \leftarrow `mytid`

- `input v[j].rel()`
 - `ic {eating = mytid and
mytid in usersj}}` // caller acquired lock at j
remove `mytid` from `usersj`
 - `oc {true}`

Inverse of distributed lock service

- DistLockServiceInverse(ADDR, v)
 - main: ... \forall_j ...
 - output doAcq(j) input $v_j.acq()$
 - $i \in oc \{ \dots \}$... $v_j.acq()$
 - $e \in ic \{ \dots \}$...
 - output doRel(j) input $v_j.rel()$
 - $i \in oc \{ \dots \}$... $v_j.rel()$
 - $e \in ic \{ \dots \}$...
 - atomicity assumption: ...
 - progress assumption condition { ... }

Some naive implementations over fifo channel

- Centralized solution
 - fixed access system, say v_0 controls lock
 - acq and rel calls send msgs to v_0
 - v_0 serves acq-call msgs in fifo order
 - Token-circulating solution
 - “token” msg cycles through access systems
 - when an access system gets token
 - if local hungry user
return an acq call; wait for rel call
 - forward token to next access system
 - Ideal solution
 - request disturbs only non-thinking access systems
 - distributed path-reversal solution
- // chapter 16