

# SESF for Time-Constrained Programs

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- **Time-constrained** program
  - statements are subject to **time constraints**
- Distributed program: local time constraints cause global effects
- Statement  $A$  can have two kinds of time constraints
  - **only-within**:  $A$  executed only within a given time interval
  - **deadline**:  $A$  executed within a given time
- **Timing assumptions**: time constraints enforced by platform
  - eliminate some evolutions of the program
- **Timing properties**: time constraints to be satisfied by program

- Let  $S$  be a time-constrained program
- $S_\tau$ : **explicit-time version** of  $S$ 
  - “regular” program whose evolutions satisfy timing assumptions
  - “regular” reasoning/compositionality applies to  $S_\tau$
- $S_\tau$  is  $S$  with following added to environment
  - $\tau\text{now}$ : real-valued “current time” variable, readable by all
  - $\tau\text{age}(\delta)$ : “ageing” function that increases  $\tau\text{now}$  by  $\delta$
  - $S_\tau$ -transitions:  $S$ -transitions and ageing transitions
- **Epoch variables**: added to  $S$  to store times of step executions
- Timing assumptions: guards in epoch variables
  - only-within: guard on  $S$ -step
  - deadline: guard on  $\tau\text{age}(\delta)$
- Timing properties: assertions in epoch variables

# Explicit-time program of S

- Time-constrained program S
  - statements can simultaneously record value of  $\tau_{\text{now}}$
  - statements can have “only-within” blocking conditions in  $\tau_{\text{now}}$
  - S can have a deadline assumption, a predicate in  $\tau_{\text{now}}$
- Explicit-time version of S

```
program S $\tau$ (.)  
  real  $\tau_{\text{now}} \leftarrow 0$ ;  
  sys  $\leftarrow$  startSystem(S(.));           // only-within: guards in S  
  return mysid  
  
output  $\tau_{\text{age}}(\delta)$   
  oc { $\delta > 0$  and  $\langle \text{deadline}[\tau_{\text{now}} \mid \tau_{\text{now}} + \delta] \rangle$ }  
   $\tau_{\text{now}} \leftarrow \tau_{\text{now}} + \delta$   
  ic {true}
```

- Let A and B be time-constrained programs
- “A implements B” is now “ $A_\tau$  implements  $B_\tau$ ”
- Compositionality theorem holds
- Program-based implements holds
- Typical timing assumptions for implementation programs
  - $\text{delay}(L,U)$ : pass within  $[L,U]$  secs from arrival
  - $\text{ddl}\{U\}$ : reach next timing assumption within U
  - $\langle \text{blocking construct} \rangle(U)$ :  
pass within U secs since arrival or since last unblocked