

There is a long path in the divisor graph.

by

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Let D_n denote the graph whose vertices are the numbers $1, 2, \dots, n$ with an edge joining a and b if either $a|b$ or $b|a$. Hegyvári has posed the problem of estimating the length of the longest simple path in the divisor graph. He and Erdős have shown that a simple path can have length at most $(1-\log 2)n$, for large n , see [1]. Recently Pomerance [2] has shown that the longest path has length which is $o(n)$. In this note we obtain an estimate in the opposite direction. We exhibit a path of length

$$n \exp(-(2 + \epsilon) \sqrt{\log n \log \log n}).$$

From now on, by a path, we shall always mean a simple path, i.e., a path which passes through a vertex at most once. The length of a path will be the number of vertices through which it passes.

The existence of this long path in D_n depends on finding a long path in the inclusion graph I_m , on subsets of $\{a_1, \dots, a_m\}$. We shall need.

Lemma. If $k < m$ there is a path in the inclusion graph I_m which visits every k -element subset of $\{a_1, \dots, a_m\}$ but which visits no subset with more than $k + 1$ elements.

Note: If $k = m - 1$, then the Gray code is such a path.

Proof. Let $M_1(a_1, \dots, a_m)$ denote the path

$$a_1 \rightarrow a_1 a_2 \rightarrow a_2 \rightarrow \dots \rightarrow a_{m-1} a_m \rightarrow a_m.$$

Let $W_i(a_1, \dots, a_m)$ denote the reversal of the path $M_i(a_1, \dots, a_m)$. We define $M_i(a_1, \dots, a_m)$ recursively:

$$\begin{aligned}
M_i(a_1, \dots, a_m) &= a_1 \rightarrow a_1 w_{i-1}(a_2, \dots, a_m) \\
&\rightarrow a_2 \rightarrow a_2 w_{i-1}(a_3, \dots, a_m) \\
&\quad \vdots \\
&\rightarrow \quad \quad \quad \vdots \\
&\rightarrow a_{m-i+1} \rightarrow a_{m-i+1} w_{i-1}(a_{m-i+2}, \dots, a_m)
\end{aligned}$$

It is easily checked that the path $M_k(a_1, \dots, a_m)$ satisfies the lemma.

We now prove:

Theorem. Given $\epsilon > 0$ there is an n_0 , so that, for all $n > n_0$ there is a simple path in D_n of length at least

$$n \exp(-(2 + \epsilon) \sqrt{\log n \log \log n}).$$

Proof. Put $k + 1 = \lceil \sqrt{\log n / \log \log n} \rceil$. Let m be the greatest integer for which there is a prime $p_m < n^{1/k+1}$. Then, by the lemma, there is a simple path which visits $2^{\binom{m}{k}}$ distinct subsets of $\{p_1, \dots, p_m\}$ and no subsets with more than $k + 1$ elements. Therefore there is a path in D_n which visits $2^{\binom{m}{k}}$ vertices.

Now $\binom{m}{k} \geq \left(\frac{m}{k}\right)^k$ for $n > n_1$ and $m > (k + 1)n^{1/k+1} / e \log n$ for $n > n_2$, by the prime number theorem.

With this choice of k we then have

$$\binom{m}{k} > n \exp(-(2 + \epsilon) \sqrt{\log n \log \log n}) \text{ for all } n > n(\epsilon).$$

References

- [1] Erdős, Freud and Hegyvári. Arithmetical properties of permutations of integers. Acta Math. Acad. Sci. Hung (to appear).
- [2] Pommerance. On the longest simple path in the divisor graph. Congressus Numerantium (to appear).

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