# START RECORDING

# Sequences, Series and Summation / Product Notation

**CMSC 250** 

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**Outlining terms** 

- Typical notation:  $a: \mathbb{N} \to \mathbb{C}$
- Examples:
  - 1, 2, 3, 4, 5, ...
  - 1.5, 2.5, 3.5, ...
  - 1, 1, 1, 1, ....
  - $\sqrt{2}, \sqrt{3}, \sqrt{4}, \sqrt{5}, \sqrt{6}, \sqrt{7}...$

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    - 1, 2, 3, 4, 5, ... • 1.5, 2.5, 3.5, ... • 1, 1, 1, 1, ... •  $\sqrt{2}, \sqrt{3}, \sqrt{4}, \sqrt{5}, \sqrt{6}, \sqrt{7}...$ •  $a_n = 2^n, n = 0, 1, 2, ...$ •  $b_k = logk + 2k, k = 1, 2, 3, ...$  "Closed form" formula

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# Recursion: Good Idea?

• Example: Fibonacci

$$F_n = \begin{cases} 1, & \text{if } n = 0, 1 \\ F_{n-1} + F_{n-2}, \text{if } n \ge 2 \end{cases}$$

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• Recomputing terms + hidden memory cost of recursion!

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- 1. Store the values of  $F_0 = 1$ ,  $F_1 = 1$  in an array A.
- 2. for i = 2 to 1000

$$F_i = A[i-1] + A[i-2]$$
$$A[i] = F_i$$
end

 This is a very elementary example of a very useful technique called <u>dynamic programming</u>.

## **Closed Formula for Fibonacci**

• The closed-form formula for  $F_n$  is:

$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n$$
  
$$\phi \qquad \qquad \psi$$
  
$$F_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n$$

• Roughly:  $F_n \approx \phi^n \approx (1.618)^n$ 

# **Recursion vs Closed Formula**

#### **1**. Computation:

- Recursion leads to a fast dynamic program.
- Classic recursion is elegant.
- Closed form: faster, but numerical issues arise.
- 2. Rate of growth:
  - Recursion gives no hint as to how big  $F_n$  is.
  - Closed form yields  $F_n \approx (1.618)^n$

## **Summation Notation**

- Suppose I have some terms of a sequence, let's say  $a_1, a_2, a_3, \dots, a_k$ .
- Their sum,  $a_1 + a_2 + a_3 + \dots + a_k$  is denoted as:



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# Examples

$$\sum_{i=1}^{2} a_i = a_1 + a_2$$

$$\sum_{i=1}^{1} a_i = a_1$$

 $\sum_{i=1}^{0} a_i = ?$ 



# Examples

1

Something

Else

$$\sum_{i=1}^{2} a_i = a_1 + a_2$$
$$\sum_{i=1}^{1} a_i = a_1$$
$$\sum_{i=1}^{0} a_i = ?$$

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So what happens if we pick  $n_1 = 0$ ? Then, for this to work, it's necessary that  $\sum_{i=1}^{0} a_i = 0$ 

## **Product Notation**

• The **product**,  $a_1 \cdot a_2 \cdot ... \cdot a_k$  is denoted as:











$$\prod_{i=1}^{0} a_i = 1$$

• The following formula has to work for all choices of  $n_1 \in \mathbb{N}$ :

$$\prod_{i=1}^{n} a_i = \prod_{i=1}^{n_1} a_i \cdot \prod_{i=n_1+1}^{n} a_i$$
  
• So, for  $n_1 = 0$ , we need  $\prod_{i=1}^{0} a_i = 1$ 

# Sum / Product Notation

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## **Sum-Product Notation**

- We can have certain *exclusionary conditions* under the  $\Sigma$  and  $\Pi$  symbols.
- Examples:





## **Series and Partial Sums**

• A *series* is the **sum** of **all** elements of an **infinite** sequence.

$$\sum_{i=0}^{+\infty} a_i = a_0 + a_1 + a_2 + \cdots$$
Or 1, if we start at 1

• A **partial sum** of a sequence, denoted  $S_n$ , is the sum ranging from the first up to (and including) the  $n^{th}$  term of a (usually infinite) sequence:

$$S_n = \sum_{i=0}^{n} a_i = a_0 + a_1 + a_2 + \dots + a_n$$
  
Or 1, if we start at 1

• Arithmetic (often called the arithmetic progression):

$$a_0, a_0 + d, a_1 + d, a_2 + d \dots$$
 where  $d \in \mathbb{R}$   
 $\alpha_1 \qquad \alpha_2 \qquad \alpha_3$ 

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 Question: which among the following is the correct characterization for a<sub>n</sub>?

$$d \cdot a_{n-1}$$
  $\alpha_0 + d \cdot a_{n-1}$   $\alpha_0 + n \cdot d$   $\alpha_0 + (n-1) \cdot d$ 

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• In the arithmetic progression:

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• Should we allow d = 0?



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It will be a pretty boring sequence, but it will still be a sequence!

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# The Gauss Story



- Gauss was a great mathematician (1777-1855)
- When Gauss was in 1<sup>st</sup> grade, the class was misbehaving.
- For punishment, the teacher made everyone compute

 $1+2+\dots+100$ 

• Gauss did it in 2 minutes. Can you?

# The Gauss Trick



$$S = 1 + 2 + \dots + 100$$
  

$$S = 100 + 99 + \dots + 1$$
  

$$2S = 101 + 101 + \dots + 101$$
  
100 terms

 $\Rightarrow 2S = 101 * 100 = 10100 \Rightarrow S = 5050$ 





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2000s	1 <sup>st</sup>	1 + 2 + + 100



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	1980	3 <sup>rd</sup>	1 + 2 + + 80
	2000s	1 <sup>st</sup>	1 + 2 + + 100
re:	2020	Nursery School	1 + 2 + + 120

Our conjecture:

• Harmonic:

$$1, \frac{1}{2}, \frac{1}{3}, \frac{1}{4}, \dots$$

• **Fibonacci**:  $F_0 = F_1 = 1$  and  $\forall n \ge 2$ ,  $F_n = F_{n-1} + F_{n-2}$ 

1, 1, 2, 3, 5, 8, 13, 21, ...

# What We'll Do Next

- We will have an intro to induction.
- The following can be proven via induction:

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$
$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

# ST()P RECORDING