

**START**

**RECORDING**

# Mathematical Induction: Introduction and Basic Problems

CMSC 250

# INTRO AND BASIC SEQUENCE PROBLEMS

# The Idea Behind Induction

- Suppose that we want to prove that a proposition  $P(n)$  is true for all natural numbers  $n$ .



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  1. For  $n = 0$ ,  $P(n)$  is true (*simplifiable to “ $P(0)$  is true”*).



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  1.  $P(0)$  is true.
  2. For all  $n \geq 1$ ,  $P(n) \Rightarrow P(n + 1)$



# How We'll Make It Work

1. Inductive **base**: We will prove (explicitly, no matter how dumb it may sometimes seem) that  $P(0)$  is true.



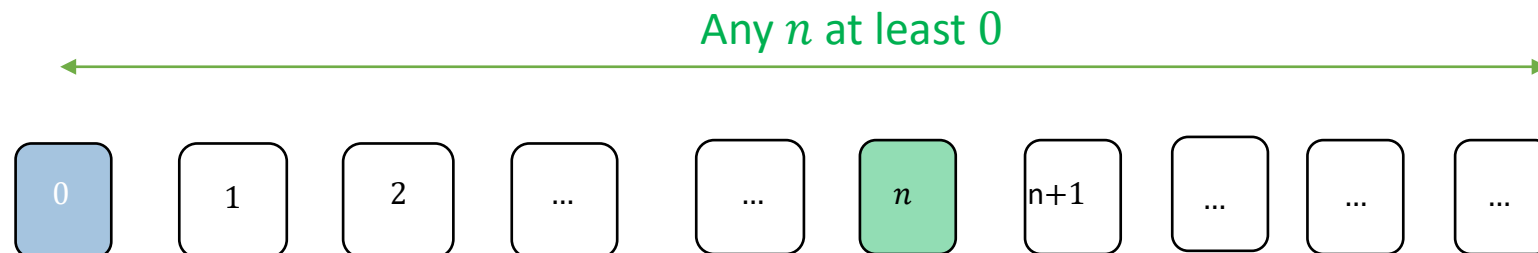
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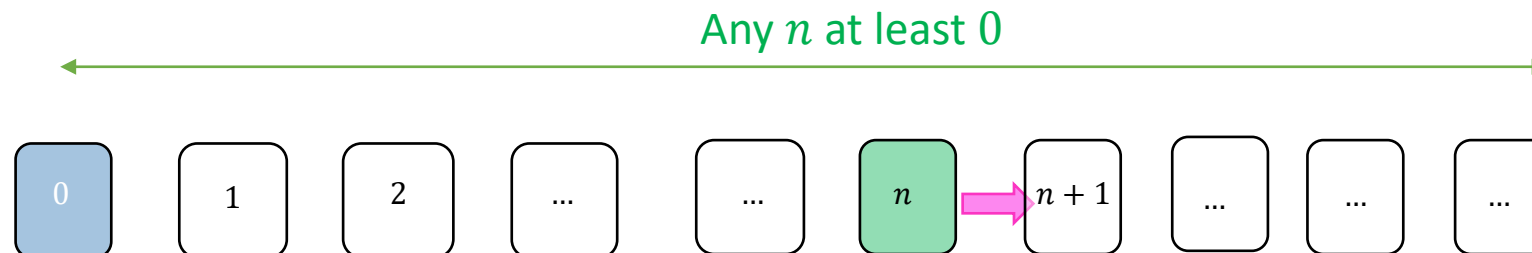
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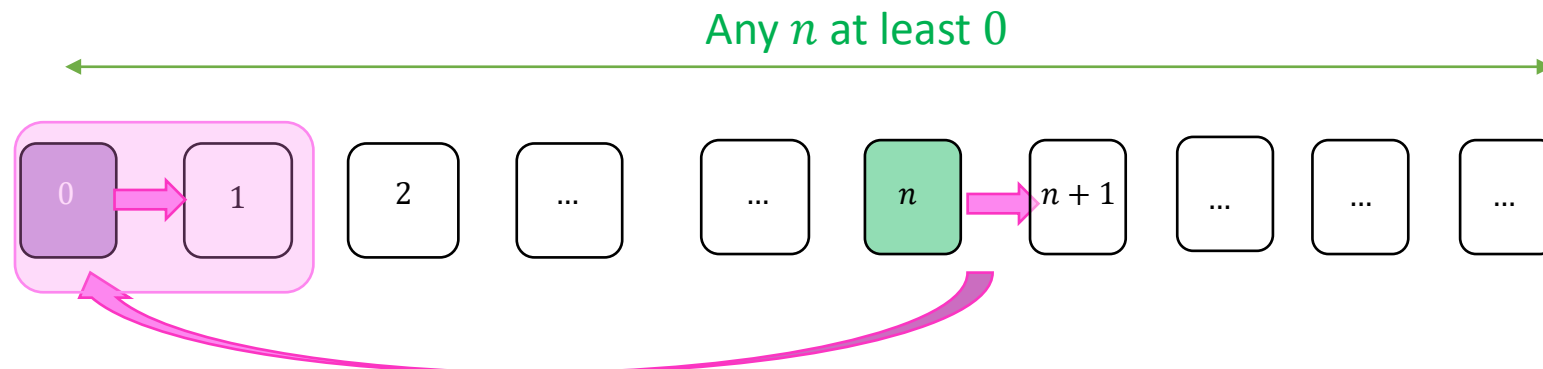
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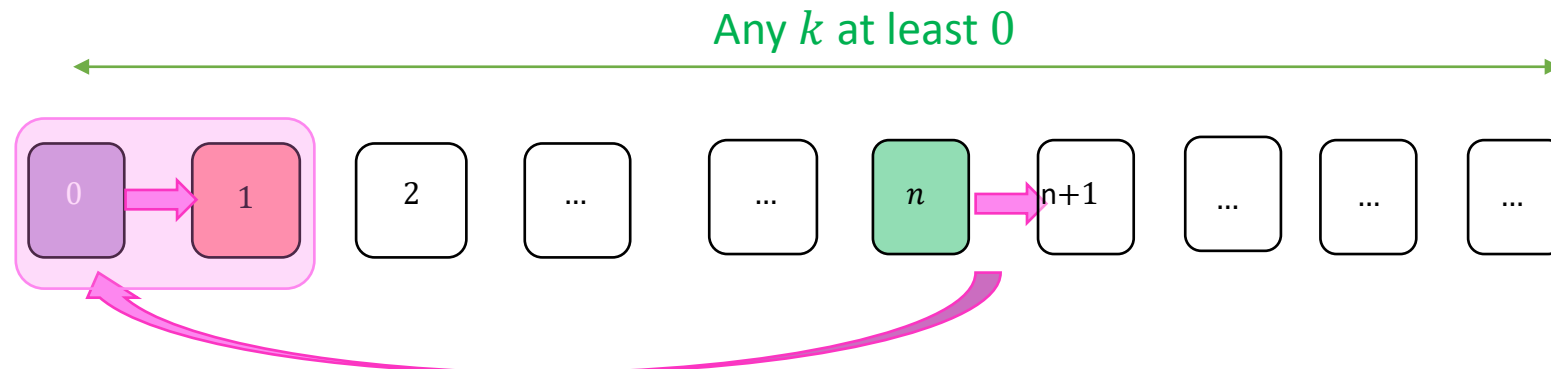
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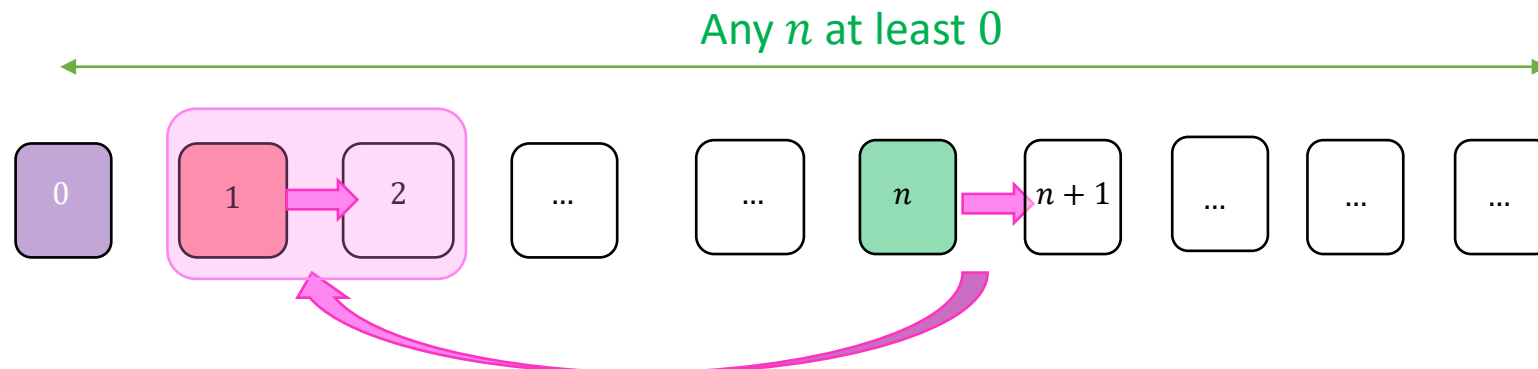
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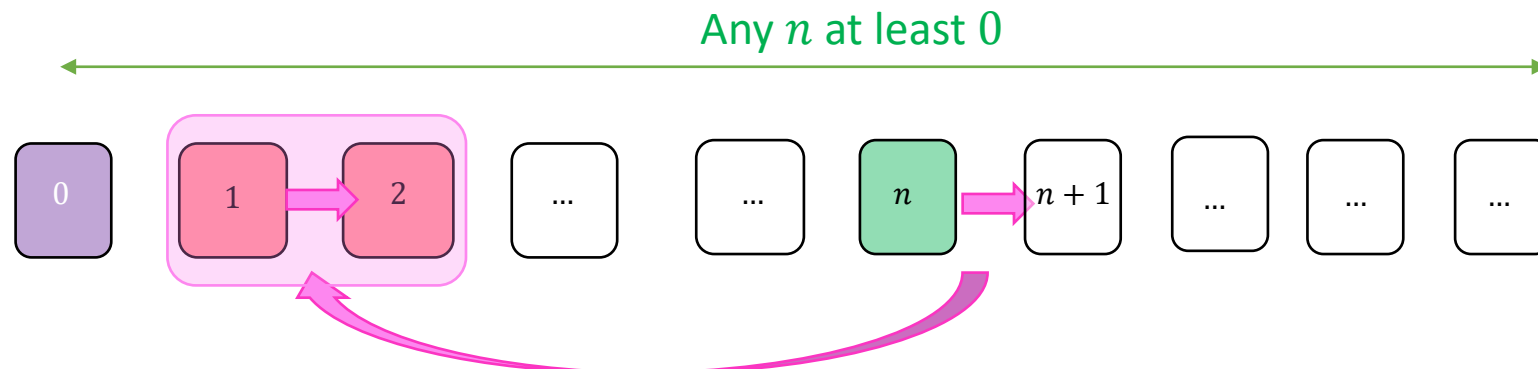
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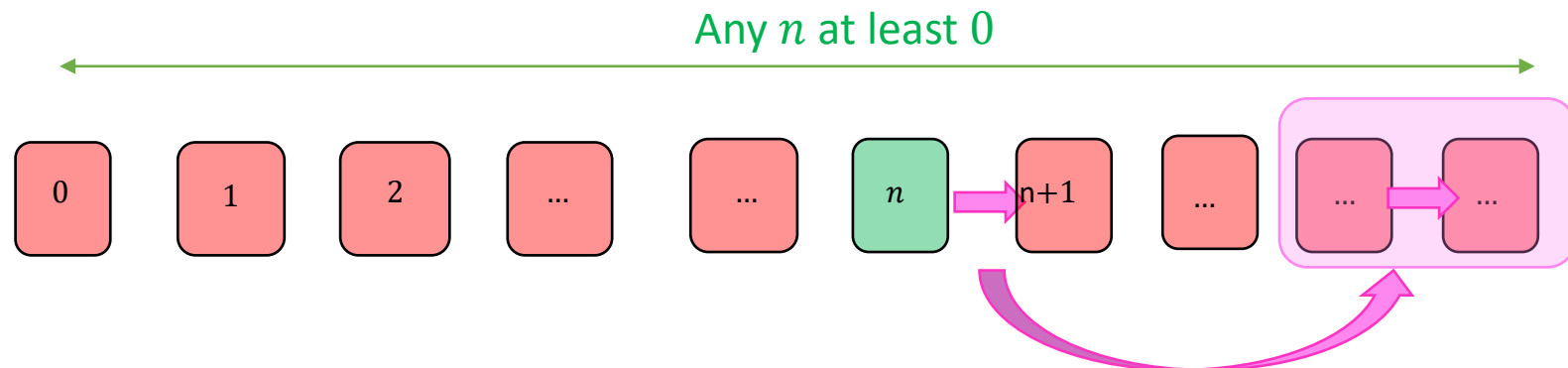
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*(We fast-forwarded here to save some time.)*

# SUM PROBLEMS

$$\sum_{i=0}^n f(n)$$

# The Gaussian Sum

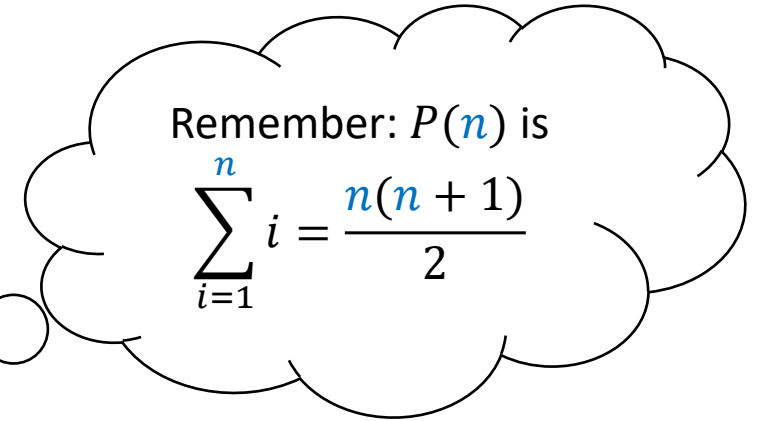
- We will prove that the sum of the first  $n$  numbers is equal to  $\frac{n(n+1)}{2}$ .
- Symbolically:

$$\underbrace{1 + 2 + 3 + \cdots + (n - 1) + n}_{\sum_{i=1}^n i} = \frac{n(n + 1)}{2}$$
$$\sum_{i=1}^n i = \frac{n(n + 1)}{2}$$

# Inductive Base

- For  $n = 0$ , we will **prove** that  $P(0)$  holds

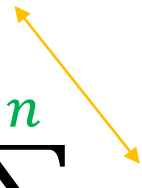
$$\sum_{i=1}^0 i = \frac{0(0+1)}{2}$$



- LHS:  $\sum_{i=1}^0 i = 0$  (recall this fact from our sequences lecture)
- RHS:  $\frac{0(0+1)}{2} = 0$
- Since LHS = RHS for  $n = 0$ ,  $P(0)$  has been proven true.

# Inductive Hypothesis

- For  $n \geq 0$ , we **assume** that  $P(n)$  **is true**:

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$


So, we **assume** that

$$P(n) \Leftrightarrow \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

is true for an arbitrary  $n \geq 0$

- Inductive Hypothesis done!

# Inductive Step

- Given that  $P(n)$  is true, we will **prove** that  $P(n + 1)$  is true.

$$\sum_{i=1}^n i = \frac{n(n+1)}{2} \Rightarrow \sum_{i=1}^{n+1} i = \frac{(n+1)(n+2)}{2}$$

# Inductive Step

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The diagram illustrates the inductive step for the sum of integers. It shows the transition from the formula for  $n$  to the formula for  $n+1$ . Two pink arrows indicate the changes: one from  $n$  to  $n+1$  and another from  $n+1$  to  $n+2$ .



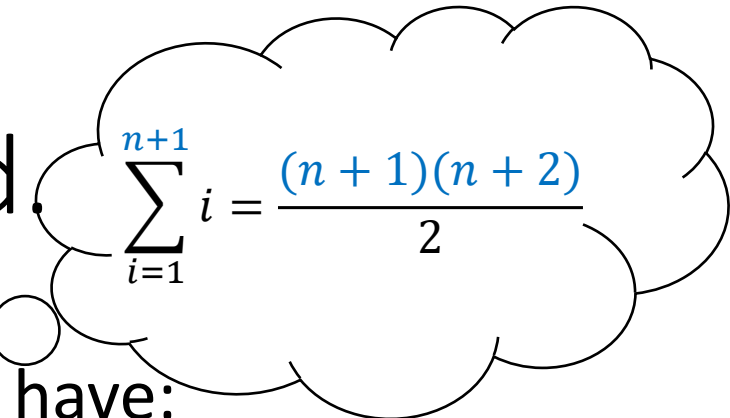
# Inductive Step

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$$\sum_{i=1}^n i = \frac{n(n+1)}{2} \Rightarrow \sum_{i=1}^{n+1} i = \frac{(n+1)(n+2)}{2}$$

**This is our goal!**

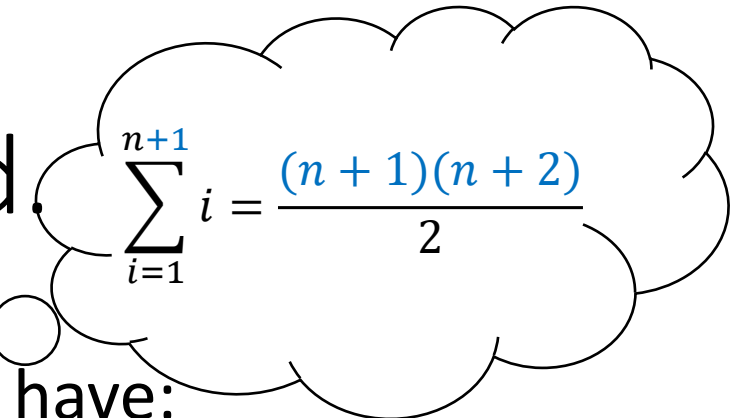
## Inductive Step, contd.


$$\sum_{i=1}^{n+1} i = \frac{(n+1)(n+2)}{2}$$

- Starting from the **LHS** of the relation to prove, we have:

$$\sum_{i=1}^{n+1} i = 1 + 2 + \cdots + n + (n + 1)$$

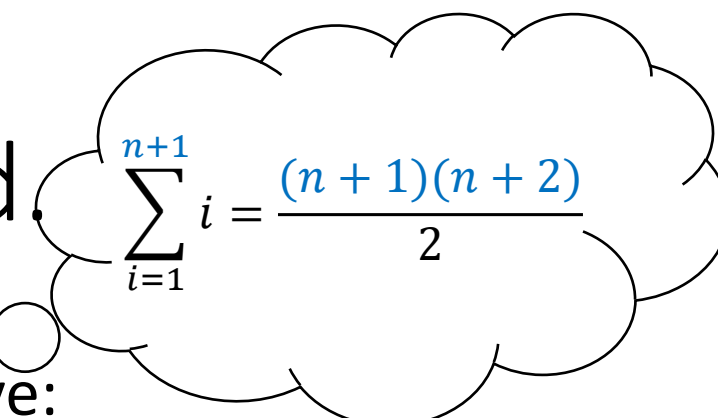
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## Inductive Step, contd.


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$$\sum_{i=1}^{n+1} i = 1 + 2 + \dots + n + (n+1) = \sum_{i=1}^n i + (n+1) \quad (1)$$

- **From the Inductive Hypothesis**, we have that

$$\sum_{i=1}^n i = \frac{n(n+1)}{2} \quad (2)$$

# Inductive Step, contd.

$$\sum_{i=1}^{n+1} i = \frac{(n+1)(n+2)}{2}$$

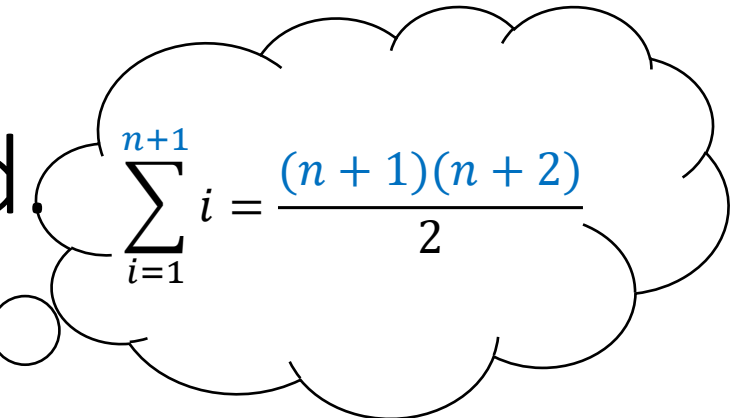
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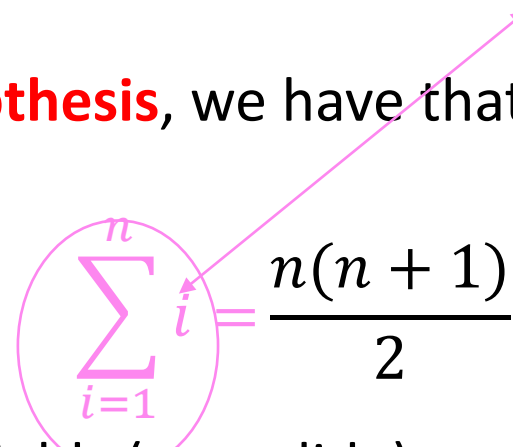
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- From the Inductive Hypothesis**, we have that


$$\sum_{i=1}^n i = \frac{n(n+1)}{2} \quad (2)$$

- Substituting (2) into (1) yields (next slide):

Inductive Step, contd.

$$\sum_{i=1}^{n+1} i = \frac{n(n+1)}{2} + (n+1) = \frac{n(n+1)}{2} + \frac{2(n+1)}{2} = \frac{(n+2)(n+1)}{2} = RHS$$

## Inductive Step, contd.

$$\sum_{i=1}^{n+1} i = \frac{n(n+1)}{2} + (n+1) = \frac{n(n+1)}{2} + \frac{2(n+1)}{2} = \frac{(n+2)(n+1)}{2} = RHS$$

- So, when  $P(n)$  is true,  $P(n+1)$  was also proven true.
- We conclude that  $P(n)$  is true  $\forall n \geq 0$ .
- WE ARE DONE.



Here's Another!

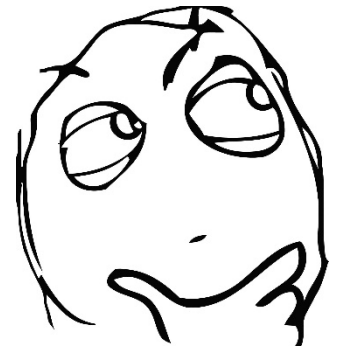
$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

# Inductive Base

- For  $n = 0$ ,  $\text{LHS} = \sum_{i=1}^0 i^2 = 0$
- $\text{RHS} = \frac{0(0+1)(2*0+1)}{2} = 0$
- Since  $\text{LHS} = \text{RHS}$ ,  $P(0)$  holds and we are done.

# Inductive Base

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- $\text{RHS} = \frac{0(0+1)(2*0+1)}{2} = 0$
- Since  $\text{LHS} = \text{RHS}$ ,  $P(0)$  holds and we are done.
- You could also start from  $n = 1$ !  $\text{LHS} = \text{RHS}$  in both cases
  - $n = 0$  sometimes makes the math easier (RHS in this case)



# Inductive Hypothesis

- Suppose that  $n \geq 0$ . (*Or 1 in the alternative scenario*)
- We will then assume  $P(n)$ , i.e:

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

# Inductive Step

- We will now attempt to prove  $P(n + 1)$ , i.e

$$\sum_{i=1}^{n+1} i^2 = \frac{(n + 1)(n + 2)(2n + 3)}{6}$$

Careful with  
factoring please!!!

# Inductive Step

- We will now attempt to prove  $P(n + 1)$ , i.e

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Careful with factoring please!!!

- By leveraging associativity of sum, the LHS can be written as follows:

$$\sum_{i=1}^{n+1} i^2 = \sum_{i=1}^n i^2 + (n + 1)^2$$

# Inductive Step

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- By leveraging associativity of sum, the LHS can be written as follows:

$$\sum_{i=1}^{n+1} i^2 = \sum_{i=1}^n i^2 + (n + 1)^2$$



We can apply the IH here!

# Inductive Step

- By IH, we can now write:

$$\sum_{i=1}^{n+1} i^2 = \frac{n(n+1)(2n+1)}{6} + (n+1)^2$$

- Remember: we **want** this to be equal to

$$\frac{(n+1)(n+2)(2n+3)}{6}$$

- We will fearlessly manipulate the algebra until it does!



## Inductive Step - Algebra

$$\begin{aligned} \frac{n(n+1)(2n+1)}{6} + (n+1)^2 &= \frac{n(n+1)(2n+1)}{6} + \frac{6(n+1)^2}{6} \\ &= \frac{(n+1)[n(2n+1) + 6(n+1)]}{6} = \frac{(n+1)[2n^2 + 7n + 6]}{6} \end{aligned}$$

- If only we could prove that  $2n^2 + 7n + 6 = (n+2)(2n+3)$ , we'd be done!
- But....  $(n+2)(2n+3) = 2n^2 + 3n + 4n + 6 = 2n^2 + 7n + 6!$  😊
- So we're done.

# Sums of Powers of 2

- Prove that the sum of the first  $n$  terms of a **geometric sequence** with  $a_1 = 1$  is equal to  $2^n - 1$ .
- Symbolically:

$$\sum_{i=0}^{n-1} 2^i = 2^n - 1$$

# Proof

- Proof : We attempt to prove  $P(n)$ ,  $\forall n \in \mathbb{N}$ . We proceed via **induction on  $n$** .
- **Inductive base:** We attempt to prove  $P(0)$ .

$$P(0): \sum_{i=0}^{1-1} 2^i = 2^0 - 1 \Leftrightarrow 0 = 0$$

So  $P(0)$  is true.

- **Inductive hypothesis:** Suppose  $n \geq 0$ . We assume  $P(n)$ , i.e

$$\sum_{i=0}^{n-1} 2^i = 2^n - 1$$

# Proof (contd.)

- **Inductive step:** We will attempt to prove  $P(n + 1)$ , i.e

$$\sum_{i=0}^{(n+1)-1} 2^i = 2^{n+1} - 1$$

From the LHS to the RHS:

$$LHS = \sum_{i=0}^n 2^i = \sum_{i=0}^{n-1} 2^i + 2^n = 2^n - 1 + 2^n = 2(2^n) - 1 = 2^{n+1} - 1 = RHS \square$$

# Sums of Powers of $m$

- Prove that the sum of the first  $n$  terms of a **geometric sequence** with  $m \in (\mathbb{R} - \{1\})$  and  $a_1 = 1$  is equal to  $\frac{m^n - 1}{m - 1}$ .
- Symbolically:

$$\sum_{i=0}^{n-1} m^i = \frac{m^n - 1}{m - 1}$$

# Proof

- Proof : We attempt to prove  $P(n)$ ,  $\forall n \in \mathbb{N}$ . We proceed via **induction on  $n$** .
- **Inductive base:** We attempt to prove  $P(0)$ .

$$P(0): \sum_{i=0}^{1-1} m^i = \frac{m^1 - 1}{m - 1} \Leftrightarrow \sum_{i=0}^0 m^i = \frac{m^1 - 1}{m - 1} \Leftrightarrow 1 = 1$$

So  $P(0)$  is true.

*Note: In the base case we are assuming  $m \neq 1$*

- **Inductive hypothesis:** Suppose  $n \geq 0$ . We assume  $P(n)$ , i.e

$$\sum_{i=0}^{n-1} m^i = \frac{m^n - 1}{m - 1}$$

# Proof (contd.)

- **Inductive step:** We will attempt to prove  $P(n + 1)$ , i.e

$$\sum_{i=0}^{(n+1)-1} m^i = \frac{m^{n+1} - 1}{m - 1}$$

From the LHS to the RHS:

$$\begin{aligned} \text{LHS} &= \sum_{i=0}^n m^i = \sum_{i=0}^{n-1} m^i + m^n = \frac{m^n - 1}{m - 1} + m^n = \frac{m - 1 + m^n(m - 1)}{m - 1} = \frac{m^{n+1} - 1}{m - 1} = \text{RHS} \quad \square \end{aligned}$$

# Base Cases

- It is standard to change your base cases to later in your index if the theorem you are trying to prove starts later



# COIN PROBLEMS!



# A Coin Problem

- We will prove that every dollar amount  $\geq 4$  cents can be exclusively paid for by 2 and/or 5 cent coins.



# Theorem Expressed in Quantifiers



- All quantifiers implicitly assumed over  $\mathbb{N}$ .

$$(\forall n \geq 4)(\exists n_1, n_2)[n = 2n_1 + 5n_2]$$

# Inductive Base



- The least amount of money we are required to prove the statement for is 4¢, so we will attempt to **prove  $P(4)$** .
- For  $n = 4$ , we have 4¢. Since  $4¢ = 2 \times 2¢$ , we are done (we have shown that the amount of 4¢ can be **exclusively** paid for by using only **2 and/or** 5 cent coins)

# Inductive Hypothesis



- Let  $n \geq 4$ .
- Assume  $P(n) \Leftrightarrow (\exists n_1, n_2)[n = 2n_1 + 5n_2]$

# Inductive Step



- We will **prove** that  $P(n) \Rightarrow P(n + 1)$ , i.e that we can pay an amount of money equal to  $n + 1$  cents using **only 2¢ or 5¢ coins**.
- In terms of algebra, what we want to prove is:

$$(\exists n_3, n_4 \in \mathbb{N}) [n + 1 = 2n_3 + 5n_4]$$

# Inductive Step



- We will **prove** that  $P(n) \Rightarrow P(n + 1)$ , i.e that we can pay an amount of money equal to  $n+1$  cents using **only 2¢ or 5¢ coins**.
- In terms of algebra, what we want to prove is:

$$(\exists n_3, n_4 \in \mathbb{N}) [n + 1 = 2n_3 + 5n_4]$$

Different variables from IH!

# Inductive Step (contd.)



- From the **Inductive Hypothesis (IH)**, we have that for some specific positive integers  $n_1$  and  $n_2$ :

$$n = 2n_1 + 5n_2$$



# Inductive Step (contd.)



- From the **Inductive Hypothesis (IH)**, we have that for some specific positive integers  $n_1$  and  $n_2$ :

$$n = 2n_1 + 5n_2$$

## 1. Case #1: $n_1 \geq 2$

- I have at least two 2¢ coins, so I can take away two 2¢ coins and add one 5¢ coin

# Inductive Step (contd.)



- From the **Inductive Hypothesis (IH)**, we have that for some specific positive integers  $n_1$  and  $n_2$ :

$$n = 2n_1 + 5n_2$$

## 1. Case #1: $n_1 \geq 2$

- I have at least two 2¢ coins, so I can take away two 2¢ coins and add one 5 ¢ coin
- By adding 1 on both sides of the IH we obtain:

$$\begin{aligned} n + 1 &= 2n_1 + 5n_2 + 1 = 2n_1 + 5n_2 + (5 - 2 * 2) = \\ &= (2n_1 - 4) + (5n_2 + 5) = 2 \underbrace{(n_1 - 2)}_{n_3} + 5 \underbrace{(n_2 + 1)}_{n_4} = 2n_3 + 5n_4 \end{aligned}$$

# Inductive Step (contd.)



- From the **Inductive Hypothesis (IH)**, we have that for some specific positive integers  $n_1$  and  $n_2$ :

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$n_1 - 2 \geq 0$  because  
 $n_1 \geq 2$

In  $\mathbb{N}$  by closure

# Inductive Step



## 2. Case #2: $n_2 \geq 1$

- I have at least one 5¢ coin so I can take away one 5¢ coin and add three 2¢ coins
- By adding 1 on both sides of the IH we obtain:

$$\begin{aligned}n + 1 &= 2n_1 + 5n_2 + 1 = 2n_1 + 5n_2 + (3 * 2 - 5) = \\ &= 2 \underbrace{(n_1 + 3)}_{n_3} + 5 \underbrace{(n_2 - 1)}_{n_4} = 2n_3 + 5n_4\end{aligned}$$

# Inductive Step



## 2. Case #2: $n_2 \geq 1$

- I have at least one 5¢ coin so I can take away one 5¢ coin and add three 2¢ coins
- By adding 1 on both sides of the IH we obtain:

$$\begin{aligned} k + 1 &= 2k_1 + 5k_2 + 1 = 2k_1 + 5k_2 + (3 * 2 - 5) = \\ &= 2(\underbrace{n_1 + 3}) + 5(\underbrace{n_2 - 1}) = 2n_3 + 5n_4 \end{aligned}$$

$(n_1 + 3) \in \mathbb{N}$   
by closure

$n_2 - 1 \geq 0$   
because  
 $n_2 \geq 1$

# Inductive Step



## 3. Case #3: $(n_1 \leq 1) \wedge (n_2 = 0)$

- This case means that we have either 0 or 2¢ at our disposal.
- But this is not possible, since we want to prove the theorem only for values  $\geq 4\text{¢}$
- So we're done.  $\square$



# A Coin Problem for You!



Prove to me that **every dollar amount  $\geq 20$  cents** can be **exclusively** paid for through combinations of **5**-cent coins and **6**-cent coins!

Go to Breakout Rooms

# TREATING INEQUALITIES



# Here's One with an Inequality!

- Prove that for all integers  $n$  at least 4,  $2^n < n!$ 
  1. **IB:** We will **prove**  $P(4) \Leftrightarrow 2^4 < 4!$  Done.
  2. **IH:** For  $n \geq 4$ , we **assume**  $P(n)$ , i.e  $2^n < n!$
  3. **IS:** We will **prove**  $P(n) \Rightarrow P(n + 1)$ , i.e

$$(2^n < n!) \Rightarrow (2^{n+1} < (n + 1)!)$$

# Inductive Step...

- Prove that for all integers  $n$  at least 4,  $2^n < n!$ 
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•  $(2) \stackrel{(3)}{\Rightarrow} 2^n \cdot 2 < (n+1)! \stackrel{(1)}{\Leftrightarrow} 2^{n+1} < (n+1)!$

**STOP**

**RECORDING**