

# The Muffin Problem

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# How it Began

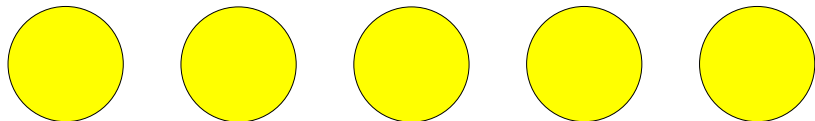
## A Recreational Math Conference (Gathering for Gardner) May 2016

I found a pamphlet:

### The Julia Robinson Mathematics Festival: A Sample of Mathematical Puzzles Compiled by Nancy Blachman

which had this problem, proposed by Alan Frank:

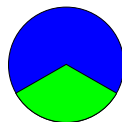
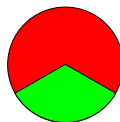
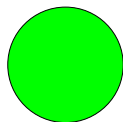
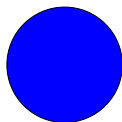
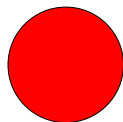
*How can you divide and distribute 5 muffins to 3 students so that every student gets  $\frac{5}{3}$  where nobody gets a tiny sliver?*



# Five Muffins, Three Students, Proc by Picture

Person	Color	What they Get
Alice	RED	$1 + \frac{2}{3} = \frac{5}{3}$
Bob	BLUE	$1 + \frac{2}{3} = \frac{5}{3}$
Carol	GREEN	$1 + \frac{1}{3} + \frac{1}{3} = \frac{5}{3}$

Smallest Piece:  $\frac{1}{3}$



# Can We Do Better?

The smallest piece in the above solution is  $\frac{1}{3}$ .

**Is there a procedure with a larger smallest piece?**

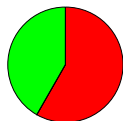
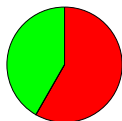
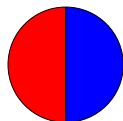
**Work on it with your neighbor**

# Five Muffins, Three People—Proc by Picture

YES WE CAN!

Person	Color	What they Get
Alice	RED	$\frac{6}{12} + \frac{7}{12} + \frac{7}{12}$
Bob	BLUE	$\frac{6}{12} + \frac{7}{12} + \frac{7}{12}$
Carol	GREEN	$\frac{5}{12} + \frac{5}{12} + \frac{5}{12} + \frac{5}{12}$

Smallest Piece:  $\frac{5}{12}$



# Can We Do Better?

The smallest piece in the above solution is  $\frac{5}{12}$ .

**Is there a procedure with a larger smallest piece?**

**Work on it with your neighbor**

## 5 Muffins, 3 People—Can't Do Better Than $\frac{5}{12}$

### NO WE CAN'T!

There is a procedure for 5 muffins, 3 students where each student gets  $\frac{5}{3}$  muffins, smallest piece  $N$ . We want  $N \leq \frac{5}{12}$ .

**Case 0:** Some muffin is uncut. Cut it  $(\frac{1}{2}, \frac{1}{2})$  and give both  $\frac{1}{2}$ -sized pieces to whoever got the uncut muffin. (Note  $\frac{1}{2} > \frac{5}{12}$ .) Reduces to other cases.

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(**Henceforth:** All muffins are cut into  $\geq 2$  pieces.)

**Case 1:** Some muffin is cut into  $\geq 3$  pieces. Then  $N \leq \frac{1}{3} < \frac{5}{12}$ .



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**Case 1:** Some muffin is cut into  $\geq 3$  pieces. Then  $N \leq \frac{1}{3} < \frac{5}{12}$ .

(**Henceforth:** All muffins are cut into 2 pieces.)

**Case 2:** All muffins are cut into 2 pieces. 10 pieces, 3 students:

**Someone** gets  $\geq 4$  pieces. He has some piece

$$\leq \frac{5}{3} \times \frac{1}{4} = \frac{5}{12} \quad \text{Great to see } \frac{5}{12}$$

# What Happened Next?

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Yada Yada Yada- in 2020:

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## MATHEMATICAL MUFFIN MORSELS: NOBODY WANTS A SMALL PIECE

William Gasarch, Erik Metz, Jacob Prinz, Daniel Smolyak  
University of Maryland, USA

In this book we consider THE MUFFIN PROBLEM: what is the best way to divide up  $m$  muffins for  $s$  students so that everyone gets  $m/s$  muffins, with the smallest pieces maximized.

This problem takes us through much mathematics of interest, for example, combinatorics and optimization theory.

228pp

978-981-121-597-1(pbk)

978-981-121-517-9

978-981-121-519-3(mbook)

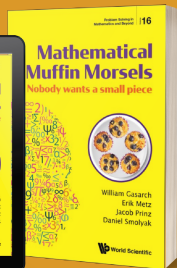
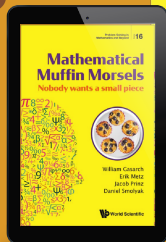
US\$28 / £25 / SGD41

US\$58 / £50 / SGD86

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*Is there a way to divide five muffins for three students so that everyone gets  $5/3$ , and all pieces are larger than  $1/3$ ?*

*Spoiler alert: Yes!*



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<https://doi.org/10.1142/11689>

 World Scientific

# General Problem

$f(m, s)$  be the smallest piece in the best procedure (best in that the smallest piece is maximized) to divide  $m$  muffins among  $s$  students so that everyone gets  $\frac{m}{s}$ .

We have shown  $f(5, 3) = \frac{5}{12}$  here.

We have two proofs that shown  $f(m, s)$  exists, is rational, and is computable.

One use Linear Programming.

One use Integer Programming.

# Amazing Results! / Amazing Theorems!

1.  $f(43, 33) = \frac{91}{264}$ .
2.  $f(52, 11) = \frac{83}{176}$ .
3.  $f(35, 13) = \frac{64}{143}$ .

**All done by hand, no use of a computer**  
**by Co-author Erik Metz is a muffin savant !**

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Have **General Theorems** from which **upper bounds** follow.  
Have **General Procedures** from which **lower bounds** follow.

What if  $m < s$ ?



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**Duality Theorem:**  $f(m, s) = \frac{m}{s} f(s, m)$ .

## What if $m < s$ ?

**Duality Theorem:**  $f(m, s) = \frac{m}{s}f(s, m)$ .

Hence we will just look at  $m > s$ .

## Floor-Ceiling Thm Generalizes $f(5, 3) \leq \frac{5}{12}$

$$f(m, s) \leq \text{FC}(m, s) = \max\left\{\frac{1}{3}, \min\left\{\frac{m}{s \lceil 2m/s \rceil}, 1 - \frac{m}{s \lfloor 2m/s \rfloor}\right\}\right\}.$$

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**Someone** gets  $\geq \lceil \frac{2m}{s} \rceil$  pieces.  $\exists$  piece  $\leq \frac{m}{s} \times \frac{1}{\lceil 2m/s \rceil} = \frac{m}{s \lceil 2m/s \rceil}$ .

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**Someone** gets  $\leq \lfloor \frac{2m}{s} \rfloor$  pieces.  $\exists$  piece  $\geq \frac{m}{s} \frac{1}{\lfloor 2m/s \rfloor} = \frac{m}{s \lfloor 2m/s \rfloor}$ .

The other piece from that muffin is of size  $\leq 1 - \frac{m}{s \lfloor 2m/s \rfloor}$ .

## FC Gives Upper Bound

Give  $m, s$ :

$$\text{FC}(m, s) = \max \left\{ \frac{1}{3}, \min \left\{ \frac{m}{s \lceil 2m/s \rceil}, 1 - \frac{m}{s \lfloor 2m/s \rfloor} \right\} \right\}$$

Gives an upper bound. So we know

$$(\forall m, s)[f(m, s) \leq \text{FC}(m, s)].$$

Is the following true?

$$(\forall m, s)[f(m, s) = \text{FC}(m, s)]$$

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Is the following true?

$$(\forall m, s)[f(m, s) = \text{FC}(m, s)]$$

**No:** If so my book would be about 20 pages.

# THREE Students

**CLEVERNESS, COMP PROGS** for the procedure.

**FC Theorem** for optimality.

$$f(1, 3) = \frac{1}{3}$$

$$f(3k, 3) = 1.$$

$$f(3k + 1, 3) = \frac{3k-1}{6k}, k \geq 1.$$

$$f(3k + 2, 3) = \frac{3k+2}{6k+6}.$$

**Note:** A Mod 3 Pattern.

**Theorem:** For all  $m \geq 3$ ,  $f(m, 3) = \text{FC}(m, 3)$ .

# FOUR Students

**CLEVERNESS, COMP PROGS** for procedures.

**FC Theorem** for optimality.

$$f(4k, 4) = 1 \text{ (easy)}$$

$$f(1, 4) = \frac{1}{4} \text{ (easy)}$$

$$f(4k + 1, 4) = \frac{4k-1}{8k}, k \geq 1.$$

$$f(4k + 2, 4) = \frac{1}{2}.$$

$$f(4k + 3, 4) = \frac{4k+1}{8k+4}.$$

**Note:** A Mod 4 Pattern.

**Theorem:** For all  $m \geq 4$ ,  $f(m, 4) = \text{FC}(m, 4)$ .

# FIVE Students

**CLEVERNESS, COMP PROGS** for procedures.

**FC Theorem** for optimality.

For  $k \geq 1$ ,  $f(5k, 5) = 1$ .

For  $k = 1$  and  $k \geq 3$ ,  $f(5k + 1, 5) = \frac{5k+1}{10k+5}$ .  $f(11, 5)$ ?

For  $k \geq 2$ ,  $f(5k + 2, 5) = \frac{5k-2}{10k}$ .  $f(7, 5) = \text{FC}(7, 5) = \frac{1}{3}$

For  $k \geq 1$ ,  $f(5k + 3, 5) = \frac{5k+3}{10k+10}$

For  $k \geq 1$ ,  $f(5k + 4, 5) = \frac{5k+1}{10k+5}$

**Note:** A Mod 5 Pattern.

**Theorem:** For all  $m \geq 5$  **except  $m=11$** ,  $f(m, 5) = \text{FC}(m, 5)$ .

# What About FIVE students, ELEVEN muffins?

1. We have a procedure which shows  $f(11, 5) \geq \frac{13}{30}$ .
2.  $f(11, 5) \leq \max\{\frac{1}{3}, \min\{\frac{11}{5\lceil 22/5 \rceil}, 1 - \frac{11}{5\lceil 22/5 \rceil}\}\} = \frac{11}{25}$ .

So

$$\frac{13}{30} \leq f(11, 5) \leq \frac{11}{25} \quad \text{Diff} = 0.006666\dots$$

Options:

1.  $f(11, 5) = \frac{11}{25}$ . Need to find procedure.
2.  $f(11, 5) = \frac{13}{30}$ . Need to find new technique for upper bounds.
3.  $f(11, 5)$  in between. Need to find both.
4.  $f(11, 5)$  unknown to science!

**Vote**

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**Vote WE SHOW**  $f(11, 5) = \frac{13}{30}$ . **Exciting** new technique!



## Terminology: Buddy

Assume that in some protocol every muffin is cut into two pieces.

Let  $x$  be a piece from muffin  $M$ .

The *other piece* from muffin  $M$  is the **buddy of  $x$** .

Note that the **buddy** of  $x$  is of size

$$1 - x.$$

## $f(11, 5) = \frac{13}{30}$ , Easy Case Based on Muffins

There is a procedure for 11 muffins, 5 students where each student gets  $\frac{11}{5}$  muffins, smallest piece  $N$ . We want  $N \leq \frac{13}{30}$ .

**Case 0:** Some muffin is uncut. Cut it  $(\frac{1}{2}, \frac{1}{2})$  and give both halves to whoever got the uncut muffin. Reduces to other cases.

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(**Negation of Case 0 and Case 1:** All muffins cut into 2 pieces.)

## $f(11, 5) = \frac{13}{30}$ , Easy Case Based on Students

**Case 2:** Some student gets  $\geq 6$  pieces.

$$N \leq \frac{11}{5} \times \frac{1}{6} = \frac{11}{30} < \frac{13}{30}.$$

**Case 3:** Some student gets  $\leq 3$  pieces.

One of the pieces is

$$\geq \frac{11}{5} \times \frac{1}{3} = \frac{11}{15}.$$

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One of the pieces is

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That piece **buddy** is of size:

$$\leq 1 - \frac{11}{15} = \frac{4}{15} < \frac{13}{30}.$$

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That piece **buddy** is of size:

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**(Negation of Cases 2 and 3:** Every student gets 4 or 5 pieces.)

## $f(11, 5) = \frac{13}{30}$ , Fun Cases

**Case 4:** Every muffin is cut in 2 pieces, every student gets 4 or 5 pieces. Number of pieces: 22. Note  $\leq 11$  pieces are  $> \frac{1}{2}$ .



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- ▶  $s_4$  is number of students who get 4 pieces
- ▶  $s_5$  is number of students who get 5 pieces

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$$4s_4 + 5s_5 = 22$$

$$s_4 + s_5 = 5$$

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$$4s_4 + 5s_5 = 22$$

$$s_4 + s_5 = 5$$

$s_4 = 3$ : There are 3 students who have 4 shares.

$s_5 = 2$ : There are 2 students who have 5 shares.

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$$4s_4 + 5s_5 = 22$$

$$s_4 + s_5 = 5$$

$s_4 = 3$ : There are 3 students who have 4 shares.

$s_5 = 2$ : There are 2 students who have 5 shares.

We call a share that goes to a person who gets 4 shares a **4-share**.

We call a share that goes to a person who gets 5 shares a **5-share**.

$f(11, 5) = \frac{13}{30}$ , Fun Cases

**Case 4.1:** Some 4-share is  $\leq \frac{1}{2}$ .

## $f(11, 5) = \frac{13}{30}$ , Fun Cases

**Case 4.1:** Some 4-share is  $\leq \frac{1}{2}$ .

Alice gets  $w, x, y, z$  and  $w \leq \frac{1}{2}$ .

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GREAT! This is where  $\frac{13}{30}$  comes from!

## $f(11, 5) = \frac{13}{30}$ , Fun Cases

**Case 4.2:** All 4-shares are  $> \frac{1}{2}$ . There are  $4s_4 = 12$  4-shares.  
There are  $\geq 12$  pieces  $> \frac{1}{2}$ . Can't occur.

## HALF Method

The above reasoning can be used to *verify* that  $f(11, 5) \leq \frac{13}{30}$  but could not *generate* the upper bound  $\frac{13}{30}$ .

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For  $f(24, 11)$  it fails!

$$f(24, 11) \leq \frac{19}{44}$$

Assume  $(24, 11)$ -procedure with smallest piece  $> \frac{19}{44}$ .

Can assume all muffin cut in two and all student gets  $\geq 2$  shares.

We show that there is a piece  $\leq \frac{19}{44}$ .



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**Case 1:** A student gets  $\geq 6$  shares. Some piece  $\leq \frac{24}{11 \times 6} < \frac{19}{44}$ .

**Case 2:** A student gets  $\leq 3$  shares. Some piece  $\geq \frac{24}{11 \times 3} = \frac{8}{11}$ .

Buddy of that piece  $\leq 1 - \frac{8}{11} \leq \frac{3}{11} < \frac{19}{44}$ .

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**Case 3:** Every muffin is cut in 2 pieces and every student gets either 4 or 5 shares. Total number of shares is 48.

## How many students get 4? 5? Where are the Shares?

*4-students*: a student who gets 4 shares.  $s_4$  is the number of them.

*5-students*: a student who gets 5 shares.  $s_5$  is the number of them.

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$$4s_4 + 5s_5 = 48$$

$$s_4 + s_5 = 11$$

$s_4 = 7$ . Hence there are  $4s_4 = 4 \times 7 = 28$  4-shares.

$s_5 = 4$ . Hence there are  $5s_5 = 5 \times 4 = 20$  5-shares.

## Case 3.1 and 3.2: Too Big or Too Small

**Case 3.1:**  $\exists$  a share  $\geq \frac{25}{44}$ . Its **buddy** is

$$\leq 1 - \frac{25}{44} = \frac{19}{44}$$



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Henceforth assume that all shares are in

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**Claim:** If some 5-share is  $\geq \frac{20}{44}$  then some share  $\leq \frac{19}{44}$ .

**Proof:** Assume that Alice 5 pieces  $A, B, C, D, E$  and  $E \geq \frac{20}{44}$ .  
Since  $A + B + C + D + E = \frac{24}{11}$  and  $E > \frac{20}{44}$

$$A + B + C + D < \frac{24}{11} - \frac{20}{44} = \frac{76}{44}$$

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Assume  $A$  is the smallest of  $A, B, C, D$ .

$$A \leq \frac{76}{44} \times \frac{1}{4} = \frac{19}{44}$$

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**Proof:** Assume that Alice 4 pieces  $A, B, C, D$  and  $D \leq \frac{21}{44}$ .

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Assume  $A$  is the largest of  $A, B, C$ .

$$A \geq \frac{75}{44} \times \frac{1}{3} = \frac{25}{44}$$

The **buddy** of  $A$  is of size

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**Case 3.5:** 4-shares in  $(\frac{21}{44}, \frac{25}{44})$ , 5-shares in  $(\frac{19}{44}, \frac{20}{44})$ .

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**Recall:** there are  $4s_4 = 4 \times 7 = 28$  4-shares.

**Recall:** there are  $5s_5 = 5 \times 4 = 20$  5-shares.

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## More Refined Picture of What is Going On

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \text{ shs} \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 28 \text{ 4-shs} \\ \frac{21}{44} \end{array} \right) \left( \begin{array}{c} \\ \frac{25}{44} \end{array} \right)$$

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The following picture captures what we know so far.

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 8 \text{ S4-shs} \\ \frac{21}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{23}{44} \end{array} \right] \left( \begin{array}{c} 20 \text{ L4-shs} \\ \frac{24}{44} \end{array} \right) \left[ \begin{array}{c} \\ \frac{25}{44} \end{array} \right)$$

S4= Small 4-shares

L4= Large 4-shares. L4 shares, 5-share: **buddies**, so  $|L4|=20$ .

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 8 \text{ S4-shs} \\ \frac{21}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{23}{44} \end{array} \right] \left( \begin{array}{c} 20 \text{ L4-shs} \\ \frac{24}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{25}{44} \end{array} \right]$$

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**Claim 2:** Every 4-student has at least 3 L4 shares.

$$\binom{20}{\frac{19}{44}} \binom{5\text{-shs}}{\frac{20}{44}} \binom{0}{\frac{21}{44}} \binom{8}{\frac{23}{44}} \binom{S4\text{-shs}}{\frac{24}{44}} \binom{0}{\frac{24}{44}} \binom{20}{\frac{25}{44}} \binom{L4\text{-shs}}{\frac{44}{44}}$$

**Claim 2:** Every 4-student has at least 3 L4 shares.

If a 4-student had  $\leq 2$  L4 shares then he has

$$< 2 \times \binom{23}{44} + 2 \times \binom{25}{44} = \frac{24}{11}.$$

$$\binom{20}{\frac{19}{44}} \binom{5\text{-shs}}{\frac{20}{44}} \binom{0}{\frac{21}{44}} \binom{8\text{ S4-shs}}{\frac{23}{44}} \binom{0}{\frac{24}{44}} \binom{20\text{ L4-shs}}{\frac{25}{44}}$$

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**Contradiction:** Each 4-student gets  $\geq 3$  L4 shares. There are  $s_4 = 7$  4-students. Hence there are  $\geq 21$  L4-shares. But there are only 20.

## More Techniques

We have shown you FC and HALF and (maybe) INT.

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We developed more:

MID, GAP,

EBM (Easy Buddy Match), HBM (Hard Buddy Match),

TRAIN

# Upshot

Let

$$A = \{(m, s) \mid 2 \leq s \leq 100 \text{ and } s < m \leq 110 \text{ and } m, s \text{ rel prime}\}$$

There are 3520 pairs  $(m, s)$  in  $A$ . We solved **all** of them!

- ▶ For 2301 of them  $f(m, s) = \text{FC}(m, s)$ . That is  $\sim 65.37\%$ .
- ▶ For 329 of them  $f(m, s) = \text{HALF}(m, s)$ . That is  $\sim 9.35\%$ .
- ▶ For 186 of them  $f(m, s) = \text{INT}(m, s)$ . That is  $\sim 5.28\%$ .
- ▶ For 111 of them  $f(m, s) = \text{MID}(m, s)$ . That is  $\sim 3.15\%$ .
- ▶ For 240 of them  $f(m, s) = \text{EBM}(m, s)$ . That is  $\sim 6.28\%$ .
- ▶ For 89 of them  $f(m, s) = \text{HBM}(m, s)$ . That is  $\sim 2.53\%$ .
- ▶ For 250 of them  $f(m, s) = \text{GAP}(m, s)$ . That is  $\sim 7.10\%$ .
- ▶ For 13 of them  $f(m, s) = \text{TRAIN}(m, s)$ . That is  $\sim 0.40\%$

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No. Did not work on

- ▶  $f(205, 178)$
- ▶  $f(226, 135)$
- ▶  $f(233, 141)$



# The Scott Huddleston Technique

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Richard Chatwin independently came up with the same algorithm and proved it worked, but never coded it up. He tells me its poly time and I believe this can be proved, though its not in his paper. His paper is here: <https://arxiv.org/abs/1907.08726>

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These are **not** well defined terms but

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4. Not Interesting and Not Important: So far  $R(5)$  has not lead to any math of interest and is also not important to find. Same for most Ramsey-type Numbers.

# Lessons Learned

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