

HW 09 Solutions

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We do the Problems in a Funny Order

We do Problem 3 then 1 then 2.

This is in order of how interesting they are.

Problem 3a

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Here is M :

1. Input(x) (this will be ignored).
2. Run $M_1(0), \dots, M_{100}(0)$ at the same time.
3. If you see that 17 of them halted then STOP

$M(0)$ halts IFF 17 of the $M_1(0), \dots, M_{100}(0)$ halt.

Note Can replace 17 with anything and get similar result.

Problem 3b

Bill gives you 100 Turing Machines M_1, \dots, M_{100} . He wants to know HOW MANY halt on 0.

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Do binary search: First ask if there exists ≥ 50 of M_1, \dots, M_{100} that halt on 0. If yes then ask ≥ 75 . If No then ask ≥ 25 . Etc. This takes $\lceil \log_2(100) \rceil = 7$.

Problem 3c

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First find HOW MANY halt on 0 by 3b. That is only 8 questions. You now KNOW x : EXACTLY x of M_1, \dots, M_{100} halt on 0. What to do now? Discuss?

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RUN $M_1(0), \dots, M_{100}(0)$ simul UNTIL x of them halt. Those x halt. Great

Key The rest DO NOT HALT on 0 since exactly x halt on 0.

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There is an entire field called [Bounded Queries in Computability Theory](#)

<https://www.amazon.com/Bounded-Queries-Recursion-Progress-Computer/dp/1461268486>

Problem 1: $3\text{COL} \leq \text{CNF-SAT}$

SOLUTION

We are given a graph $G = (V, E)$. We assume $V = \{1, \dots, n\}$.

For every vertex i we have 3 Boolean variables. We list them and what they mean

x_{iR} : T if $\text{COL}(i) = R$.

x_{iB} : T if $\text{COL}(i) = B$.

x_{iG} : T if $\text{COL}(i) = G$.

Formula has two parts

3COL \leq CNF-SAT, PART ONE

Making sure that a satisfying assignment really is a (not necessarily proper) coloring

Every vertex has at least one color:

$$\bigwedge_{i=1}^n (x_{iR} \vee x_{iB} \vee x_{iG})$$

Every vertex has at most one color:

$$\bigwedge_{i=1}^n \neg(x_{iR} \wedge x_{iB}) \wedge \neg(x_{iR} \wedge x_{iG}) \wedge \neg(x_{iB} \wedge x_{iG})$$

3COL \leq CNF-SAT, PART TWO

Make sure it's a proper coloring

$$\bigwedge_{(i,j) \in E} \neg(x_{iR} \wedge x_{jR}) \wedge \neg(x_{iB} \wedge x_{jB}) \wedge \neg(x_{iG} \wedge x_{jG})$$

Prob 2a: Coding TMs into Numbers

All TMs: $\Sigma = \{1, 2, 3\}$, Q is $\{1, \dots, n\}$.

Describe a procedure to code Turing Machines into \mathbb{N} such that the following holds:

- ▶ Two diff Turing Machines map to diff numbers. (Some numbers do not get mapped to.)
- ▶ The following should be computable:
Input: $x, y \in \mathbb{N}$
Output:
If x does not code a TM than output THATSBSMAN.

HINT Do not over think this. Any way you code a TM into numbers should work.

Prob 2a: Coding TMs into Numbers, SOLUTION

THE CODING: Let $M = (Q, \{a, b, \#\}, \delta, s, h)$

The number will be the product of the following numbers

1. $2^{|Q|}$.
2. 3^s (Recall that s , the start state, is a number)
3. 5^h (Recall that h , the halt state, is a number)
4. For coding the transitions, next slide

Prob 2a: Coding TMs into Numbers, SOLUTION

There will be $n = (Q - 1) \times \Sigma$ rules. Let p_1, \dots, p_n be the n primes after 5 (so $p_1 = 7$). (It's $Q - 1$ since there are no transitions out of h .) Order the rules lexicographically by $Q \times \Sigma$, so

$\delta(1, 1)$

$\delta(1, 2)$

$\delta(1, 3)$

$\delta(2, 1)$

\vdots

$\delta(|Q| - 1, 3)$.

For $1 \leq i \leq n$ take rule i and form the following number.

1. $\delta(p, \sigma) = (q, \sigma')$ maps to $2^p \times 3^\sigma \times 5^q \times 7^{\sigma'}$. ($\sigma' \in \{1, 2, 3\}$).
2. $\delta(p, \sigma) = (q, L)$ maps to $2^p \times 3^\sigma \times 5^q \times 7^4$. ($4 \notin \{1, 2, 3\}$).
3. $\delta(p, \sigma) = (q, R)$ maps to $2^p \times 3^\sigma \times 5^q \times 7^5$. ($5 \notin \{1, 2, 3\}$).

Problem 2c

Let M be the TM: $Q = \{1, 2, 3\}$, $\Sigma = \{1, 2, 3\}$, $s = 1$, $h = 3$,

$$\delta(1, 1) = (1, L).$$

$$\delta(1, 2) = (1, 2).$$

$$\delta(1, 3) = (2, R).$$

$$\delta(2, 1) = (1, 1).$$

$$\delta(2, 2) = (3, 3).$$

$$\delta(2, 3) = (3, L).$$

Code this TM into a number using your procedure.

Problem 2c Solution

The number will be the product of many numbers:

$$Q = \{1, 2, 3\} \text{ (so } 2^3\text{),}$$

$$\Sigma = \{1, 2, 3\},$$

$$s = 1 \text{ (so } 3^1\text{),}$$

$$h = 3 \text{ (so } 5^3\text{).}$$

$$\delta(1, 1) = (1, L). \text{ This is coded by } 7^{2^1 3^1 5^1 7^4}$$

$$\delta(1, 2) = (1, 2). \text{ This is coded by } 11^{2^1 3^2 5^1 7^2}$$

$$\delta(1, 3) = (2, R). \text{ This is coded by } 13^{2^1 3^3 5^2 7^5}$$

$$\delta(2, 1) = (4, 1). \text{ This is coded by } 17^{2^2 3^1 5^4 7^1}$$

$$\delta(2, 2) = (3, 3). \text{ This is coded by } 23^{2^2 3^2 5^3 7^3}$$

$$\delta(2, 3) = (3, L). \text{ This is coded by } 29^{2^2 3^3 5^3 7^4}$$

Take the product of all of the above numbers.