

BILL AND NATHAN START RECORDING

CFL \subset P

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We will obtain $p(n) = O(n^3)$.

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We will assume that $e \notin L$ and hence we do not have the rule

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Our proof can be modified to accommodate this case.

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We are really asking: Is $S \in \text{GEN}[1, n]$?

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I will solve a harder problem:

Find $\text{GEN}[1, n]$

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Why solve this harder problem?

We will use Dynamic programming so having some $\text{GEN}[i, j]$ solved will help us solve later ones.

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The only way for $A \Rightarrow \sigma_i \sigma_{i+1}$ is if

$A \rightarrow BC$

$B \rightarrow \sigma_i$ (AH- then $B \in \text{GEN}[i, i]$.)

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$$\text{GEN}[i, i + 1] =$$

$$\{A : A \rightarrow BC \quad \wedge \quad B \in \text{GEN}[i, i] \quad \wedge \quad C \in \text{GEN}[i + 1, i + 1]\}$$

GEN[$i, i + 2$]

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We can write GEN[$i, i + 2$] in terms of GEN[i, i], GEN[$i, i + 1$],
GEN[$i + 1, i + 2$].

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We can write GEN[$i, i + 2$] in terms of GEN[i, i], GEN[$i, i + 1$], GEN[$i + 1, i + 2$].

More important: we can easily FIND GEN[$i, i + 2$] if we KNOW GEN[i, i], GEN[$i, i + 1$], GEN[$i + 1, i + 2$].

GEN[$i, i + k$]

We find GEN[$i, i + k$].

The only way for $A \Rightarrow \sigma_i \cdots \sigma_{i+k}$ is if there exists j such that

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Can use this recurrence bottom up to get a DYN PROGRAM for the problem

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- 2) For $k = 1$ to $n - 1$ we will look at $\text{GEN}[i, i + k]$
For $i = 1$ to $n - k$
 $\text{GEN}[i, i + k] =$

$$\bigcup_{i < j < k} \{A : A \rightarrow BC \quad \wedge \quad B \in \text{GEN}[i, j] \quad \wedge \quad C \in \text{GEN}[j + 1, k]\}$$

- 3) If $S \in \text{GEN}[1, n]$ then output YES, else output NO.