# BILL, RECORD LECTURE!!!!

#### BILL RECORD LECTURE!!!



# Tricks for Divisibility and DFA's

For this Slide Packet  $\Sigma = \{0, \ldots, 9\}$ 

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Strings are numbers in base 10. The string

 $d_n \cdots d_0$ 

is the number

$$d_n \times 10^n + \cdots + d_1 \times 10^1 + d_0 \times 10^0.$$

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We feed a number into a DFA right-to-left:  $d_0$ , then  $d_1$  etc.

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 $d_n \times 10^n + \cdots + d_1 \times 10 + d_0 = 10(d_n \times 10^{n-1} + \cdots + d_1) + d_0 \equiv d_0.$ 

# DFA for Mod 2

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# DFA for Mod 2



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 $\equiv d_n + \cdots + d_1 + d_0$ 

# DFA for Divisible by 3

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**Did you Know?**  $n \equiv 0$  iff



**Did you Know?**  $n \equiv 0$  iff last 2 digits are a number  $\equiv 0$ .

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**Did you Know?**  $n \equiv 0$  iff last 2 digits are a number  $\equiv 0$ . **Thm**  $d_n \cdots d_0 \equiv 2d_1 + d_0$ .

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 $d_n \times 10^n + \cdots + d_1 \times 10 + d_0$ 

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$$d_n \times 10^n + \cdots + d_1 \times 10 + d_0$$

 $\equiv d_1 \times 10 + d_0$ 

**Did you Know?**  $n \equiv 0$  iff last 2 digits are a number  $\equiv 0$ . **Thm**  $d_n \cdots d_0 \equiv 2d_1 + d_0$ . **Pf** I'll have 250H prove by induction  $(\forall n \ge 2)[10^n \equiv 0]$ . Hence

$$d_n imes 10^n + \dots + d_1 imes 10 + d_0$$
  
 $\equiv d_1 imes 10 + d_0$   
 $\equiv 2d_1 + d_0.$ 

# **DFA** for Divisible by 4

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For all of these problems we need to find a pattern of  $10^n \pmod{a}$ .

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# Tricks for Mod 5 and Mod 6

These may be on a HW.



Is there a trick for mod 11? Did you Know?  $n \equiv 0$  iff  $\pm$  sum of digits is  $\equiv 0$ .

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Is there a trick for mod 11? **Did you Know?**  $n \equiv 0$  iff  $\pm$  sum of digits is  $\equiv 0$ . **Thm**  $d_n \cdots d_0 \equiv d_0 - d_1 + d_2 - \cdots + d_n$ .
#### Trick for Mod 11. All $\equiv$ are Mod 11

Is there a trick for mod 11? **Did you Know?**  $n \equiv 0$  iff  $\pm$  sum of digits is  $\equiv 0$ . **Thm**  $d_n \cdots d_0 \equiv d_0 - d_1 + d_2 - \cdots + d_n$ . Proof may be on HW or Midterm or Final or some combination.

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Need to keep track of both the running weighted sum mod 11 and if you are reading an even or odd place.

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 $Q=\{0,\ldots,10\}\times\{0,1\}$ 

Need to keep track of both the running weighted sum mod 11 and if you are reading an even or odd place.

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$$Q = \{0, \dots, 10\} \times \{0, 1\}$$
  
s = (0, 0).

Need to keep track of both the running weighted sum mod 11 and if you are reading an even or odd place.

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$$Q = \{0, \ldots, 10\} \times \{0, 1\}$$

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Final state: Not going to have these, this is DFA-classifier.

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$$Q = \{0, \dots, 10\} \times \{0, 1\}$$

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$$\delta((i,j),\sigma) \begin{cases} (i+\sigma \pmod{11}, j+1 \pmod{2}) \text{ if } j=0\\ (i-\sigma \pmod{11}, j+1 \pmod{2}) \text{ if } j=1 \end{cases}$$
(1)

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We keep track of a running weighted sum mod 11 and position of the symbol mod 2.

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22 states.

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**Classifier** If end in (i, 0) or (i, 1) then number is  $\equiv i$ .

Is there a trick for mod 7? VOTE

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Is there a trick for mod 7? VOTE Answer Depends what you call a trick.

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Want to know what 3876554 is mod 7. All arith is mod 7.  $4\times1+5\times3+5\times2+6\times6+7\times4+8\times5+3\times1$ 

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Want to know what 3876554 is mod 7. All arith is mod 7.  $4 \times 1 + 5 \times 3 + 5 \times 2 + 6 \times 6 + 7 \times 4 + 8 \times 5 + 3 \times 1$ We do this mod 7 so the numbers do not get that big

4 + 15 + 10 + 36 + 28 + 40 + 3

$$\equiv 4+1+3+1+0+5+3 \equiv (4+3+1)+(3+1+5+3) \equiv 1+5 \equiv 6.$$

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 $\equiv 4+1+3+1+0+5+3 \equiv (4+3+1)+(3+1+5+3) \equiv 1+5 \equiv 6.$ DFA States will keep track of

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**DFA** States will keep track of Running weighted sum mod 7

Want to know what 3876554 is mod 7. All arith is mod 7.  $4 \times 1 + 5 \times 3 + 5 \times 2 + 6 \times 6 + 7 \times 4 + 8 \times 5 + 3 \times 1$ We do this mod 7 so the numbers do not get that big

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**DFA** States will keep track of Running weighted sum mod 7 Position of digit mod 6 so know which weights to use.

Want to know what 3876554 is mod 7. All arith is mod 7.  $4 \times 1 + 5 \times 3 + 5 \times 2 + 6 \times 6 + 7 \times 4 + 8 \times 5 + 3 \times 1$ We do this mod 7 so the numbers do not get that big

4 + 15 + 10 + 36 + 28 + 40 + 3

 $\equiv 4+1+3+1+0+5+3 \equiv (4+3+1)+(3+1+5+3) \equiv 1+5 \equiv 6.$ 

**DFA** States will keep track of Running weighted sum mod 7 Position of digit mod 6 so know which weights to use. So there are  $7 \times 6 = 42$  states.

### Is the Method a Trick?

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# Is the Method a Trick?

YES A DFA can do it.



## Is the Method a Trick?

YES A DFA can do it.

**NO** A human can't do it easily- the pattern is not like 1,1,1,... or mostly 0's.

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**BILL:** Saadiq, please draw a DFA for  $\{n : n \equiv 0 \pmod{7}\}$ .



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Might make it a HW to do as a table.

What is the fastest way to determine  $n \pmod{7}$ ?

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What is the fastest way to determine  $n \pmod{7}$ ? Method One Divide and take remainder.

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What is the fastest way to determine *n* (mod 7)? Method One Divide and take remainder. Method Two Use the DFA.

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What is the fastest way to determine n (mod 7)?Method One Divide and take remainder.Method Two Use the DFA.Question Which is faster?

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What is the fastest way to determine n (mod 7)?
Method One Divide and take remainder.
Method Two Use the DFA.
Question Which is faster?
Might be hard to tell because today's computers are so fast!

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## BILL, STOP RECORDING LECTURE!!!!

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#### BILL STOP RECORDING LECTURE!!!