BILL AND NATHAN, RECORD LECTURE!!!!

BILL RECORD LECTURE!!!

Factoring
Is Probably Not NPC

BILL START RECORDING

Factoring: Some History

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I think it is unlikely that anyone aside from myself will ever know.

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  We know about the role of computers to speed up calculations, but it's reasonable it never dawned on him.
- **Conclusion**
  - His arrogance: assumed the world would not change much.
  - Our arrogance: knowing how much the world did change.

Factoring Algorithms

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- We measure the run time as a function of $\lg N$ which is the **length** of the input. We may use L for this.

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- New SVP algorithm: Unclear!

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This is an informal diff between Factoring and SAT.

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Hence

$$\mathrm{FACT}\ \text{NPC} \rightarrow \mathrm{NP} = \text{co-NP}, \text{ which is unlikely.}$$

Primality in NP

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We will present PRIMES in NP and that is all we will need in our proof that $\text{FACT} \in \text{co-NP}$.

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We abbreviate **certificate** by **cert**.

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then n is prime.

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So need the cert to contain a cert that the claimed prime factors of $n-1$ are prime.

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So it's a recursive cert.
Need to check that the cert is short, but this is not difficult.

Back to Factoring

$\overline{\mathrm{FACT}} \in \mathrm{NP}$

$$\mathrm{FACT} = \{(n,a) : (\exists b \le a)[b \text{ divides } n]\}$$

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Verifier has to check

1. $n = p_1^{c_1} \cdots p_k^{c_k}$.
2. $a < p_1$.
3. Each p_i is prime.

Recap What We Know

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Next slide.

The Future of Factoring

I paraphrase **The Joy of Factoring** by Wagstaff:
The best factoring algorithms have time complexity of the form

$$e^{c(\ln N)^t(\ln \ln N)^{1-t}}$$

with Q.Sieve using $t = \frac{1}{2}$ and N.F.Sieve using $t = \frac{1}{3}$.
Moreover, any method that uses B-factoring must take this long.

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  - Anthony, Davin, Erika, Jacob, and Nathan have not yet applied Ramsey theory to this problem.

BILL AND NATHAN STOP RECORDING