

Hard Cases for SAT Solvers

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Example

The AND of the following:

1. $x_{11} \vee x_{12}$

2. $x_{21} \vee x_{22}$

3. $x_{31} \vee x_{32}$

4. $\neg x_{11} \vee \neg x_{21}$

5. $\neg x_{11} \vee \neg x_{31}$

6. $\neg x_{21} \vee \neg x_{31}$

7. $\neg x_{12} \vee \neg x_{22}$

8. $\neg x_{12} \vee \neg x_{32}$

9. $\neg x_{22} \vee \neg x_{32}$

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Can't put 3 pigeons into 2 holes! So Fml is NOT satisfiable.

PHP: Pigeon Hole Principle

Let $n < m$. n is NUMBER OF HOLES, m is NUMBER OF PIGEONS. x_{ij} will be thought of as Pigeon i IS in Hole j .

Definition

PHP_n^m is the AND of the following:

1. For $1 \leq i \leq m$

$$x_{i1} \vee x_{i2} \vee \cdots \vee x_{in}$$

(Pigeon i is in SOME Hole.)

2. For $1 \leq i_1 < i_2 \leq m$ and $1 \leq j \leq n$

$$\neg x_{i_1 j} \vee \neg x_{i_2 j}$$

(Hole j does not have BOTH Pigeon i_1 and Pigeon i_2 .)

NOTE: PHP_n^m has nm **VARS** and $O(mn^2)$ **CLAUSES** and is NOT satisfiable.

What is Known

1. If $n < m$ then PHP_n^m is not satisfiable.
2. The proof of this is by the Pigeon hole principle and not by Truth Table, it was by mathematical reasoning.
3. There is a proof technique called **Resolution** that is used to show formulas are **not satisfiable**. It is known that resolution proofs that PHP_n^m is not satisfiable are large.
4. Our speculation is that the SAT Solvers we have been studying will take a long time on PHP_n^m .
5. Try our out SAT solvers on PHP_n^{n+1} , PHP_n^{n+2} , ... and see if it takes a long time. See what happens as the m gets bigger.