## BILL, RECORD LECTURE!!!!

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## Today: Admin, <br> <br> Intro to Theory of <br> <br> Intro to Theory of Computation

## Admin

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## Necessary Administrative

Everything in these slides is also on the written syllabus on the course website.

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Course Webpage:
https://www.cs.umd.edu/users/gasarch/COURSES/452/S24/
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1. Taught by William Gasarch. Tu-Th 3:30-4:45 in CSI 3117.
2. TAed by Cheng-Yuan 'Sam' Lee, Adam Melrod and Isaac Mammel.

## Necessary administrative stuff

- Course Website: Will post slides, notes, and HW there.
- Elms: will post recording.
- Gradescope: you will submit HW there.
- Gradescope: we will grade HW there.
- Regrade requests due within a week of the HW being graded.
- Piazza is great for asking questions.

IF you are auditing this class for whatever reason- perhaps you are having a hard time getting permission to take it, or perhaps you like the material but don't want to take it, let me know and I will put you on the class email list and invite you to join the Piazza.

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- Appointments (possibly on zoom).


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- There will be one short programming project. (This is not a course like CMSC 412 where the project IS the course.)

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I am not sure why you are telling me about time stamps, but, as the kids say, whatever.

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## Elementary Theory of Computation

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This question permeates all branches of mathematics and computer science.

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Example: Gauss invented the Fast Fourier Transform but never told anyone since he did not think it was that important.

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6. HALT is undecidable (Turing, 1950's.)

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5. We will define problems that are HARDER THAN HALT.

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