BILL, RECORD LECTURE!!!!

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Today: Admin, Intro to Theory of Computation

Admin

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Necessary Administrative

Everything in these slides is also on the written syllabus on the course website.

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- 1. Taught by William Gasarch. Tu-Th 3:30-4:45 in CSI 3117.
- 2. TAed by Cheng-Yuan 'Sam' Lee, Adam Melrod and Isaac Mammel.

Necessary administrative stuff

- ► Course Website: Will post slides, notes, and HW there.
- Elms: will post recording.
- Gradescope: you will submit HW there.
- ► Gradescope: we will grade HW there.
- Regrade requests due within a week of the HW being graded.
- Piazza is great for asking questions.

IF you are auditing this class for whatever reason- perhaps you are having a hard time getting permission to take it, or perhaps you like the material but don't want to take it, let me know and I will put you on the class email list and invite you to join the Piazza.

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Office hours

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- Office hours Bill: Tu 11:00-12:15, 2:00-3:15, an Th same times. IRB 2242.
 Sam: Monday 4-6 in AVW 4160.
 Isaac: Wedensday 5-7 in AVW 4160
 Adam: Friday 1-3 in AVW 4160

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Sam: Monday 4-6 in AVW 4160. Isaac: Wedensday 5-7 in AVW 4160 Adam: Friday 1-3 in AVW 4160

Email us- put 250 in the subject line.
 Bill: gasarch@umd.edu or (301) 503-3157.
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- Appointments (possibly on zoom).

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Mathematical maturity.

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- There will be one short programming project. (This is not a course like CMSC 412 where the project IS the course.)

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3.5 Recording might not work that day.

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I am not sure why you are telling me about **time stamps**, but, as the kids say, whatever.

Textbook

Required Text None. Recommended Text None. If you really want a text then buy used (cheap) or borrow:

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Introduction to The Theory of Computation by Michael Sipser There will be notes, slides, and recordings of lecture online.

You are INVITED to talk to us

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If **for whatever reason** you are falling behind in the class, or are having trouble with the HW, see us in office hours or **you can make an appointment to see us!** Either in person or on zoom.

Elementary Theory of Computation

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Our Key Question

Given a problem, classify how hard it is.

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Given a problem, **classify** how hard it is.

This question permeates all branches of mathematics and computer science.

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 There was an awareness of computational problems taking an amount of time, but it was not rigorous.
 Example: Gauss invented the Fast Fourier Transform but never told anyone since he did not think it was that important.

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 Given a real, we want to know how hard it is in terms of the lowest degree polynomial over Z that it is a root of.
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3. There is no set of axioms from which one can derive all the truths of arithmetic. (Godel's Incompleteness Theorem, 1933.)

The Age of Computers.



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6. HALT is undecidable (Turing, 1950's.)

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- 3. Regular languages have many closure properties. We will prove this.
- 4. We will prove some sets are **not regular**.

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 - 3.2 Sets CSGs can generate that CFGs cannot generate.

- 1. Grammars generate sets of strings.
- 2. Two kinds we study: Context Free Grammars (CFG) and Context Sensitive Grammars (CSG).
- 3. We will study grammars briefly, emphasizing:
 - 3.1 Sets CFGs can generate that are not regular.
 - 3.2 Sets CSGs can generate that CFGs cannot generate.
 - 3.3 Problems that can be done with a large DFA or NFA, but only need a small CFG.

Polynomial Time and Non-Deterministic Polynomial Time.

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- 4. P, NP have many closure properties. We will prove this.

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Decidable sets and undecidable sets.

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- 4. DEC and UNDEC have many closure properties. We will prove this.
- 5. We will define problems that are HARDER THAN HALT.

BILL, STOP RECORDING LECTURE!!!!

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