An Application of Ramsey's Theorem to Proving Programs Terminate: An Exposition

William Gasarch-U of MD

Who is Who

1. Work by
   1.1 **Floyd**,
   1.2 **Byron Cook, Andreas Podelski, Andrey Rybalchenko**,
   1.3 **Lee, Jones, Ben-Amram**
   1.4 Others

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3. **Pre-Brag**: Not my area-some things may be understandable.

Overview I

Problem: Given a program we want to prove it terminates no matter what user does (called TERM problem).

1. **Impossible in general**- Harder than Halting.
2. **But** can do this on some simple progs. (We will.)

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4. Do another example with **Ramsey Theory**.
5. Do example with **Ramsey Theory** and Matrices.

Notation

1. Will use psuedo-code progs.
2. **KEY:** If A is a set then the command

   ```
   x = input(A)
   ```

   means that x gets some value from A that the user decides.
3. **Note:** we will want to show that **no matter what the user does** the program will halt.
4. The code

   ```
   (x,y) = (f(x,y),g(x,y))
   ```

   means that simultaneously x gets f(x,y) and y gets g(x,y).

Example of Traditional Method

```
(x,y,z) = (input(INT), input(INT), input(INT))
While x>0 and y>0 and z>0
        control = input(1,2,3)
        if control == 1 then
                (x,y,z)=(x+1,y-1,z-1)
        else
        if control == 2 then
                (x,y,z)=(x-1,y+1,z-1)
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Discuss Can you prove this program **always** terminates?

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Discuss Can you prove this program **always** terminates?
Whatever the user does x+y+z is decreasing.
Eventually x+y+z=0 so prog terminates there or earlier.

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General method due to **Floyd:** Find a function f(x,y,z) from the values of the variables to N such that

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2. if f(x,y,z) is ever 0 then the program **must have halted.**

Note: Method is more general- can map to a well founded order such that in every iteration f(x,y,z) decreases in that order, and if f(x,y,z) is ever a min element then program must have halted.

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Use Lex Order: $(0,0,0) < (0,0,1) < \cdots < (0,1,0) \cdots$.
Note: $(4, 10^{100}, 10^{10!}) < (5,0,0)$.

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Note: $(4, 10^{100}, 10^{10!}) < (5, 0, 0)$.
In every iteration (x, y, z) **decreases in this ordering.**
If hits bottom then all vars are 0 so **must halt then or earlier.**

Well Ordering is Key!

Def An ordering $(X, \preceq)$ is a **well founded** if there are no infinite decreasing sequences. (Induction proofs can be done on such orderings.)

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Z in its usual ordering is NOT well founded.

Lex order on $\mathsf{N} \times \mathsf{N} \times \mathsf{N}$ is well founded. Discuss.

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2. **Good News:** We only had to reason about what happens in **one** iteration.

Keep these in mind- our later proof will use a **nice** ordering but will need to reason about a **block** of instructions.

Digression Into Ramsey Theory (Parties!)

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3. If you have 2^{2k-1} people at a party then either k of them mutually know each other of k of them mutually do not know each other.
4. If you have an **infinite** number of people at a party then either there exists an **infinite** subset that all know each other or an **infinite** subset that all do not know each other.

Digression Into Ramsey Theory (Math!)

Def Let $c, k, n \in \mathsf{N}$. K_n is the **complete graph on n vertices (all pairs are edges).** K_{N} is the **infinite complete graph.** A **c-coloring of K_n** is a c-coloring of the edges of K_n. A **homog set** is a subset H of the vertices such that every pair has the same color (e.g., 10 people all of whom know each other).

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3. For all c-colorings of the K_{N} there exists an infinite homog set.

Alt Proof Using Ramsey

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Proof of termination

If program does not halt then there is infinite sequence $(x_1, y_1, z_1), (x_2, y_2, z_2), \ldots,$ representing state of vars.

Reasoning about Blocks

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Look at $(x_i, y_i, z_i), \ldots, (x_j, y_j, z_j)$.

1. If control is ever 1 then $x_i > x_j$.
2. If control is never 1 then $y_i > y_j$.

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1. If control is ever 1 then $x_i > x_j$.
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Upshot: For all $i < j$ either $x_i > x_j$ or $y_i > y_j$.

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For all $i < j$ either $\mathrm{x}_i > \mathrm{x}_j$ or $\mathrm{y}_i > \mathrm{y}_j$.
Define a 2-coloring of the edges of K_{N}:

$$COL(i,j) = \begin{cases} X \text{ if } \mathrm{x}_i > \mathrm{x}_j \\ Y \text{ if } \mathrm{y}_i > \mathrm{y}_j \end{cases} \tag{1}$$

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If color is X then $\mathrm{x}_{i_1} > \mathrm{x}_{i_2} > \mathrm{x}_{i_3} > \cdots$

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If color is X then $x_{i_1} > x_{i_2} > x_{i_3} > \cdots$
If color is Y then $y_{i_1} > y_{i_2} > y_{i_3} > \cdots$
In either case will have eventually have a var ≤ 0 and hence program must terminate. **Contradiction.**

Compare and Contrast

1. Trad. proof used lex order on $\mathbb{N}^3$–complicated!
2. Ramsey Proof used natural ordering on $\mathbb{N}$–simple!
3. Trad. proof only had to reason about single steps–simple!
4. Ramsey Proof had to reason about blocks of steps–complicated!

What do YOU think?

VOTE:

1. Traditional Proof!
2. Ramsey Proof!

Another Example

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Reasoning about Blocks

If program does not halt then there is infinite sequence $(\mathrm{x}_1, \mathrm{y}_1), (\mathrm{x}_2, \mathrm{y}_2), \ldots,$ representing state of vars.
We look at a block $(x_i, y_i), \ldots, (x_j, y_j)$.

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We look at a block $(x_i, y_i), \ldots, (x_j, y_j)$.
Case 1 If control is never 1 then $x_i + y_i > x_j + y_j$.

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Case 1 If control is never 1 then $x_i + y_i > x_j + y_j$.
Case 2 If control is ever 1 then assume there are a 2's first. After a 2's we have $(\mathrm{x}_i - a, \mathrm{x}_i)$. Then with the one 1 we have $(\mathrm{x}_i - 2, \mathrm{x}_i - a + 1)$. Can show that $\mathrm{x}_i > \mathrm{x}_j$.

Use Ramsey!

Define a 2-coloring of the edges of K_N:

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Comments

1. The condition
   $x_i > x_j$ OR $x_i + y_i > x_j + y_j$.
   in the last proof is called a **Termination Invariant**. It is used to strengthen the induction hypothesis.
2. The proof was **found by the system** of B. Cook et al.
3. Looking for a Termination Invariant is the hard part to automate but they have automated it.
4. Can we use these techniques to solve a fragment of Termination Problem?

Model control=1 via a Matrix

```
if control == 1 then (x,y)=(x-1,x)
```

Model as a matrix A indexed by x,y,x+y.

$$\begin{pmatrix} -1 & 0 & \infty \\ \infty & \infty & \infty \\ \infty & \infty & \infty \end{pmatrix}$$

For a,b $\in$ {x,y,x+y}
Entry (a,b) is difference between NEW b and OLD a.
Entry (a,a) is most interesting- if neg then a decreased.

Model control=2 via a Matrix

```
if control == 2 then (x,y)=(y-2,x+1)
```

Model as a matrix B indexed by x,y,x+y.

$$\begin{pmatrix} \infty & 1 & \infty \\ -2 & \infty & \infty \\ \infty & \infty & -1 \end{pmatrix}$$

Redefine Matrix Mult

A and B matrices, C=AB defined by

$$c_{ij} = \min_k \{a_{ik} + b_{kj}\}.$$

Lemma

If matrix A models a statement s_1 and matrix B models a statement s_2 then matrix AB models what happens if you run $s_1; s_2$.

Matrix Proof that Program Terminates

- A is matrix for control=1. B is matrix for control=2.
- Show: any prod of A's and B's some diag is negative.
- Hence in any finite seg one of the vars decreases.
- Hence, by Ramsey proof, the program always terminates

General Program

```
X = (input(INT),...,input(INT))
While x[1]>0 and x[2]>0 and ... x[n]>0
 control = input(1,2,3,...,m)
 if control==1
    X = F1(X,input(INT),...,input(INT)))
  else
  if control==2
    X = F2(X,input(INT),...,input(INT))
  else...
  else
  if control==m
    X = Fm(X,input(INT),...,input(INT))
```

Fragment of TERM decidable?

Def The **TERMINATION PROBLEM:** Given $F_1, \ldots, F_m$ can we determine if the following holds:

For all ω-seq of inputs the program halts

Much Easier Problem Undecidable

History Lesson: In 1900 David Hilbert proposed 23 problems for mathematicians to work on over the next 100 years.
Hilberts Tenth Problem (in modern terminology):
Give an algorithm that will, given a polynomial $p(x_1, \ldots, x_n)$ over Z, determines if there exists $a_1, \ldots, a_n \in Z$ such that $p(a_1, \ldots, a_n) = 0$.

- Hilbert thought there was such an algorithm and that this was a problem in Number Theory.
- Over time (next slide) it was proven that there is NO such algorithm and that this is a problem in Logic.

Computable and C.E. Sets

Def: A set A is **computable** if there is a Java program (Turing Machine, other models) J (on one var) that halts on all inputs such that
If $x \in A$ then $J(x)$=YES
If $x \notin A$ then $J(x)$=NO
Def: A set A is **computably enumerable (c.e.)** (also called Σ_1) if there is a Java program J (on two vars) that halts on all inputs such that
If $x \in A$ then $(\exists y)[J(x,y) = \text{YES}]$.
If $x \notin A$ then $(\forall y)[J(x,y) = \text{NO}]$.
Known: There are sets that are c.e. but not computable. Here is one: Let J_x be the xth Java program in some reasonable ordering.

$$\{(x,y) \colon J_x(y) \text{ halts }\} = \{(x,y) \colon (\exists t)[J_x(y) \text{ halts in } \leq t \text{ steps}]\ \}$$

Back to Hilbert's Tenth

1. In 1959 Davis-Putnam-Robinson showed that for **every** c.e. set A there exists an exp-poly (so can include vars as exponents) $p(x, x_1, \ldots, x_n)$ such that

   $$A = \{a \colon (\exists a_1, \ldots, a_n)[p(a, a_1, \ldots, a_n)]\}$$

   Needed just ONE step to get down to polynomials.
2. In 1970 Yuri Matiyasevich supplies that one missing step. So ALL c.e. sets (including undecidable ones) can be written in terms of solutions to polynomials.
3. From all of this you can conclude Hilbert's Tenth Problem is Unsolvable.
4. From this you can conclude that TERM is undecidable.

Termination Problem More Than Undecidable

The **TERMINATION PROBLEM:** Given $F_1, \ldots, F_m$ can we determine if the following holds:

For all ω-seq of inputs the program halts

1. This is **HARDER** than **HALT.** This is Σ_1^1-complete. Infinitely harder than HALT!
2. **EASY** to show is **HARD:** use polynomials and Hilbert's Tenth Problem. This shows a much easier version of the problem undecidable.
3. **OPEN:** Determine which subsets of F_i make this decidable? Σ_1^1-complete? Other?

(New Topic) Didn't Need Full Strength of Ramsey

The colorings we applied Ramsey to were of a certain type:

Def A coloring of the edges of K_n or $K_{\mathbb{N}}$ is **transitive** if, for every $i < j < k$, if $COL(i,j) = COL(j,k)$ then both equal $COL(i,k)$.

1. Our colorings were transitive.
2. **Transitive Ramsey Thm** is weaker than **Ramsey's Thm**.

Transitive Ramsey Weaker than Ramsey

TR is Transitive Ramsey, R is Ramsey.

1. **Combinatorially:** $R(k,c) = c^{\Theta(ck)}$, $TR(k,c) = (k-1)^c + 1$. This may look familiar

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2. **Computability:** There exists a computable 2-coloring of K_{N} with no computable homog set (can even have no Σ_2 homog set). For every transitive computable c-coloring of K_{N} there exists a computable homog set (folklore).

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3. **Proof Theory:** Over the axiom system RCA_0, R implies TR, but TR does not imply R.

Summary

1. Ramsey Theory can be used to prove some simple programs terminate that seem harder to do by traditional methods. Interest to **PL.**
2. Some subcases of **TERMINATION PROBLEM** are decidable. Of interest to **PL** and **Logic.**
3. Full strength of Ramsey not needed. Interest to **Logicians** and **Combinatorists.**