

# The Muffin Problem

Guangi Cui - Montgomery Blair HS

John Dickerson- University of MD

Naveen Durvasula - Montgomery Blair HS

William Gasarch - University of MD

Erik Metz - University of MD

Jacob Prinz-University of MD

Naveen Raman - Richard Montgomery HS

Daniel Smolyak- University of MD

Sung Hyun Yoo - Bergen County Academies (in NJ)

# How it Began

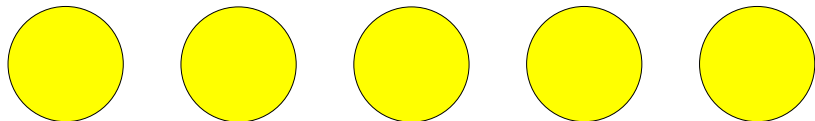
## A Recreational Math Conference (Gathering for Gardner) May 2016

I found a pamphlet:

### The Julia Robinson Mathematics Festival: A Sample of Mathematical Puzzles Compiled by Nancy Blachman

which had this problem, proposed by Alan Frank:

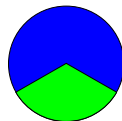
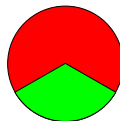
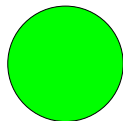
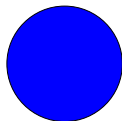
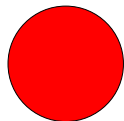
*How can you divide and distribute 5 muffins to 3 students so that every student gets  $\frac{5}{3}$  where nobody gets a tiny sliver?*



# Five Muffins, Three Students, Proc by Picture

Person	Color	What they Get
Alice	RED	$1 + \frac{2}{3} = \frac{5}{3}$
Bob	BLUE	$1 + \frac{2}{3} = \frac{5}{3}$
Carol	GREEN	$1 + \frac{1}{3} + \frac{1}{3} = \frac{5}{3}$

Smallest Piece:  $\frac{1}{3}$



## Can We Do Better?

The smallest piece in the above solution is  $\frac{1}{3}$ .

**Is there a procedure with a larger smallest piece?**

**VOTE**

## Can We Do Better?

The smallest piece in the above solution is  $\frac{1}{3}$ .

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**VOTE**

- ▶ **YES**
- ▶ **NO**

## Can We Do Better?

The smallest piece in the above solution is  $\frac{1}{3}$ .

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**VOTE**

- ▶ **YES**
- ▶ **NO**

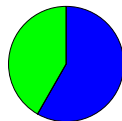
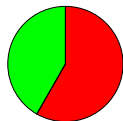
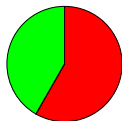
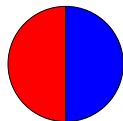
**YES WE CAN!**

We use **!** since we are excited that we can!

## Five Muffins, Three People—Proc by Picture

Person	Color	What they Get
Alice	RED	$\frac{6}{12} + \frac{7}{12} + \frac{7}{12}$
Bob	BLUE	$\frac{6}{12} + \frac{7}{12} + \frac{7}{12}$
Carol	GREEN	$\frac{5}{12} + \frac{5}{12} + \frac{5}{12} + \frac{5}{12}$

**Smallest Piece:**  $\frac{5}{12}$



# Can We Do Better?

The smallest piece in the above solution is  $\frac{5}{12}$ .

**Is there a procedure with a larger smallest piece?**

**VOTE**

- ▶ **YES**
- ▶ **NO**



# Can We Do Better?

The smallest piece in the above solution is  $\frac{5}{12}$ .

**Is there a procedure with a larger smallest piece?**

**VOTE**

- ▶ **YES**
- ▶ **NO**

**NO WE CAN'T!**

We use **!** since we are excited to prove we can't do better!

## Five Muffins, Three People—Can't Do Better Than $\frac{5}{12}$

There is a procedure for 5 muffins, 3 students where each student gets  $\frac{5}{3}$  muffins, smallest piece  $N$ . We want  $N \leq \frac{5}{12}$ .

**Case 0:** Some muffin is uncut. Cut it  $(\frac{1}{2}, \frac{1}{2})$  and give both  $\frac{1}{2}$ -sized pieces to whoever got the uncut muffin. (Note  $\frac{1}{2} > \frac{5}{12}$ .) Reduces to other cases.

(**Henceforth:** All muffins are cut into  $\geq 2$  pieces.)

**Case 1:** Some muffin is cut into  $\geq 3$  pieces. Then  $N \leq \frac{1}{3} < \frac{5}{12}$ .

(**Henceforth:** All muffins are cut into 2 pieces.)

**Case 2:** All muffins are cut into 2 pieces. 10 pieces, 3 students: **Someone** gets  $\geq 4$  pieces. He has some piece

$$\leq \frac{5}{3} \times \frac{1}{4} = \frac{5}{12} \quad \text{Great to see } \frac{5}{12}$$

# Be Amazed Now! And Later!

1. Procedure for 5 muffins, 3 people, smallest piece  $\frac{5}{12}$ .
2. NO Procedure for 5 muffins, 3 people, smallest piece  $> \frac{5}{12}$ .

**Amazing That Have Exact Result!**

# Be Amazed Now! And Later!

1. Procedure for 5 muffins, 3 people, smallest piece  $\frac{5}{12}$ .
2. NO Procedure for 5 muffins, 3 people, smallest piece  $> \frac{5}{12}$ .

**Amazing That Have Exact Result!**

Prepare To Be More Amazed! On Next Page!

## Amazing Results!

1. Procedure for 47 muffins, 9 people, smallest piece  $\frac{111}{234}$ .
2. NO Procedure for 47 muffins, 9 people, smallest piece  $> \frac{111}{234}$ .

# Amazing Results!

1. Procedure for 47 muffins, 9 people, smallest piece  $\frac{111}{234}$ .
2. NO Procedure for 47 muffins, 9 people, smallest piece  $> \frac{111}{234}$ .
1. Procedure for 52 muffins, 11 people, smallest piece  $\frac{83}{176}$ .
2. NO Procedure for 52 muffins, 11 people, smallest piece  $> \frac{83}{176}$ .

# Amazing Results!

1. Procedure for 47 muffins, 9 people, smallest piece  $\frac{111}{234}$ .
2. NO Procedure for 47 muffins, 9 people, smallest piece  $> \frac{111}{234}$ .
1. Procedure for 52 muffins, 11 people, smallest piece  $\frac{83}{176}$ .
2. NO Procedure for 52 muffins, 11 people, smallest piece  $> \frac{83}{176}$ .
1. Procedure for 35 muffins, 13 people, smallest piece  $\frac{64}{143}$ .
2. NO Procedure for 35 muffins, 13 people, smallest piece  $> \frac{64}{143}$ .

# Amazing Results!

1. Procedure for 47 muffins, 9 people, smallest piece  $\frac{111}{234}$ .
2. NO Procedure for 47 muffins, 9 people, smallest piece  $> \frac{111}{234}$ .
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1. Procedure for 35 muffins, 13 people, smallest piece  $\frac{64}{143}$ .
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**All done by hand, no use of a computer**



# Amazing Results!

1. Procedure for 47 muffins, 9 people, smallest piece  $\frac{111}{234}$ .
2. NO Procedure for 47 muffins, 9 people, smallest piece  $> \frac{111}{234}$ .
1. Procedure for 52 muffins, 11 people, smallest piece  $\frac{83}{176}$ .
2. NO Procedure for 52 muffins, 11 people, smallest piece  $> \frac{83}{176}$ .
1. Procedure for 35 muffins, 13 people, smallest piece  $\frac{64}{143}$ .
2. NO Procedure for 35 muffins, 13 people, smallest piece  $> \frac{64}{143}$ .

**All done by hand, no use of a computer**

**Co-author Erik Metz is a *muffin savant***

## General Problem

*How can you divide and distribute  $m$  muffins to  $s$  students so that each student gets  $\frac{m}{s}$  AND the MIN piece is MAXIMIZED?*

An  $(m, s)$ -*procedure* is a way to divide and distribute  $m$  muffins to  $s$  students so that each student gets  $\frac{m}{s}$  muffins.

An  $(m, s)$ -procedure is *optimal* if it has the largest smallest piece of any procedure.

$f(m, s)$  be the smallest piece in an optimal  $(m, s)$ -procedure.

We have shown  $f(5, 3) = \frac{5}{12}$ .

**Note:**  $f(m, s) \geq \frac{1}{s}$ : divide each muffin into  $s$  pieces of size  $\frac{1}{s}$  and give each student  $m$  of them.

$$f(3, 5) \geq ?$$

Clearly  $f(3, 5) \geq \frac{1}{5}$ . Can we get  $f(3, 5) > \frac{1}{5}$ ?  
Think about it at your desk.

$$f(3, 5) \geq ?$$

Clearly  $f(3, 5) \geq \frac{1}{5}$ . Can we get  $f(3, 5) > \frac{1}{5}$ ?

Think about it at your desk.

$$f(3, 5) \geq \frac{1}{4}$$

1. Divide 2 muffin  $[\frac{6}{20}, \frac{7}{20}, \frac{7}{20}]$
2. Divide 1 muffin  $[\frac{5}{20}, \frac{5}{20}, \frac{5}{20}, \frac{5}{20}]$
3. Give 4 students  $(\frac{5}{20}, \frac{7}{20})$
4. Give 1 students  $(\frac{6}{20}, \frac{6}{20})$

$$f(3, 5) \geq ?$$

Clearly  $f(3, 5) \geq \frac{1}{5}$ . Can we get  $f(3, 5) > \frac{1}{5}$ ?

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4. Give 1 students  $(\frac{6}{20}, \frac{6}{20})$

Can we do better? Vote:

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4. Give 1 students  $(\frac{6}{20}, \frac{6}{20})$

Can we do better? Vote:

**YES**

**NO**

**UNKNOWN TO SCIENCE**

$$f(3, 5) \geq ?$$

Clearly  $f(3, 5) \geq \frac{1}{5}$ . Can we get  $f(3, 5) > \frac{1}{5}$ ?

Think about it at your desk.

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3. Give 4 students  $(\frac{5}{20}, \frac{7}{20})$
4. Give 1 students  $(\frac{6}{20}, \frac{6}{20})$

Can we do better? Vote:

**YES**

**NO**

**UNKNOWN TO SCIENCE**

**NO** Proof on next slide.

$$f(3, 5) \leq \frac{1}{4}$$

There is a procedure for 3 muffins, 5 students where each student gets  $\frac{3}{5}$  muffins, smallest piece  $N$ . We want  $N \leq \frac{1}{4}$ .

**Case 0:** Some student gets 1 piece, so size  $\frac{3}{5}$ . Cut that piece in half and give both  $\frac{3}{10}$ -sized pieces to that student. (Note  $\frac{3}{10} > \frac{1}{4}$ .)  
Reduces to other cases.

(**Henceforth:** All students get  $\geq 2$  pieces.)

**Case 1:** Some student gets  $\geq 3$  pieces. Then  $N \leq \frac{3}{5} \times \frac{1}{3} = \frac{1}{5} < \frac{1}{4}$ .  
(**Henceforth:** All students get 2 pieces.)

**Case 2:** All students get 2 pieces. 5 students, so 10 pieces.  
**Some muffin** gets cut into  $\geq 4$  pieces. Some piece  $\leq \frac{1}{4}$ .



## 3 People, 5 Muffins VS 5 People, 3 Muffins

$$f(5, 3) \geq \frac{5}{12}$$

1. Divide 4 muffins  $[\frac{5}{12}, \frac{7}{12}]$
2. Divide 1 muffin  $[\frac{6}{12}, \frac{6}{12}]$
3. Give 2 students  $(\frac{6}{12}, \frac{7}{12}, \frac{7}{12})$
4. Give 1 student  $(\frac{5}{12}, \frac{5}{12}, \frac{5}{12}, \frac{5}{12})$

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1. Divide 4 muffins  $[\frac{5}{12}, \frac{7}{12}]$
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4. Give 1 students  $(\frac{5}{12}, \frac{5}{12}, \frac{5}{12}, \frac{5}{12})$

$$f(3, 5) \geq \frac{1}{4}$$

1. Divide 2 muffin  $[\frac{6}{20}, \frac{7}{20}, \frac{7}{20}]$
2. Divide 1 muffin  $[\frac{5}{20}, \frac{5}{20}, \frac{5}{20}, \frac{5}{20}]$
3. Give 4 students  $(\frac{5}{20}, \frac{7}{20})$
4. Give 1 students  $(\frac{6}{20}, \frac{6}{20})$

## 3 People, 5 Muffins VS 5 People, 3 Muffins

$$f(5, 3) \geq \frac{5}{12}$$

1. Divide 4 muffins  $[\frac{5}{12}, \frac{7}{12}]$
2. Divide 1 muffin  $[\frac{6}{12}, \frac{6}{12}]$
3. Give 2 students  $(\frac{6}{12}, \frac{7}{12}, \frac{7}{12})$
4. Give 1 students  $(\frac{5}{12}, \frac{5}{12}, \frac{5}{12}, \frac{5}{12})$

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3. Give 4 students  $(\frac{5}{20}, \frac{7}{20})$
4. Give 1 students  $(\frac{6}{20}, \frac{6}{20})$

$f(3, 5)$  proc is  $f(5, 3)$  proc but swap Divide/Give and mult by 3/5.

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3. Give 2 students  $(\frac{6}{12}, \frac{7}{12}, \frac{7}{12})$
4. Give 1 students  $(\frac{5}{12}, \frac{5}{12}, \frac{5}{12}, \frac{5}{12})$

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1. Divide 2 muffin  $[\frac{6}{20}, \frac{7}{20}, \frac{7}{20}]$
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$f(3, 5)$  proc is  $f(5, 3)$  proc but swap Divide/Give and mult by 3/5.

**Theorem:**  $f(m, s) = \frac{m}{s} f(s, m)$ .

## Floor-Ceiling Theorem (Generalize $f(5, 3) \leq \frac{5}{12}$ )

$$f(m, s) \leq \max \left\{ \frac{1}{3}, \min \left\{ \frac{m}{s \lceil 2m/s \rceil}, 1 - \frac{m}{s \lfloor 2m/s \rfloor} \right\} \right\}.$$

**Case 0:** Some muffin is uncut. Cut it  $(\frac{1}{2}, \frac{1}{2})$  and give both halves to whoever got the uncut muffin, so reduces to other cases.

**Case 1:** Some muffin is cut into  $\geq 3$  pieces. Some piece  $\leq \frac{1}{3}$ .

**Case 2:** Every muffin is cut into 2 pieces, so  $2m$  pieces.

**Someone** gets  $\geq \lceil \frac{2m}{s} \rceil$  pieces.  $\exists$  piece  $\leq \frac{m}{s} \times \frac{1}{\lceil 2m/s \rceil} = \frac{m}{s \lceil 2m/s \rceil}$ .

**Someone** gets  $\leq \lfloor \frac{2m}{s} \rfloor$  pieces.  $\exists$  piece  $\geq \frac{m}{s} \frac{1}{\lfloor 2m/s \rfloor} = \frac{m}{s \lfloor 2m/s \rfloor}$ .

The other piece from that muffin is of size  $\leq 1 - \frac{m}{s \lfloor 2m/s \rfloor}$ .

## THREE Students

**CLEVERNESS, COMP PROGS** for the procedure.

**Floor-Ceiling Theorem** for optimality.

$$f(1, 3) = \frac{1}{3}$$

$$f(3k, 3) = 1.$$

$$f(3k + 1, 3) = \frac{3k-1}{6k}, k \geq 1.$$

$$f(3k + 2, 3) = \frac{3k+2}{6k+6}.$$

## FOUR Students

**CLEVERNESS, COMP PROGS** for procedures.

**Floor-Ceiling Theorem** for optimality.

$$f(4k, 4) = 1 \text{ (easy)}$$

$$f(1, 4) = \frac{1}{4} \text{ (easy)}$$

$$f(4k + 1, 4) = \frac{4k-1}{8k}, k \geq 1.$$

$$f(4k + 2, 4) = \frac{1}{2}.$$

$$f(4k + 3, 4) = \frac{4k+1}{8k+4}.$$

**Is FIVE student case a Mod 5 pattern?**

**VOTE YES or NO**

## FOUR Students

**CLEVERNESS, COMP PROGS** for procedures.

**Floor-Ceiling Theorem** for optimality.

$$f(4k, 4) = 1 \text{ (easy)}$$

$$f(1, 4) = \frac{1}{4} \text{ (easy)}$$

$$f(4k + 1, 4) = \frac{4k-1}{8k}, k \geq 1.$$

$$f(4k + 2, 4) = \frac{1}{2}.$$

$$f(4k + 3, 4) = \frac{4k+1}{8k+4}.$$

**Is FIVE student case a Mod 5 pattern?**

**VOTE YES or NO**

**YES but with some exceptions**



## FIVE Students, $m = 1, \dots, 11$

$$f(1, 5) = \frac{1}{5} \text{ (easy or use } f(1, 5) = \frac{5}{1}f(5, 1).)$$

$$f(2, 5) = \frac{1}{5} \text{ (easy or use } f(2, 5) = \frac{5}{2}f(5, 2).)$$

$$f(3, 5) = \frac{1}{4} \text{ (use } f(3, 5) = \frac{3}{5}f(5, 3).)$$

$$f(4, 5) = \frac{3}{10} \text{ (use } f(4, 5) = \frac{4}{5}f(5, 4).)$$

$$f(5, 5) = 1 \text{ (Easy and fits pattern)}$$

$$f(6, 5) = \frac{2}{5} \text{ (Use Floor-Ceiling Thm, fits pattern)}$$

$$f(7, 5) = \frac{1}{3} \text{ (Use Floor-Ceiling Thm, NOT pattern)}$$

$$f(8, 5) = \frac{2}{5} \text{ (Use Floor-Ceiling Thm, fits pattern)}$$

$$f(9, 5) = \frac{2}{5} \text{ (Use Floor-Ceiling Thm, fits pattern)}$$

$$f(10, 5) = 1 \text{ (Easy and fits pattern)}$$

$$f(11, 5) = \text{(Will come back to this later)}$$

# FIVE Students

**CLEVERNESS, COMP PROGS** for procedures.

**Floor-Ceiling Theorem** for optimality.

For  $k \geq 1$ ,  $f(5k, 5) = 1$ .

For  $k = 1$  and  $k \geq 3$ ,  $f(5k + 1, 5) = \frac{5k+1}{10k+5}$

For  $k \geq 2$ ,  $f(5k + 2, 5) = \frac{5k-2}{10k}$

For  $k \geq 1$ ,  $f(5k + 3, 5) = \frac{5k+3}{10k+10}$

For  $k \geq 1$ ,  $f(5k + 4, 5) = \frac{5k+1}{10k+5}$

## What About FIVE students, ELEVEN muffins?

$$f(11, 5) \geq \frac{13}{30}.$$

### Procedure:

1. Divide 8 muffins into  $(\frac{13}{30}, \frac{17}{30})$ .
2. Divide 2 muffins into  $(\frac{14}{30}, \frac{16}{30})$ .
3. Divide 1 muffin into  $(\frac{15}{30}, \frac{15}{30})$ .
4. Give 2 students  $[\frac{14}{30}, \frac{13}{30}, \frac{13}{30}, \frac{13}{20}, \frac{13}{20}]$
5. Give 1 student  $[\frac{17}{30}, \frac{17}{30}, \frac{16}{30}, \frac{16}{20}]$
6. Give 2 students  $[\frac{17}{30}, \frac{17}{30}, \frac{17}{30}, \frac{15}{30}]$ .

## What About FIVE students, ELEVEN muffins?

$$f(m, s) \leq \max \left\{ \frac{1}{3}, \min \left\{ \frac{m}{s \lceil 2m/s \rceil}, 1 - \frac{m}{s \lceil 2m/s \rceil} \right\} \right\} \leq 0.44.$$

So

$$\frac{13}{30} \leq f(11, 5) \leq \frac{11}{25} \quad \text{Diff} = 0.006666 \dots$$

## What About FIVE students, ELEVEN muffins?

$$f(m, s) \leq \max \left\{ \frac{1}{3}, \min \left\{ \frac{m}{s \lceil 2m/s \rceil}, 1 - \frac{m}{s \lfloor 2m/s \rfloor} \right\} \right\} \leq 0.44.$$

So

$$\frac{13}{30} \leq f(11, 5) \leq \frac{11}{25} \quad \text{Diff} = 0.006666 \dots$$

### VOTE:

1.  $f(11, 5) = \frac{13}{30}$ : Needs NEW methods to bound  $f(m, s)$ .
2.  $f(11, 5) = \frac{11}{25}$ : Needs NEW better procedure.
3.  $f(11, 5) = \alpha$  where  $\frac{13}{30} < \alpha < \frac{11}{25}$ . Needs both:
4. **UNKNOWN TO SCIENCE!**

## What About FIVE students, ELEVEN muffins?

$$f(m, s) \leq \max \left\{ \frac{1}{3}, \min \left\{ \frac{m}{s \lceil 2m/s \rceil}, 1 - \frac{m}{s \lfloor 2m/s \rfloor} \right\} \right\} \leq 0.44.$$

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4. **UNKNOWN TO SCIENCE!**

$$\text{KNOWN: } f(11, 5) = \frac{13}{30}$$

**HAPPY:** New opt tech more interesting than new proc.

## $f(11, 5) = \frac{13}{30}$ , Easy Case Based on Muffins

There is a procedure for 11 muffins, 5 students where each student gets  $\frac{11}{5}$  muffins, smallest piece  $N$ . We want  $N \leq \frac{13}{30}$ .

**Case 0:** Some muffin is uncut. Cut it  $(\frac{1}{2}, \frac{1}{2})$  and give both halves to whoever got the uncut muffin. Reduces to other cases.

**Case 1:** Some muffin is cut into  $\geq 3$  pieces.  $N \leq \frac{1}{3} < \frac{13}{30}$ .

**(Negation of Case 0 and Case 1:** All muffins cut into 2 pieces.)

## $f(11, 5) = \frac{13}{30}$ , Easy Case Based on Students

**Case 2:** Some student gets  $\geq 6$  pieces.

$$N \leq \frac{11}{5} \times \frac{1}{6} = \frac{11}{30} < \frac{13}{30}.$$

**Case 3:** Some student gets  $\leq 3$  pieces.

One of the pieces is

$$\geq \frac{11}{5} \times \frac{1}{3} = \frac{11}{15}.$$

Look at the muffin it came from to find a piece that is

$$\leq 1 - \frac{11}{15} = \frac{4}{15} < \frac{13}{30}.$$

**(Negation of Cases 2 and 3:** Every student gets 4 or 5 pieces.)



## $f(11, 5) = \frac{13}{30}$ , Fun Cases

**Case 4:** Every muffin is cut in 2 pieces, every student gets 4 or 5 pieces. Number of pieces: 22. Note  $\leq 11$  pieces are  $> \frac{1}{2}$ .

- ▶  $s_4$  is number of students who get 4 pieces
- ▶  $s_5$  is number of students who get 5 pieces

$$4s_4 + 5s_5 = 22$$

$$s_4 + s_5 = 5$$

$s_4 = 3$ : There are 3 students who have 4 pieces.

$s_5 = 2$ : There are 2 students who have 5 pieces.

$$f(11, 5) = \frac{13}{30}, \text{ Fun Cases}$$

$$\begin{array}{ccccc} \diamond & \diamond & \diamond & \diamond & \diamond & (\text{Sums to } 11/5) \\ \diamond & \diamond & \diamond & \diamond & \diamond & (\text{Sums to } 11/5) \end{array}$$

$$\begin{array}{ccccc} \circ & \circ & \circ & \circ & (\text{Sums to } 11/5) \\ \circ & \circ & \circ & \bigcirc & (\text{Sums to } 11/5) \\ \circ & \circ & \circ & \bigcirc & (\text{Sums to } 11/5) \end{array}$$

**Case 4.1:** One of (say)

$$\circ \quad \circ \quad \circ \quad \bigcirc \quad (\text{Sums to } 11/5)$$

is  $\leq \frac{1}{2}$ . Then there is a piece

$$\geq \frac{(11/5) - (1/2)}{3} = \frac{17}{30}.$$

The other piece from the muffin is

$$\leq 1 - \frac{17}{30} = \frac{13}{30} \quad \text{Great to see } \frac{13}{30}.$$

$$f(11, 5) = \frac{13}{30}, \text{ Fun Cases}$$

### Case 4.2: All

○	○	○	○	(Sums to 11/5)
○	○	○	○	(Sums to 11/5)
○	○	○	○	(Sums to 11/5)

are  $> \frac{1}{2}$ .

There are  $\geq 12$  pieces  $> \frac{1}{2}$ . Can't occur.

## The Techniques Generalizes- We Do $f(24, 11)$

The technique for  $f(11, 5) \leq \frac{13}{30}$  has a generalization with a baker's dozen subcases. We do one concrete example:

**Definition:** Assume we have a protocol where all muffins are cut into two pieces. If  $x$  is a piece then the other piece in the muffin it came from is its **buddy**. Note that  $B(x) = 1 - x$ .

$$f(24, 11) \leq \frac{19}{44}$$

**Theorem:**  $f(24, 11) \leq \frac{19}{44}$  ( $\geq$  also known)

Assume  $(24, 11)$ -procedure with smallest piece  $> \frac{19}{44}$ .

Can assume all muffin cut in two and all student gets  $\geq 2$  shares.

We show that there is a piece  $\leq \frac{19}{44}$ .

**Case 1:** A student gets  $\geq 6$  shares. Some piece  $\leq \frac{24}{11 \times 6} < \frac{19}{44}$ .

**Case 2:** A student gets  $\leq 3$  shares. Some piece  $\geq \frac{24}{11 \times 3} = \frac{8}{11}$ .

Buddy of that piece  $\leq 1 - \frac{8}{11} \leq \frac{3}{11} < \frac{19}{44}$ .

**Case 3:** Every muffin is cut in 2 pieces and every student gets either 4 or 5 shares. Total number of shares is 48.

## How many students get 4? 5? Where are the Shares?

Let  $s_4$  ( $s_5$ ) be the number of 4-students (5-students).

$$4s_4 + 5s_5 = 48$$

$$s_4 + s_5 = 11 \quad \text{Get } s_4 = 7 \text{ and } s_5 = 4$$

**Case 3.1:**  $(\exists)$  4-sh  $\leq \frac{21}{44}$ . Rm. Now: 3 shares  $\geq \frac{24}{11} - \frac{21}{44}$ .  $(\exists)$  share

$$\geq \frac{(24/11) - (21/44)}{3} = \frac{25}{44}.$$

Buddy is

$$\leq 1 - \frac{25}{44} = \frac{19}{44}.$$

SO can assume all 4-shares are  $> \frac{21}{44}$ .

By similar reasoning:

**Case 3.2:** 4-shares in  $(\frac{21}{44}, \frac{25}{44})$ , 5-shares in  $(\frac{19}{44}, \frac{20}{44})$ .

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \text{ shs} \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 28 \text{ 4-shs} \\ \frac{25}{44} \end{array} \right)$$

## More Refined Picture of What is Going On

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \text{ shs} \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 28 \text{ 4-shs} \\ \frac{21}{44} \end{array} \right) \left[ \begin{array}{c} \\ \frac{25}{44} \end{array} \right)$$

**Claim 1:** There are no shares  $x \in [\frac{23}{44}, \frac{24}{44}]$ .

If there was such a share then  $B(x) \in [\frac{20}{44}, \frac{21}{44}]$ .

The following picture captures what we know so far.

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 8 \text{ S4-shs} \\ \frac{21}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{23}{44} \end{array} \right] \left( \begin{array}{c} 20 \text{ L4-shs} \\ \frac{24}{44} \end{array} \right) \left[ \begin{array}{c} \\ \frac{25}{44} \end{array} \right)$$

$$\left( \begin{array}{c} 20 \text{ 5-shs} \\ \frac{19}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{20}{44} \end{array} \right] \left( \begin{array}{c} 8 \text{ S4-shs} \\ \frac{23}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{21}{44} \end{array} \right] \left( \begin{array}{c} 20 \text{ L4-shs} \\ \frac{24}{44} \end{array} \right) \left[ \begin{array}{c} 0 \\ \frac{25}{44} \end{array} \right]$$

**Claim 2:** Every 4-student has at least 3 L4 shares.

If a 4-student had  $\leq 2$  L4 shares then he has

$$< 2 \times \left( \frac{23}{44} \right) + 2 \times \left( \frac{25}{44} \right) = \frac{24}{11}.$$

**Contradiction:** There are at least  $3 \times s_4 = 3 \times 7 = 21$  L4 shares.  
But there are only 20.



## What Else do we Have -Concrete

1. Formulas for  $f(m, 6)$  and  $f(m, 7)$ .
2. For  $s = 8, \dots, 100$  conjectures for  $f(m, s)$ .  $f(m, s)$  seems to be a mod  $s$  pattern.
3. Formulas for  $f(s + 1, s)$ ,  $f(s + 2, s)$ ,  $f(s + 3, s)$ ,  $f(s + 4, s)$ .  $f(s + d, s)$  seems to have a mod  $3d$  pattern.
4. A computer program that, on input  $m, s$  uses our theorems to find  $\alpha$  with  $f(m, s) \leq \alpha$  and then tries to prove  $f(m, s) \geq \alpha$  using linear algebra.
5. For  $1 \leq m, s \leq 50$  have all  $f(m, s)$  (Need to check that.)
6. Mixed integer program that always solves the problem but it is slow and has not been that useful.

## What Else do we have-Abstract

1. For fixed  $s$ , for  $m \geq \frac{s^3+2s^2+s}{2}$   $f(m, s)$  matches the Floor-ceiling bound.
2.  $f(m, s)$  always exists and is rational. Provable by compactness argument OR a large number of Linear Programs, OR one MIP. The last two proofs also give that  $f(m, s)$  is computable. Nice synergy – applied math tools helping us prove theorems in pure math!

# Open Problems-Complexity

Consider:

Given  $m, s$  in binary, compute  $f(m, s)$ .

1. Is the problem in P? We keep on finding techniques that we think cover all cases (so it would be in P) but then finding a case not covered.
2. Is it in NP? The procedure might be very large compared to the input.
3. Is it NP-complete or NP-hard?
4. Given  $m, s$  is there a bound on the denominators of the sizes of shares used?

## Open Problems-Misc

1. Show that for all  $m \geq s$ ,  $f(m, s) \geq \frac{1}{3}$ .
2. Prove that we ALWAYS get mod  $s$  pattern for  $f(m, s)$  (true for large enough  $m$ ).
3. Prove that we ALWAYS get mod  $3d$  pattern for  $f(s + d, s)$ .

# Accomplishment I Am Most Proud of

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Accomplishment I Am Most Proud of:

## Convinced

- ▶ 4 High School students (Guang, Naveen, Naveen, Sunny)
- ▶ 3 college student (Erik, Jacob, Daniel)
- ▶ 1 professor (John D)

that the most important field of Mathematics is **Muffinry**.