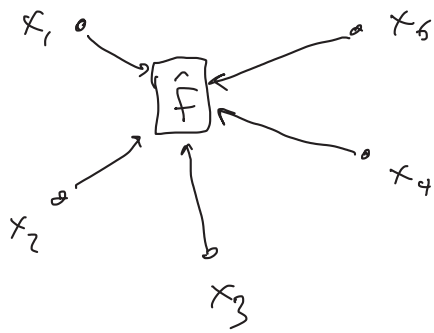
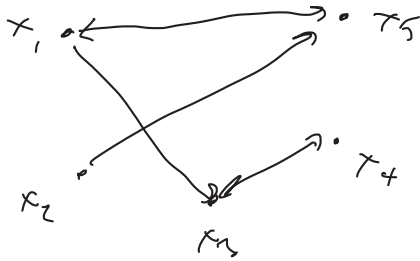


- scribes?
- lecture recording



- Want to compute $f(x_1, \dots, x_n) = \sum x_i$
- instead compute approximation $\hat{f}(\vec{x}) = f(\vec{x}) + \text{Lap}(\gamma \epsilon)$
 - to avoid a central authority, use MPC to compute \hat{f}

\Rightarrow protocol Π that computes a d.p. approx to f



Info-theoretic diff. privacy of Π :

For any set of τ parties & any neighboring inputs \vec{x}, \vec{x}' (that are equal for the τ corrupted), & any set of views V

$$\Pr(\text{View}_{\tau}^{\Pi}(\vec{x}) \in V) \leq e^{\epsilon} \cdot \Pr(\text{View}_{\tau}^{\Pi}(\vec{x}') \in V)$$

Computational version of diff. privacy of Π

for all efficient distinguishers D :

$$\Pr(D(\text{view}_{\tau}^{\Pi}(\vec{x})) = 1) \leq e^{\epsilon} \cdot \Pr(D(\text{view}_{\tau}^{\Pi}(\vec{x}')) = 1) + \underline{\underline{\delta(n)}}$$

Centralized protocol for summation: $\sum_i x_i + \text{Lap}(\gamma \epsilon)$

Local protocol for summation: $\sum_i (x_i + \text{Lap}(\gamma \epsilon))$

(un) Computationally, D.P. protocol for summation:

- Parties set up a threshold homomorphic encryption scheme
 - public key pk
 - given $Enc_{pk}(x_1), Enc_{pk}(x_2) \Rightarrow Enc_{pk}(x_1 \oplus x_2)$
 - threshold: every party holds a share sk_i of secret key sk
- Every party sets $\hat{x}_i = x_i + \text{noise}$
publish $Enc_{pk}(\hat{x}_i)$

- Parties compute $Enc_{pk}(\sum \hat{x}_i)$

- Parties collectively decrypt to get $\sum \hat{x}_i$

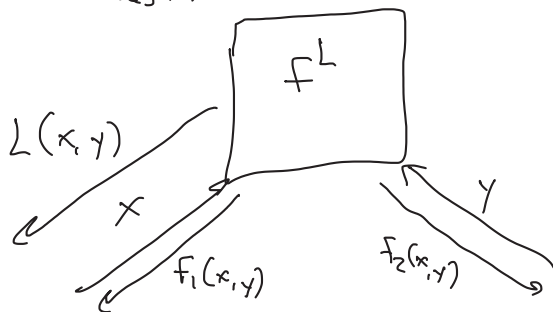
\Rightarrow noise per party can be much lower than in the LDP

Here: use MPC to compute a diff private functionality F

Another possibility: Could we a (ϵ, δ) -DP MPC protocol to compute F

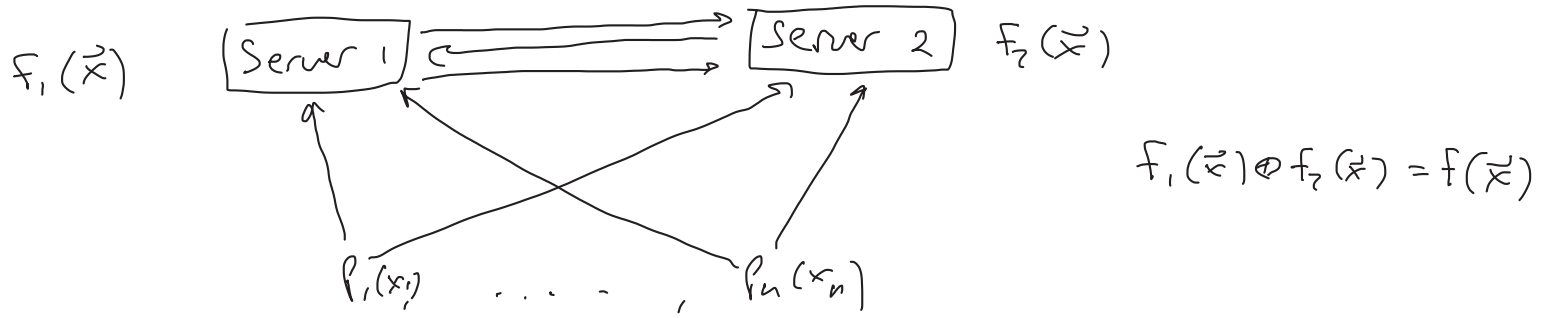
3rd possibility: use a (ϵ, δ) -DP MPC protocol to compute F

Semi-honest:

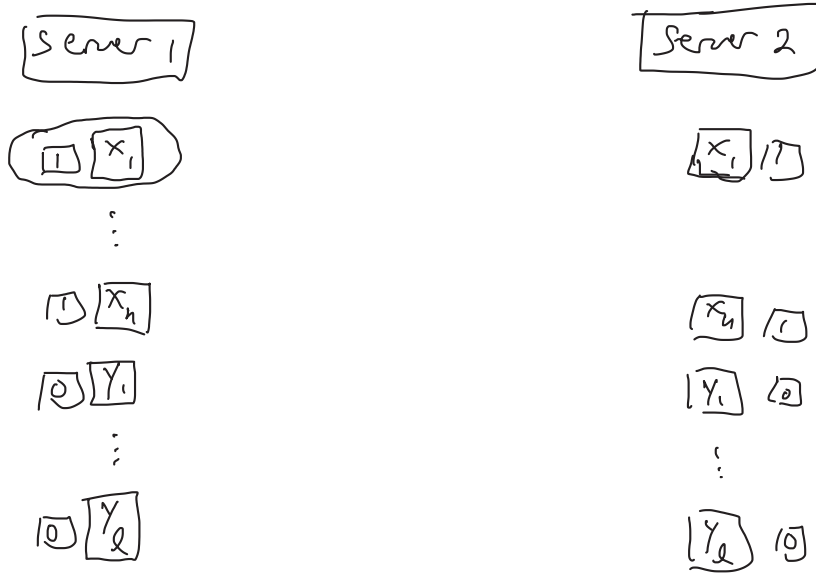


Protocol ϵ -diff privately computes F, f :

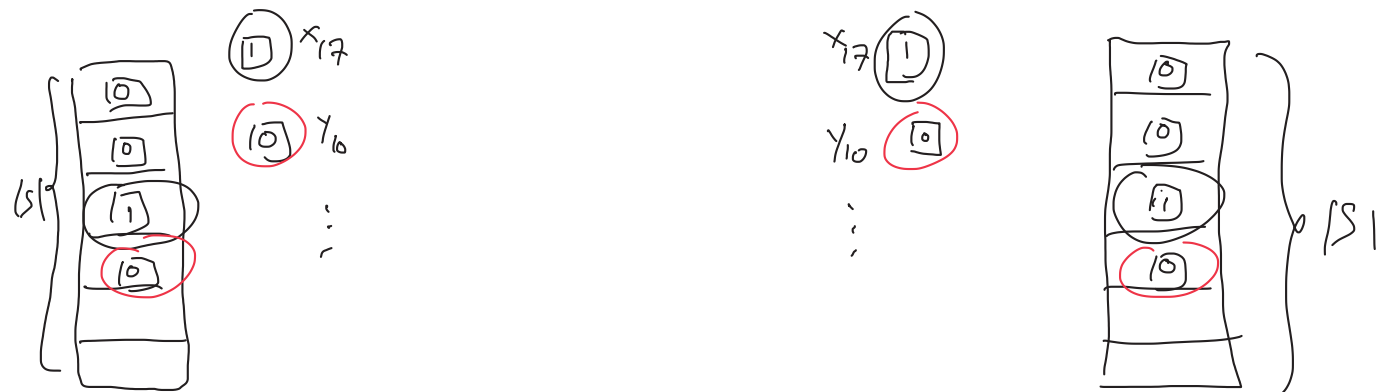
- π securely computes F^L
- L is ϵ -d.p.



$$f_1(\vec{x}) \oplus f_2(\vec{x}) = f(\vec{x})$$

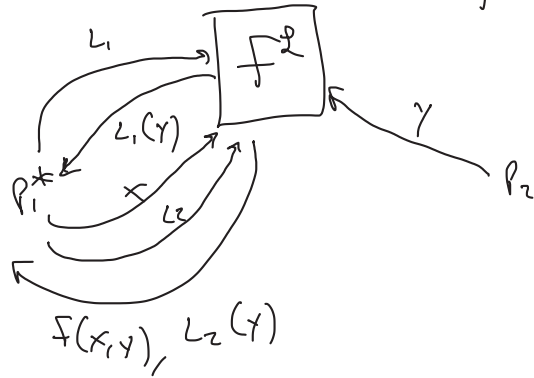


↓ randomly permute



diff. private computation, malicious case

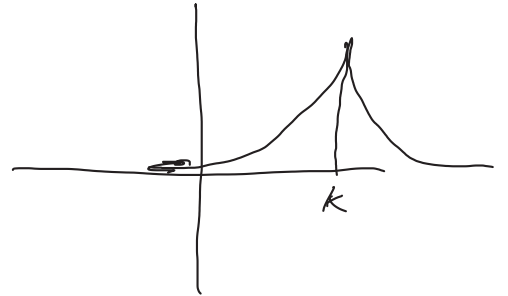
$$\mathcal{L} = \{ (L_1, L_2) \}$$



Protocol π is ϵ -dp. if

- π securely realizes $f^{\mathcal{L}}$

- Every $(L_1, L_2) \in \mathcal{L}$ is Σ -diff. private



$$\underline{P_1}(x_1, \dots, x_n)$$

$$\underline{P_2}(y_1, \dots, y_n)$$

$$d'_2, d'_1, x_1, x_{12}, x_{25} \longleftrightarrow y_1, y_2, d_1, d_2.$$

$$x_2, x_7, x_{20}$$

$$y_0, y_{21}, y_{30}$$

\vdots

y_3