# CMSC216: Binary Floating Point Numbers 

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## Logistics

## Reading Bryant/O'Hallaron

- Ch 2.1-3: Integers
- Ch 2.4-5: Floats
- 2021 Quick Guide to GDB


## Goals

- Finish Ints / Bitwise Ops
- gdb introduction
- Floating Point layout
- Thu: Assembly


## Assignments

- Lab05: Bits and GDB
- Minor bugs in Makefile and Quiz code line numbers
- HW05: Assembly Intro
- Project 2: Bitwise Ops, GDB, C Application
- Delayed for release, out later today
- 10 days to compelte

Grading on Exam 1 / Project 1 ongoing, release grades towards end of week

## Don't Give Up, Stay Determined!



- If Project 1 / Exam 1 went awesome, count yourself lucky
- If things did not go well, Don't Give Up
- Spend some time contemplating why things didn't go well, talk to course staff about it, learn from any mistakes
- There is a LOT of semester left and plenty of time to recover from a bad start


## GDB: The GNU Debugger

- P2 will include a "debugging problem" called puzzlebox
- Easiest to solve this problem using GDB (or some other debugger)
- Debuggers allow one to stop time in a program, inspect variables, pause execution at certain points and skip forwards
- If you've added tons of printf ()'s to your code and still can't figure out what's going on, a Debugger is your next option
- Phase 03 of puzzlebox is part of Lab05, makes a good demo of GDB basics
- Associated Reading: 2021 Quick Guide to GDB


## Note on Float Coverage

- Floating point layout is complex and interesting but...
- It's not a core topic that will appear on any exams, only tangentially on assignments
- Our coverage will be brief, examine slides / textbook if you want more depth
- GOAL: Demonstrate that (1) Real numbers can be approximated and (2) doing so uses bits in a very different way than integer representations


## Parts of a Fractional Number

The meaning of the "decimal point" is as follows:

$$
\begin{array}{rlrl}
123.406_{10}= & 1 \times 10^{2}+2 \times 10^{1}+3 \times 10^{0}+ & 123 & =100+20+3 \\
& 4 \times 10^{-1}+0 \times 10^{-2}+6 \times 10^{-3} & 0.406=\frac{4}{10}+\frac{6}{1000} \\
= & 123.406_{10} & &
\end{array}
$$

Changing to base 2 induces a "binary point" with similar meaning:

$$
\begin{array}{rlrl}
110.101_{2}= & 1 \times 2^{2}+1 \times 2^{1}+0 \times 2^{0}+ & 6 & =4+2 \\
& 1 \times 2^{-1}+0 \times 2^{-2}+1 \times 2^{-3} & 0.625=\frac{1}{2}+\frac{1}{8} \\
= & 6.625_{10} & &
\end{array}
$$

One could represent fractional numbers with a fixed point e.g.

- 32 bit fractional number with
- 10 bits left of Binary Point (integer part)
- 22 bits right of Binary Point (fractional part)

BUT most applications require a more flexible scheme

## Scientific Notation for Numbers

"Scientific" or "Engineering" notation for numbers with a fractional part is

| Standard | Scientific | printf("\%.4e",x); |
| ---: | ---: | :--- |
| 123.456 | $1.23456 \times 10^{2}$ | $1.2346 \mathrm{e}+02$ |
| 50.01 | $5.001 \times 10^{1}$ | $5.0010 \mathrm{e}+01$ |
| 3.14159 | $3.14159 \times 10^{0}$ | $3.1416 \mathrm{e}+00$ |
| 0.54321 | $5.4321 \times 10^{-1}$ | $5.4321 \mathrm{e}-01$ |
| 0.00789 | $7.89 \times 10^{-3}$ | $7.8900 \mathrm{e}-03$ |

- Always includes one non-zero digit left of decimal place
- Has some significant digits after the decimal place
- Multiplies by a power of $\mathbf{1 0}$ to get actual number

Binary Floating Point Layout Uses Scientific Convention

- Some bits for integer/fractional part
- Some bits for exponent part
- All in base 2: 1's and 0's, powers of 2


## Conversion Example

Below steps convert a decimal number to a fractional binary number equivalent then adjusts to scientific representation.
float fl = -248.75;

$$
\begin{aligned}
-248.75= & -(128+64+32+16+8+0+0+0) \cdot(1 / 2+1 / 4) \\
= & -11111000.11 * 2^{\wedge} 0 \\
& 76543210 \\
= & -12
\end{aligned}
$$

MANTISSA EXPONENT
$=-1.111100011 * 2 へ 7$
0123456789
Mantissa $\equiv$ Significand $\equiv$ Fractional Part

## Principle and Practice of Binary Floating Point Numbers

- In early computing, computer manufacturers used similar principles for floating point numbers but varied specifics
- Example of Early float data/hardware
- Univac: 36 bits, 1 -bit sign, 8 -bit exponent, 27 -bit significand ${ }^{1}$
- IBM: 32 bits, 1-bit sign, 7 -bit exponent, 24 -bit significand ${ }^{2}$
- Manufacturers implemented circuits with different rounding behavior, with/without infinity, and other inconsistencies
- Troublesome for reliability: code produced different results on different machines
- This was resolved with the adoption of the IEEE 754 Floating Point Standard which specifies
- Bit layout of 32-bit float and 64-bit double
- Rounding behavior, special values like Infinity
- Turing Award to William Kahan for his work on the standard

[^0]
## IEEE 754 Format: The Standard for Floating Point

| float | double | Property |
| ---: | ---: | :--- |
| 32 | 64 | Total bits |
| 1 | 1 | Bits for sign (1 neg / 0 pos) |
| 8 | 11 | Bits for Exponent multiplier (power of 2) |
| 23 | 52 | Bits for Fractional part or mantissa |
| 7.22 | 15.95 | Decimal digits of accuracy ${ }^{3}$ |

- Most commonly implemented format for floating point numbers in hardware to do arithmetic: processor has physical circuits to add/mult/etc. for this bit layout of floats
- Numbers/Bit Patterns divided into three categories

| Category | Description | Exponent |
| :--- | :--- | :--- |
| Normalized | most common like 1.0 and -9.56 e 37 | mixed 0/1 |
| Denormalized | very close to zero and 0.0 | all 0's |
| Special | extreme/error values like Inf and NaN | all 1's |

[^1]
## Example float Layout of -248.75: float_examples.c

FLOATING POINT FORMAT IEEE-754, 32 BITS


SIGN BIT
1 = NEGATIVE $0=$ POSITIVE

EXAMPLE: -248.75
HEXADECIMAL: C3 78 C0 00

Source: IEEE-754 Tutorial, www.puntoflotante.net
Color: 8-bit blocks, Negative: highest bit, leading 1

Exponent: high 8 bits, $2^{7}$ encoded with bias of -127

1000_0110-0111_1111
$=128+4+2-127$
$=134-127$
$=7$

Fractional/Mantissa portion is

```
1.111100011...
- ||||||||
    | explicit low 23 bits
|
implied leading 1
not in binary layout
```


## Normalized Floating Point: General Case

- A "normalized" floating point number is in the standard range for float/double, bit layout follows previous slide
- Example: $-248.75=-1.111100011 * 2 へ 7$

Exponent is in Bias Form (not Two's Complement)

- Unsigned positive integer minus constant bias number
- Consequence: exponent of 0 is not bitstring of 0 's
- Consequence: tiny exponents like -125 close to bitstring of 0 's; this makes resulting number close to 0
- 8-bit exponent $10000110=128+4+2=134$ so exponent value is $134-127=7$


## Integer and Mantissa Parts

- The leading 1 before the binary point is implied so does not show up in the bit string
- Remaining fractional/mantissa portion shows up in the low-order bits


## Fixed Bit Standards for Floating Point

## IEEE Standard Layouts

| Kind | Sign | Exponent |  |  | Mantissa |
| :--- | :--- | :--- | ---: | :--- | :--- |
|  | Bit | Bits | Bias | Exp Range | Bits |
| float | $31(1)$ | $30-23(8$ bits $)$ | -127 | -126 to +127 | $22-0(23$ bits $)$ |
| double | $63(1)$ | $62-52(11$ bits $)$ | -1023 | -1022 to +1023 | $51-0(52$ bits $)$ |

Standard allows hardware to be created that is as efficient as possible to do calculation on these numbers

## Consequences of Fixed Bits

- Since a fixed \# of bit is used, some numbers cannot be exactly represented, happens in any numbering system:
- Base 10 and Base 2 cannot represent $\frac{1}{3}$ in finite digits
- Base 2 cannot represent $\frac{1}{10}$ in finite digits

```
float f = 0.1;
```

printf("0.1 = \%.20e\n", f);
$0.1=1.00000001490116119385 \mathrm{e}-01$

Try show_float.c to see this in action

## Exercise: Quick Checks

1. What distinct parts are represented by bits in a floating point number (according to IEEE)
2. What is the "bias" of the exponent for 32-bit floats
3. Represent 7.125 in binary using "binary point" notation
4. Lay out 7.125 in IEEE-754 format
5. What does the number 1.0 look like as a float?

FLOATING POINT FORMAT IEEE-754, 32 BITS


SIGN BIT
1 = NEGATIVE
$0=$ POSITIVE

Source: IEEE-754 Tutorial, www.puntoflotante.net
The diagram above may help in recalling IEEE 754 layout

## Special Cases: See float_examples.c

## Special Values

- Infinity: exponent bits all 1, fraction all 0 , sign bit indicates $+\infty$ or $-\infty$
- Infinity results from overflow/underflow or certain ops like float $x=1.0 / 0.0$;
- \#include <math.h> gets macro INFINITY and -INFINITY
- NaN: not a number, exponent bits all 1, fraction has some 1s
- Errors in floating point like $0.0 / 0.0$


## Denormalized values: Exponent bits all 0

- Fractional/Mantissa portion evaluates without implied leading one, still an unsigned integer though
- Exponent is Bias $+1: 2^{-126}$ for float
- Result: very small numbers close to zero, smaller than any other representation, degrade uniformly to 0
- Zero: bit string of all 0s, optional leading 1 (negative zero);


## Other Float Notes



Source: XKCD \#217

## Approximations and Roundings

- Approximate $\frac{2}{3}$ with 4 digits, usually 0.6667 with standard rounding in base 10
- Similarly, some numbers cannot be exactly represented with fixed number of bits: $\frac{1}{10}$ approximated
- IEEE 754 specifies various rounding modes to approximate numbers


## Clever Engineering

- IEEE 754 allows floating point numbers to sort using signed integer sorting routines
- Bit patterns for float follows are ordered nearly the same as bit patterns for signed int
- Integer comparisons are usually fewer clock cycles than floating comparisons


## Sidebar: The Weird and Wonderful Union

- Bitwise operations like \& are not valid for float/double
- Can use pointers/casting to get around this OR...
- Use a union: somewhat unique construct to $C$
- Defined like a struct with several fields
- BUT fields occupy the same memory location (!?!)
- Allows one to treat a byte position as multiple different types, ex: int / float / char []
- Memory size of the union is the max of its fields

```
// union.c
typedef union { // shared memory
    float fl; // an float
    int in; // a int
    char ch[4]; // char array
} flint_t; // 4 bytes total
int main(){
    flint_t flint;
    flint.in = 0xC378C000;
    printf("%.4f\n", flint.fl);
    printf("%08x %d\n",flint.in,flint.in);
    for(int i=0; i<4; i++){
        unsigned char c = flint.ch[i];
        printf("%d: %02x '%c'\n",i,c,c);
    }
}
\begin{tabular}{|c|c|c|}
\hline | Symbol & Mem & Val I \\
\hline | flint.ch[3] & \#1027 & 0xC3 \\
\hline | flint.ch[2] & \#1026 & 0x78 \\
\hline | flint.ch[1] & | \#1025 & \(0 \times \mathrm{CO}\) \\
\hline | flint.in/fl/ch[0] & \#1024 & 0x00 \\
\hline | i & \#1020 & ? \\
\hline
\end{tabular}
```


## Floating Point Operation Efficiencies

- Floating Point Operations per Second, FLOPS is a major measure for numerical code/hardware efficiency
- Often used to benchmark and evaluate scientific computer resources, (e.g. top super computers in the world)
- Tricky to evaluate because of
- A single FLOP (add/sub/mul/div) may take 3 clock cycles to finish: latency 3
- Another FLOP can start before the first one finishes: pipelined
- Enough FLOPs lined up can get average 1 FLOP per cycle
- FP Instructions may automatically operate on multiple FPs stored in memory to feed pipeline: vectorized ops
- Generally referred to as superscalar
- Processors schedule things out of order too
- All of this makes micro-evaluation error-prone and pointless
- Run a real application like an N -body simulation and compute

$$
\text { FLOPS }=\frac{\text { number of floating ops done }}{\text { time taken in seconds }}
$$

## Top 5 Super Computers Worldwide, June 2023

| Rank | System | \#Cores | $\begin{aligned} & \mathrm{Rmax} \\ & \text { (TFlop/s) } \\ & \hline \end{aligned}$ | $\begin{array}{r} \text { Rpeak } \\ \text { (TFlop/s) } \end{array}$ | Power* <br> (kW) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | Frontier, USA / Oak Ridge Cray EX235a, AMD EPYC 2 GHz (×86-64) | 8,699,904 | 1,194.00 | 1,679.82 | 22,703 |
| 2 | Fugaku, Japan / Fujitsu Fujitsu A64FX 2.2GHz (Arm) | 7,630,848 | 442,010.0 | 537.21 | 29,899 |
| 3 | LUMI Finland / EuroHPC Cray EX235a, AMD EPYC 2 GHz (×86-64) | 2,220,288 | 309.10 | 428.70 | 6,016 |
| 4 | Leonardo Italy / EuroHPC | 1,824,768 | 238.70 | 304.47 | 7,404 |
| 5 | Summit United States IBM POWER9 22C 3.07 GHz (Power) | 2,414,592 | 148,600.0 | 200,794.9 | 10,096 |

https://www.top500.org/lists/top500/2022/06/
*: An average US Home uses 909 kWh of power per month

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| Rank | System | \#Cores | Rmax <br> (TFlop/s) | Rpeak <br> (TFlop/s) | Power* <br> (kW) |
| ---: | :--- | :--- | :--- | :--- | :--- |
| 1 | Frontier, USA / Oak Ridge <br> Cray EX235a, AMD EPYC 2GHz <br> (x86-64) | $8,730,112$ | $1,102.00$ | $1,685.65$ | 21,100 |
| 2 | Fugaku, Japan / Fujitsu <br> Fujitsu A64FX 2.2GHz <br> (Arm) | $7,630,848$ | $442,010.0$ | $537,212.0$ | 29,899 |
| 3 | LUMI Finland / EuroHPC <br> Cray EX235a, AMD EPYC 2GHz <br> (x86-64) | $1,110,144$ | 151.90 | 214.35 | 2,942 |
| 4Summit United States <br> IBM POWER9 22C 3.07GHz <br> (Power) | $2,414,592$ | $148,600.0$ | $200,794.9$ | 10,096 |  |
| 5 | Sierra United States <br> IBM POWER9 22C 3.1GHz <br> (Power) | $1,572,480$ | $94,640.0$ | $125,712.0$ | 7,438 |

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## Top 5 Super Computers Worldwide, June 2021

| Rank | System | \#Cores | Rmax (TFlop/s) | Rpeak (TFlop/s) | Power (kW) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | Fugaku, Japan / Fujitsu Fujitsu A64FX 2.2GhZ (Arm) | 7,630,848 | 442,010.0 | 537,212.0 | 29,899 |
| 2 | Summit United States IBM POWER9 22C 3.07GHz (Power) | 2,414,592 | 148,600.0 | 200,794.9 | 10,096 |
| 3 | Sierra United States IBM POWER9 22C 3.1GHz (Power) | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |
| 4 | Sunway TaihuLight China <br> Sunway SW26010 <br> (custom RISC) | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 5 | Perlmutter, United States AMD EPYC 2.45 GHz , Cray (×86-64) | 706,304 | 64,590.0 | 89,794.5 | 2,528 |

https://www.top500.org/lists/top500/2021/06/

## Top 5 Super Computers Worldwide, Nov 2020

| Rank | System | \#Cores | Rmax <br> (TFlop/s) | Rpeak <br> (TFlop/s) | Power <br> (kW) |
| ---: | :--- | :--- | :--- | :--- | :--- |
| 1 | Fugaku, Japan / Fujitsu <br> Fujitsu A64FX 2.2GhZ <br> (Arm) | $7,299,072$ | $415,530.0$ | $513,854.7$ | 28,335 |
| 2 | Summit United States <br> IBM POWER9 22C 3.07GHz <br> (Power) | $2,397,824$ | $143,500.0$ | $200,794.9$ | 10,096 |
| 3 | Sierra United States <br> IBM POWER9 22C 3.1GHz <br> (Power) | $1,572,480$ | $94,640.0$ | $125,712.0$ | 7,438 |
| 4 | Sunway TaihuLight China <br> Sunway SW26010 <br> (custom RISC) | $10,649,600$ | $93,014.6$ | $125,435.9$ | 15,371 |
| 5 | Selene USA, NVIDIA/AMD <br> AMD EPYC 7742 64C 2.25GHz <br> (x86-64) | 555,520 | $63,460.0$ | $79,215.0$ | 2,646 |

https://www.top500.org/lists/top500/2020/06/

## Top 5 Super Computers Worldwide, June 2020

| Rank | System | \#Cores | Rmax (TFlop/s) | Rpeak (TFlop/s) | Power (kW) |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | Fugaku, Japan / Fujitsu Fujitsu A64FX 2.2GhZ (Arm) | 7,299,072 | 415,530.0 | 513,854.7 | 28,335 |
| 2 | Summit United States IBM POWER9 22C 3.07GHz (Power) | 2,397,824 | 143,500.0 | 200,794.9 | 10,096 |
| 3 | Sierra United States IBM POWER9 22C 3.1GHz (Power) | 1,572,480 | 94,640.0 | 125,712.0 | 7,438 |
| 4 | Sunway TaihuLight China <br> Sunway SW26010 <br> (custom RISC) | 10,649,600 | 93,014.6 | 125,435.9 | 15,371 |
| 5 | Tianhe-2A China Intel Xeon 2.2 GHz (×86-64) | 4,981,760 | 61,444.5 | 100,678.7 | 18,482 |

https://www.top500.org/lists/top500/2020/06/

## Top 5 Super Computers Worldwide, Nov 2019

| Rank | System | \#Cores | Rmax <br> (TFlop/s) | Rpeak <br> (TFlop/s) | Power <br> (kW) |
| ---: | :--- | :--- | :--- | :--- | :--- |
| 1 | Summit United States <br> IBM POWER9 22C 3.07GHz | $2,397,824$ | $143,500.0$ | $200,794.9$ | 9,783 |
| 2 | Sierra United States <br> IBM POWER9 22C 3.1GHz, | $1,572,480$ | $94,640.0$ | $125,712.0$ | 7,438 |
| 3 | Sunway TaihuLight China <br> Sunway MPP | $10,649,600$ | $93,014.6$ | $125,435.9$ | 15,371 |
| 4 | Tianhe-2A China <br> Xeon 2.2GHz | $4,981,760$ | $61,444.5$ | $100,678.7$ | 18,482 |
| 5 | Frontera, United States <br> Dell 6420, Xeons 2.7GHz | 448,448 | $23,516.4$ | $38,745.9$ | $? ?$ |

https://www.top500.org/list/2019/11/

## Top 5 Super Computers Worldwide, Nov 2018

| Rank | System | \#Cores | Rmax <br> (TFlop/s) | Rpeak <br> (TFlop/s) | Power <br> (kW) |
| ---: | :--- | :--- | :--- | :--- | :--- |
| 1 | Summit United States <br> IBM POWER9 22C 3.07GHz | $2,397,824$ | $143,500.0$ | $200,794.9$ | 9,783 |
| 2 | Sierra United States <br> IBM POWER9 22C 3.1GHz, | $1,572,480$ | $94,640.0$ | $125,712.0$ | 7,438 |
| 3 | Sunway TaihuLight China <br> Sunway MPP | $10,649,600$ | $93,014.6$ | $125,435.9$ | 15,371 |
| 4 | Tianhe-2A China | $4,981,760$ | $61,444.5$ | $100,678.7$ | 18,482 |
|  | TH-IVB-FEP Cluster |  |  |  |  |
| 5 | Piz Daint Switzerland <br> Cray XC50, Xeon E5-2690v3 | 387,872 | $21,230.0$ | $27,154.3$ | 2,384 |

https://www.top500.org/list/2018/11/

## Top 5 Super Computers Worldwide, Nov 2017

| Rank | System | \#Cores | Rmax <br> (TFlop/s) | Rpeak <br> (TFlop/s) | Power <br> (kW) |
| ---: | :--- | :--- | :--- | :--- | :--- |
| 1 | Sunway TaihuLight China <br> Sunway MPP | $10,649,600$ | $93,014.6$ | $125,435.9$ | 15,371 |
| 2 | Tianhe-2 (MilkyWay-2) China | $3,120,000$ | $33,862.7$ | $54,902.4$ | 17,808 |
|  | TH-IVB-FEP Cluster |  |  |  |  |

https://www.top500.org/lists/2017/11/


[^0]:    ${ }^{1}$ Floating Point Arithmetic
    ${ }^{2}$ IBM Hexadecimal Floats

[^1]:    ${ }^{3}$ Wikipedia: IEEE 754

